

Harvesting Rate Optimized Transmit Power Control Policies in Wireless Energy Harvesting Sensors with Retransmissions

Adithya M. Devraj, Mohit K. Sharma, and Chandra R. Murthy

Signal Processing for Communications Lab
Dept. of ECE, IISc

December 21, 2013

Outline

- 1 System Model
- 2 Previous Works
- 3 Harvesting Rate Optimized Fixed Energy Transmission Scheme
 - Slow Fading
 - Fast Fading
- 4 Harvesting Rate Optimized Variable Transmit Energy Scheme
 - Slow Fading
 - Fast Fading
- 5 Finite Battery Capacity
 - Exact Modification
 - Heuristic Modification
- 6 Simulations and Results

Transmission Timeline

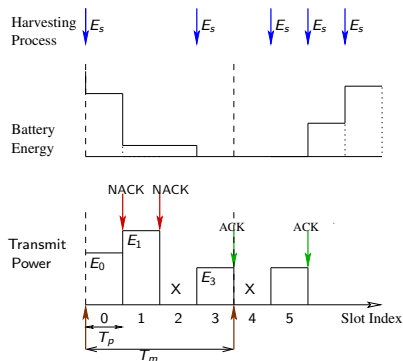


Figure : Transmission timeline of the EH node for $K = 4$, showing the random energy harvesting process (\downarrow) and periodic data arrival (\uparrow). The marker “X” denotes slots where the EHS does not transmit data

- Energy Injection Process:
 - Every slot, E_s energy is harvested with probability ρ
 - No energy is harvested with probability $1-\rho$
- Channel:
 - Slow fading: $T_c = T_m$
 - Fast fading: $T_c = T_p$
- Packet error probability:

$$\text{For BPSK: } P_e(E_i, \gamma) = 1 - \left(1 - Q \left(\sqrt{\frac{2\gamma E_i}{N_0}} \right) \right)^L \quad (1)$$

$$\text{Convex approximation: } P_e(E_i, \gamma) \approx \exp \left(-\frac{E_i \gamma}{N_0} \right) \quad (2)$$

- Feedback:
 - If packet is in error: (NAK) is sent
 - If packet is successfully decoded: (ACK) is sent
- Performance metric:

$$\text{Outage probability} = \frac{\text{number of packets discarded}}{\text{number of frames}} \quad (3)$$

Previous Works

- Heuristic policies
- Decision-theoretic policies (using dynamic programming):
 - Completely observable MDP
 - Partially observable MDP (POMDP)
 - Channel independent MDP
- Issues with decision-theoretic policies:
 - Hard to compute; except under certain conditions
 - Accurate knowledge of the battery energy level
 - Lack of insight into the structure of optimal solution

This Work

- *Harvesting rate optimized fixed energy transmission scheme*
 - Optimal fixed energy transmission scheme for a given ρ
 - Transmit energy fixed through all K slots for all the frames
- *Harvesting rate optimized variable transmit energy scheme*
 - Transmit energy varies from one retransmission index to the other
 - The transmit energy vector for a given ρ remains same throughout

Infinite Battery Capacity Assumption

- Replace **instantaneous power constraint** with an **average energy constraint**
 - Valid only when we have an infinite energy buffer
 - Completely absorbs the randomness in energy harvesting
- As long as $\sum_{k=1}^K E_k \ll B_{\max}$, infinite battery capacity is a reasonable approximation

Harvesting Rate Optimized Fixed Energy Transmission Scheme

- Transmit energy is fixed through out the transmission process
- Every slot, node harvests E_s with probability ρ
- Average energy harvested per frame: $\bar{E}_f = K\rho E_s$
- Fixed transmit energy: $E_{tx} = \epsilon K\rho E_s$
 - Find the optimum ϵ value

Slow Fading

- Outage probability:

$$P_{\text{out}} = \mathbb{E}_{\gamma} \left\{ P_e^K(\epsilon K \rho E_s, \gamma) \right\} \quad (4)$$

