

12 May 2021

Last time:

- Analysis of greedy algo via RIC.

Today:

- To show that Gaussian matrices satisfy RIP.

Recap: Two Lemmas

Lemma 1: $A \in \mathbb{R}^{m \times n}$, $A_{ij} \stackrel{iid}{\sim} \mathcal{N}(0, \frac{1}{n})$.

Then, for any fixed $x \in \mathbb{R}^n$,

$$\Pr \{ | \|Ax\|_2^2 - \|x\|_2^2 | \geq \epsilon \|x\|_2^2 \} \leq 2e^{-m c_\epsilon(\epsilon)}$$

$$\forall 0 < \epsilon < 1, \text{ where } c_\epsilon(\epsilon) = \frac{\epsilon^2}{8} - \frac{\epsilon^3}{6}.$$

[$c_\epsilon(\epsilon) = \frac{\epsilon^2}{12}$ also works].

Remark: $\mathbb{E} \{ \|Ax\|_2^2 \} = \|x\|_2^2$; $\text{var}(A_{ij}) = \frac{1}{n}$ measures the length of x in the expected sense.

Lemma 2: Let $\epsilon \in (0, 1)$ be given. Then, \exists a set of points ("net of points") $Q \subset \mathbb{R}^k$ s.t.

$$\|q\|_2 = 1 \quad \forall q \in Q, \text{ and } |Q| \leq \left(\frac{2}{1-\epsilon}\right)^k,$$

for any $x \in \mathbb{R}^k$, $\|x\|_2 = 1$, $\min_{q \in Q} \|x - q\|_2 \leq \epsilon$.

Remark: Upper b.d. on the min. # pts we need to take on the unit k -ball s.t. an arbitrary pt on the unit k -ball is within ϵ of at least one of the pts.

Thm. Let $0 < \delta < 1$, $A \in \mathbb{R}^{m \times n}$, $A_{ij} \stackrel{iid}{\sim} \mathcal{N}(0, \frac{1}{n})$.

If $m \geq \frac{2}{\delta^2} k \log \left(\frac{2}{\delta}\right)$, then A satisfies RIP of order k

with the prescribed δ with prob. exceeding

$$1 - 2e^{-\frac{\delta^2 m}{8} \left(\frac{\delta^2}{8}\right)^k}$$

where κ_1 is arbitrary and $\kappa_2 = \frac{1}{24} - \frac{1}{\kappa_1}$.

Remark: If $m = \frac{2}{\delta^2} k \log \left(\frac{2}{\delta}\right)$, the prob. becomes

$$1 - 2e^{-\frac{\delta^2 m}{8} \left(\frac{\delta^2}{8}\right)^k} = \left(\frac{\delta^2}{8}\right)^{-\left(\frac{\delta^2 m}{8} - 1\right)k} \left(\frac{\delta^2}{8}\right)^k = \left(\frac{\delta^2}{8}\right)^{\frac{\delta^2 m}{8} - k}$$

Remark: By choosing κ_1 large enough, can assume $\kappa_2 > 0$, and further, for large enough m , $e^{-\kappa_2 m} \left(\frac{\delta^2}{8}\right)^k \ll 1$ holds.

Proof: Define $\mathcal{I}_k = \{x \in \mathbb{R}^n, \|x\|_2 \leq k\}$.

Need t.s.t.

$$(1-\delta) \|x\|_2 \leq \|Ax\|_2 \leq (1+\delta) \|x\|_2 \quad \forall x \in \mathcal{I}_k$$

(1) Suffices to consider $\|x\|_2 = 1$.

(2) Fix $T \subset [n]$, $|T| = k$, and let $\mathcal{X}_T = \{x \in \mathbb{R}^n, \text{supp}(x) = T\} \subset \mathcal{I}_k$.

Choose a finite net of pts $Q_T \subset \mathcal{X}_T$ s.t. $\|q\|_2 = 1$

$\forall q \in Q_T$ and $\forall x \in \mathcal{X}_T$ with $\|x\|_2 = 1$,

$$\min_{q \in Q_T} \|x - q\|_2 \leq \frac{\delta}{14}.$$

Lemma 2 \Rightarrow Can choose Q_T s.t. $|Q_T| \leq \left(\frac{2}{1-\delta}\right)^k$.

