

- Last time:
- Algo for l_1 min.
 - Coordinate descent ("shooting algo")
 - Homology based methods: LARS
 - Proximal and gradient projection methods

Today:

- Complete the discussion on proximal & gradient proj. methods
- Gelfand widths and lower bounds on the # meas. needed to satisfy RIP.

Proximal and Gradient Projection Methods

Objective fun: $f(x) = \underbrace{L(x)}_{\text{Liss}} + \underbrace{R(x)}_{\text{Rygotinjor}}$

To develop intuition, suppose $A=I$. Then,
 $f(x) = \frac{1}{2} \|y-x\|_2^2 + R(x)$.

Soln: $x_{opt} = \text{prox}_R(y) = \arg \min \left\{ \frac{1}{2} \|y-x\|_2^2 + R(x) \right\}$
proximal operator Minimize R, but minimize distance to y

Will use the above in an iterative algo:
 $\text{prox}_R(x_k) = \arg \min \left\{ \frac{1}{2} \|x_k - x\|_2^2 + R(x) \right\}$ - (2)
prox iterat.

- (a) How to solve the above pb?
 (b) How to extend to a general A .

(a) Proximal operator

If $R(x) = \lambda \|x\|_1$,
 $\text{prox}_R(x) = \text{soft}(x, \lambda)$ (component-wise)
 where $\text{soft}(a, \lambda) = \text{sign}(a) \max(|a| - \lambda, 0)$,
 $x_i = \max\{x_i, 0\}$.

If $R(x) = \lambda \|x\|_2$,
 $\text{prox}_R(x) = \text{hard}(x, \lambda/\|x\|_2)$
 where $\text{hard}(x, a) = \frac{x}{\|x\|_2} \max\{a, \|x\|_2\}$.

(b) Proximal gradient method
 $x_{k+1} = \arg \min \left\{ \frac{1}{2} \|R(x_k) + L(x_k) + g_k^T(x - x_k)\|_2^2 + \frac{\lambda}{2} \|x - x_k\|_2^2 \right\}$

⊙ Alternative: if $R \cdot g^T$ exist.
 For e.g., if we replace $\|x\|_2^2$ with a Bregman divergence, we get "mirror method".

⊙ No approximate $\nabla^2 L(x_k) \approx \frac{1}{\lambda} I$.

Thus, $x_{k+1} = \arg \min \left\{ \frac{1}{2} \|R(x_k) + \frac{1}{\lambda} \nabla L(x_k)\|_2^2 \right\}$
 $= \text{prox}_R(x_k)$

where $u_k = x_k - \frac{1}{\lambda} g_k$, $g_k = \nabla L(x_k)$.

If $R(x) = 0$, it is the gradient descent method.
 If $R(x) = \lambda \|x\|_1$, get iterative soft thresholding.

Choice of $\frac{1}{\lambda}$: natural choice $\left[\frac{1}{\lambda} I \approx \nabla^2 L \right]$
 $\frac{1}{\lambda} I \approx \frac{\nabla^2 L}{g(x_k) \cdot g(x_k)}$
 $\frac{1}{\lambda} \approx \frac{g_k^T g_k}{\|g_k\|_2^2}$

Hence, we compute
 $u_k = \arg \min_x \|x - (x_k - \frac{1}{\lambda} g_k)\|_2^2$
 $= \frac{(x_k - \frac{1}{\lambda} g_k)(g_k^T g_k - \frac{1}{\lambda})}{(g_k - \frac{1}{\lambda} g_k)^T (g_k - \frac{1}{\lambda} g_k)}$

and set $\frac{1}{\lambda} = \frac{1}{g_k^T g_k}$.
 This is known as the Bregman-Bregman (BB) or spectral stepsize method.

BB stepsize + iterative soft thresholding + continuation method (gradually reduce λ) \Rightarrow SparsSA [Wright 2007]
 Sparse reconstruction by separable approx.

Nesterov's Method

$$x_{k+1} = \text{prox}_R \left(\frac{g_k - \frac{1}{\lambda} g_k}{2} \right)$$

(with $x_k = \text{prox}_R(x_k - \frac{1}{\lambda} g_k)$)

$$g_k = \frac{1}{2} \left(\frac{g_k}{\|g_k\|_2} (x_k - x_{k-1}) \right)$$

$$g_k = \nabla L(x_k)$$

This + IST ($R(x) = \lambda \|x\|_1$) + Continuation \Rightarrow FISTA (fast iterative shrinkage & thresholding algo).