- Average energy used per frame:

$$\begin{aligned} \bar{E}_{\text{tx}} = & \mathbb{E}_{\gamma} \left\{ \epsilon K \rho E_s (1 - P_e(\epsilon K \rho E_s, \gamma)) \right. \\ & + 2\epsilon K \rho E_s (P_e(\epsilon K \rho E_s, \gamma))(1 - P_e(\epsilon K \rho E_s, \gamma)) + \dots \\ & + k\epsilon K \rho E_s (P_e(\epsilon K \rho E_s, \gamma))^{k-1} (1 - P_e(\epsilon K \rho E_s, \gamma)) + \dots \\ & \left. + K^2 \epsilon \rho E_s (P_e(\epsilon K \rho E_s, \gamma))^{K-1} \right\} \quad (5) \end{aligned}$$

- Energy unconstrained regime ($\bar{E}_{\text{tx}} \leq \bar{E}_f$):

$$\mathbb{E}_{\gamma} \left\{ \sum_{k=0}^{K-1} \epsilon K \rho E_s (P_e(\epsilon K \rho E_s, \gamma))^k \right\} \leq K \rho E_s \quad (6)$$

Remark 1:

Given that we make all K attempts of transmission in a frame with non-zero transmit energy, for an optimum value of $\epsilon \triangleq \epsilon^*$ resulting in the minimum outage probability, (6) holds by equality:

$$\mathbb{E}_\gamma \left\{ \sum_{k=0}^{K-1} \epsilon^* K \rho E_s \left(P_e(\epsilon^* K \rho E_s, \gamma) \right)^k \right\} = K \rho E_s \quad (7)$$

- If $\bar{E}_{\text{tx}} < \bar{E}_f$:
 - Not using enough energy
- If $\bar{E}_{\text{tx}} > \bar{E}_f$:
 - Unable to make K transmissions

Fast Fading

- Outage probability:

$$P_{\text{out}} = \left(\mathbb{E}_{\gamma} \{ P_e(\epsilon K \rho E_s, \gamma) \} \right)^K \quad (8)$$

- Average energy used per frame:

$$\bar{E}_{\text{tx}} = \sum_{k=0}^{K-1} \epsilon K \rho E_s \left(\mathbb{E}_{\gamma} \{ P_e(\epsilon K \rho E_s, \gamma) \} \right)^k \quad (9)$$

- Energy unconstrained regime ($\bar{E}_{\text{tx}} \leq \bar{E}_f$):

$$\sum_{k=0}^{K-1} \epsilon K \rho E_s \left(\mathbb{E}_{\gamma} \{ P_e(\epsilon K \rho E_s, \gamma) \} \right)^k \leq K \rho E_s \quad (10)$$

- Remark 1 follows

Harvesting Rate Optimized Variable Transmit Energy Scheme

- Choose different ϵ for different k
 - Transmit with $E_k = \epsilon_k K \rho E_s$ in the k^{th} slot
- ϵ_k for a given k remains constant for all the frames
- Find the optimum transmit energy vector: $[\epsilon_1 K \rho E_s, \dots, \epsilon_K K \rho E_s]$ that minimizes the outage

Slow Fading

- Outage probability:

$$P_{\text{out}} = \mathbb{E}_{\gamma} \left\{ \prod_{i=1}^K P_e(\epsilon_i K \rho E_s, \gamma) \right\} \quad (11)$$

- Average energy used per frame:

$$\bar{E}_{\text{tx}} = \sum_{k=1}^K \epsilon_k K \rho E_s \left\{ \mathbb{E}_{\gamma} \left\{ \prod_{i=1}^{k-1} P_e(\epsilon_i K \rho E_s, \gamma) \right\} \right\} \quad (12)$$

- Constraint:

$$\sum_{k=1}^K \epsilon_k K \rho E_s \left\{ \mathbb{E}_{\gamma} \left\{ \prod_{i=1}^{k-1} P_e(\epsilon_i K \rho E_s, \gamma) \right\} \right\} \leq K \rho E_s \quad (13)$$

Remark 2:

The optimal transmit energy vector to minimize the outage probability given by (11) satisfies (13) by equality:

$$\sum_{k=1}^K \epsilon_k K \rho E_s \left\{ \mathbb{E}_\gamma \left\{ \prod_{i=1}^{k-1} P_e(\epsilon_i K \rho E_s, \gamma) \right\} \right\} = K \rho E_s \quad (14)$$

- Optimization problem:
 - minimize:

$$P_{\text{out}} = \mathbb{E}_\gamma \left\{ \prod_{i=1}^K P_e(\epsilon_i K \rho E_s, \gamma) \right\}$$