Repeat $\binom{n}{k}$ choices for T , and collect Q_T together

$$Q = \bigcup_{T: |T|=k} Q_T$$

$$\binom{n}{k} = \frac{n(n-1) \dots (n-k+1)}{k!} \leq \frac{n^k}{k!} \leq \left(\frac{en}{k}\right)^k$$

$$\Rightarrow |Q| \leq \left(\frac{en}{k}\right)^k \left(\frac{2}{1-\delta}\right)^k = \left(\frac{4en}{k\delta}\right)^k$$

Lemma 1 \Rightarrow For each $q \in Q$, $\forall 0 < \epsilon < 1$,

$$\Pr \left\{ (1-\epsilon) \|q\|_2 \leq \|Aq\|_2 \leq (1+\epsilon) \|q\|_2 \right\} \geq \frac{1 - 2e^{-m c_\epsilon(\epsilon)}}{2}$$

$$\Rightarrow \text{Taking the union bound,} \\ (1-\delta) \|q\|_2 \leq \|Aq\|_2 \leq (1+\delta) \|q\|_2 \quad \forall q \in Q \\ \text{w.p. } \geq \frac{1 - 2e^{-m c_\epsilon(\epsilon)}}{2} \cdot \frac{1 - 2e^{-m c_\epsilon(\epsilon)}}{2}$$

$$\text{Let } \epsilon = \frac{\delta}{\sqrt{2}}. \\ \text{Also note } \log \left[\left(\frac{4en}{k\delta}\right)^k \right] \leq k \left(\log \left(\frac{en}{k}\right) + \log \frac{4}{\delta} \right)$$

$$\therefore m \geq \frac{2}{\delta^2} k \log \frac{4}{\delta} \leq \frac{2}{\delta^2} k \left(\log \frac{en}{k} + \log \frac{4}{\delta} \right)$$

$$1 - 2e^{-m c_\epsilon(\epsilon)} \geq \frac{1 - 2e^{-\frac{m \delta^2}{8}}}{2} \geq 1 - 2e^{-\frac{m \delta^2}{8}} \cdot e^{\frac{2}{\delta^2} k \log \frac{4}{\delta}}$$

$$\therefore c_\epsilon(\epsilon) = \frac{\epsilon^2}{12} \\ \geq 1 - 2 \left(\frac{4en}{k\delta}\right)^k e^{-m \left(\frac{\delta^2}{24} - \frac{1}{\kappa_1}\right)}$$

So we have the req'd. failure prob. But we need t.s.t. A satisfies RIP with the prescribed δ .

Let S_k be the smallest # s.t. $\|Ax\|_2 \leq \sqrt{1+\delta_k} \|x\|_2 \quad \forall x \in \sum_k, \|x\|_2 \leq 1$.

Need t.s.t. $S_k \leq \frac{2}{\delta}$.

Lemma 2 $\Rightarrow \forall x \in \sum_k, \|x\|_2 = 1$, can pick $q \in Q$ s.t. $\|x - q\|_2 \leq \frac{\delta}{14}$.

In fact, $q \in \sum_k$ as well ($\because \text{supp}(q) = T$, the $x = x_T \Rightarrow$ can pick $q \in Q_T$ s.t. $\|x - q\|_2 \leq \frac{\delta}{14}$, $\text{supp}(q) = T$ as well.)

$$\|Ax\|_2 \leq \|Aq\|_2 + \|A(x-q)\|_2 \quad (a_4)$$

$$\leq \sqrt{1+\delta} \|q\|_2 + \sqrt{1+\delta_k} \|x-q\|_2$$

$$\leq \left(\sqrt{1+\delta} + \sqrt{1+\delta_k} \frac{\delta}{14} \right)$$

Since, by defn. δ_k is the smallest # s.t.

$$\|Ax\|_2 \leq \sqrt{1+\delta_k} \|x\|_2 \quad \forall x \in \Sigma_k,$$

$$\sqrt{1+\delta_k} \leq \sqrt{1+\delta} + \sqrt{1+\delta_k} \frac{\delta}{14}.$$

$$\Rightarrow \sqrt{1+\delta_k} \leq \frac{\sqrt{1+\delta}}{1-\frac{\delta}{14}} \leq \sqrt{1+\delta} \Rightarrow \delta_k \leq \delta.$$

(Why we chose $\frac{\delta}{14}$)

For the lower bound,

$$\|Ax\|_2 \geq \|Aq\|_2 - \|A(x-q)\|_2$$

$$\geq \sqrt{1-\delta} - \sqrt{1+\delta} \frac{\delta}{14}.$$

$$\geq \sqrt{1-\delta}$$

Thus, given $0 < \delta < 1$, $A \in \mathbb{R}^{m \times n}$, $A_{ij} \stackrel{i.i.d.}{\sim} \mathcal{N}(0, \frac{1}{n})$,
 $\forall x \in \Sigma_k$, $(1-\delta)\|x\|_2^2 \leq \|Ax\|_2^2 \leq (1+\delta)\|x\|_2^2$

holds w.p. exceeding

$$1 - 2\left(\frac{2\delta}{\delta}\right)^k e^{-k_2 m}$$

where $k_2 = \frac{\delta^2}{24} - \frac{1}{k_1}$, provided $m \geq k_1 k \log\left(\frac{m}{k}\right)$.

□