Chapter 10: Gelfand Widths of l_1 Balls

Defn. The Gelfand m -width of a subset K of a normed space X is

$$d^m(K, X) = \inf_{A \in \mathbb{R}^{m \times m}} \sup_{x \in K \cap \text{ker}(A)} \|x\|_X$$

$\{d^m(K, X)\}_{m=0}^N$ is nonincreasing.

$$d^0(K, X) = \sup_{x \in K} \|x\|_X$$

If K contains the all zero vector,
 $d^m(K, X) = 0 \ \forall \ m \geq N$.

Defn. The compressive m -width of a subset K of a normed space X is defined as

$$E^m(K, X) = \inf_{\substack{A: X \rightarrow \mathbb{R}^m \text{ linear} \\ \Delta: \mathbb{R}^m \rightarrow X}} \sup_{x \in K} \|Ax\|_X$$

Δ : Reconstruction map $\mathbb{R}^m \rightarrow X$, arbitrary (Could even be the NP-hard l_1 min. solver)
 A : nonadaptive $A = \begin{bmatrix} a_{11} & \dots & a_{1N} \\ \vdots & \ddots & \vdots \\ a_{m1} & \dots & a_{mN} \end{bmatrix}$ chosen once and for all.

In the context of CS, $K =$ set of compressible vecs.

Prop. 1.1 For any $\gamma > 0$, and any $x \in \mathcal{C}^m$,

$$g(x)_\gamma \leq \frac{1}{\sqrt{2}} \|x\|_1$$

\Rightarrow the nonconvex unit balls \mathcal{C}^m with $p < 1$ are good models for compressible vecs.

$\tau = \frac{1}{\sqrt{2}}$ if $1 < p < 2$, $\|x\|_p \leq 1$, then $\|x\|_1 \leq \tau \|x\|_p$.

$$\|g(x)\| \leq \frac{r}{\delta}$$

For simplicity, we consider $p=1$ here.

We have the following result:

Thm 105 For $1 < p \leq 2$ and $m < n$, \exists const. $c_1, c_2 > 0$ depending only on p s.t.
 $c_1 \min\left\{1, \frac{\lambda_1(\frac{\Sigma}{n})}{n}\right\}^{1/p} \leq d^*(x^*, x_p^*) \leq c_2 \min\left\{1, \frac{\lambda_1(\frac{\Sigma}{n})}{n}\right\}^{1/p}$

Matching upper and lower bounds.

Leads to this result on compressive widths:

Cor. 104 For $1 < p \leq 2$ and $m < n$, the compressive width satisfies
 $c_1 \min\left\{1, \frac{\lambda_1(\frac{\Sigma}{n})}{n}\right\}^{1/p} \leq c^*(\frac{\Sigma}{n}, p) \leq c_2 \min\left\{1, \frac{\lambda_1(\frac{\Sigma}{n})}{n}\right\}^{1/p}$

for some const. $c_1, c_2 > 0$ that dep. only on p .

Remark: We have seen that, if $m \geq c \lambda_1(\frac{\Sigma}{n})$,

there are matrices $A \in \mathbb{R}^{m \times n}$ with small RIC and

reconstruction maps A (e.g. SP, HT, OMP, etc) s.t.

$$\|x - \Delta(Ax)\|_p \leq \frac{c}{\lambda_1(\frac{\Sigma}{n})} \|x\|_p$$

[$m \geq c \lambda_1(\frac{\Sigma}{n})$ is sufficient]. Conversely, using

the above corollary, we can s.t. the existence of

A & Δ providing such a stability result in the

reconstruction error forces $m \geq c \lambda_1(\frac{\Sigma}{n})$ to

be satisfied:

Prop. 107 Let $1 < p \leq 2$. Suppose $\exists A \in \mathbb{R}^{m \times n}$

and $\Delta: \mathbb{R}^m \rightarrow \mathbb{R}^n$ s.t. $\forall x \in \mathbb{R}^n$

$$\|x - \Delta(Ax)\|_p \leq \frac{c}{\lambda_1(\frac{\Sigma}{n})} \|x\|_p$$

Then, for some const. $c_1, c_2 > 0$ dep. only on c ,

$$m \geq c_1 \lambda_1(\frac{\Sigma}{n}), \text{ provided } m \geq c_2.$$