- subject to:

$$\sum_{k=1}^K \epsilon_k K \rho E_s \left\{ \mathbb{E}_\gamma \left\{ \prod_{i=1}^{k-1} P_e(\epsilon_i K \rho E_s, \gamma) \right\} \right\} = K \rho E_s$$

- Need a closed form expression for $P_e(E_i, \gamma)$ to evaluate the optimization problem
- Even if $P_e(E_i, \gamma)$ is approximated to a strict convex function (2), Hessian of the constraint \bar{E}_{tx} is **not necessarily positive semidefinite**
 - Can not use standard convex optimization tools (e.g., Lagrange multiplier method)
- Following algorithm can be used to obtain the transmit energy values:

Algorithm 1 To find ϵ_k^* values for a slow fading channel

Define: $\epsilon_0^* = 0$; Therefore, $P_e(\epsilon_0 K \rho E_s, \gamma) = 1$;

for $k = 1$ to K **do**

▷ Solve for ϵ_k^* using the values of $\epsilon_0^*, \dots, \epsilon_{k-1}^*$ in:

$$\epsilon_k K \rho E_s \left\{ \mathbb{E}_\gamma \left\{ \prod_{i=0}^{k-1} P_e(\epsilon_i^* K \rho E_s, \gamma) \right\} \right\} = \rho E_s \quad (15)$$

end for

- Based on the fact that on an average, ρE_s arrives every slot
- The solution in closed form:

$$E_k^* = \epsilon_k^* K \rho E_s = \frac{\rho E_s}{\mathbb{E}_\gamma \left\{ \prod_{n=0}^{k-1} f_{Pe}(n, \gamma) \right\}} \quad (16)$$

- Where:

$$f_{Pe}(n, \gamma) = P_e \left(\frac{\rho E_s}{\mathbb{E}_\gamma \left\{ \prod_{i=1}^{n-1} f_{Pe}(i, \gamma) \right\}}, \gamma \right) \quad \forall n \geq 1 \quad (17)$$

with $f_{Pe}(0, \gamma) = 1$;

Proposition

For a slow fading channel and the approximate probability of error function $P_e(E_i, \gamma) \approx \exp\left(-\frac{E_i \gamma}{N_0}\right)$, the optimal transmit energy level for each (re)transmission attempt is obtained using (16).

- For the approximation, the optimization problem to be solved is:

$$\text{constraint: } \sum_{k=1}^K \frac{E_k}{\left(1 + \sum_{i=1}^{k-1} E_i/N_0\right)} = \bar{E}_f \quad (18)$$

$$\text{minimize: } P_{\text{out}} = \frac{1}{1 + \sum_{k=1}^K E_k/N_0} \quad (19)$$

$$\text{or maximize: } E_{\text{sum}} = \sum_{k=1}^K E_k \quad (20)$$

- To prove that the optimal solution to the problem is given by:

$$E_k^* = \frac{\bar{E}_f}{K} \left(1 + \sum_{i=1}^{k-1} E_i^* / N_0 \right) \quad (21)$$

- The problem is a K dim. constrained optimization problem
- We can write E_K in terms of $\bar{E}_f, E_1, \dots, E_{K-1}$
 - Reduces to $K - 1$ dim. unconstrained optimization problem
- For optimality, it is sufficient to prove that (21) holds with:

$$\frac{\partial E_{\text{sum}}}{\partial E_k^*} = 0 \quad \forall \quad 1 \leq k \leq K - 1 \quad (22)$$

- Proof is by mathematical induction; consider $K = 2$:

$$\begin{aligned}\bar{E}_f &= E_1 + \frac{E_2}{1 + E_1/N_0} \\ E_2 &= (\bar{E}_f - E_1)(1 + E_1/N_0)\end{aligned}\quad (23)$$

- Substituting in (20) for $K = 2$:

$$E_{\text{sum}} = E_1 + (\bar{E}_f - E_1)(1 + E_1/N_0)\quad (24)$$

- Differentiating E_{sum} w.r.t E_1 :

$$\frac{\partial E_{\text{sum}}}{\partial E_1} = 1 - (1 + E_1/N_0) + (\bar{E}_f - E_1)/N_0 = 0\quad (25)$$

$$\Rightarrow E_1 = \frac{\bar{E}_f}{2}\quad (26)$$

$$\Rightarrow E_2 = \frac{\bar{E}_f}{2}(1 + E_1/N_0)\quad (27)$$

- (21) holds for $K = 2$

- Assume that (21) holds for $K = K'$; the problem for $K = K' + 1$ is:

$$\max_{E_1, \dots, E_{K'+1}} \sum_{k=1}^{K'+1} E_k \quad (28)$$

$$\text{subject to: } \sum_{k=1}^{K'+1} \frac{E_k}{\left(1 + \sum_{i=1}^{k-1} E_i/N_0\right)} = \bar{E}_f \quad (29)$$

- Fix E_1 : $0 \leq E_1 \leq \bar{E}_f$; above problem reduces to:

$$\max_{E_2, \dots, E_{K'+1}} \sum_{k=2}^{K'+1} E_k \quad (30)$$

$$\text{subject to: } \sum_{k=2}^{K'+1} \frac{E_k}{\left(1 + \sum_{i=2}^{k-1} E_i/N'_0\right)} = (\bar{E}_f - E_1)(1 + E_1/N_0) \quad (31)$$

where: $N'_0 = N_0(1 + E_1/N_0)$

- The problem is a $K = K'$ case, with $k = 2, \dots, K' + 1$
 - Solution (21) holds for $K = K'$:

$$E_k^* = \frac{(\bar{E}_f - E_1)(1 + E_1/N_0)}{K'} \left(1 + \sum_{i=2}^{k-1} E_i^*/N_0' \right) \forall 2 \leq k \leq K' + 1 \quad (32)$$

- Reduces to:

$$E_k^* = \frac{(\bar{E}_f - E_1)(1 + E_1/N_0)}{K'} \left(1 + \frac{\bar{E}_f - E_1}{K'N_0} \right)^{k-2} \quad (33)$$

- Only variable left to be maximized in (28) is E_1 :

$$\max_{E_1, \dots, E_{K'+1}} \sum_{k=1}^{K'+1} E_k = \max_{0 \leq E_1 \leq \bar{E}_f} E_1 + \sum_{k=2}^{K'+1} E_k^* \quad (34)$$

$$= \max_{0 \leq E_1 \leq \bar{E}_f} E_1 + \frac{(\bar{E}_f - E_1)(1 + E_1/N_0)}{K'} \times \sum_{k=2}^{K'+1} \left(1 + \frac{\bar{E}_f - E_1}{K'N_0}\right)^{k-2} \quad (35)$$

$$= \max_{0 \leq E_1 \leq \bar{E}_f} E_1 + \frac{(\bar{E}_f - E_1)(1 + E_1/N_0)}{K'} \times \sum_{k=0}^{K'-1} \left(1 + \frac{\bar{E}_f - E_1}{K'N_0}\right)^k \quad (36)$$

$$= \max_{0 \leq E_1 \leq \bar{E}_f} E_1 - N_0(1 + E_1/N_0) \times \left(1 - \left(1 + \frac{\bar{E}_f - E_1}{K'N_0}\right)^{K'}\right) \quad (37)$$

Differentiating w.r.t E_1 :

$$\frac{\partial E_{\text{sum}}}{\partial E_1} = \left(1 + \frac{\bar{E}_f - E_1}{K'N_0}\right)^{K'} - (1 + E_1/N_0) \left(1 + \frac{\bar{E}_f - E_1}{K'N_0}\right)^{K'-1} = 0 \quad (38)$$

$$E_1^* = \frac{\bar{E}_f}{K' + 1} \quad (39)$$

Remarks

Remark 3:

For an infinite battery capacity, the optimal transmit energy vector depends **only on the average energy arrival per slot** ρE_s and not on the individual values of ρ and E_s .

Remark 4:

For an optimal policy, the energy with which a (re)transmission attempt is made, is simply **the expected energy that will be available** in the battery at the start of that attempt.

Remark 5:

In any given frame, the optimal transmit energy level at a higher retransmission attempt is higher; i.e., the optimal transmit energy policy is a **strictly increasing policy** ($E_1^* < E_2^* \dots < E_K^*$), with non-zero transmit energy levels.

Fast Fading

- Outage probability:

$$P_{\text{out}} = \prod_{i=1}^K \mathbb{E}_{\gamma} \left\{ P_e(\epsilon_i K \rho E_s, \gamma) \right\} \quad (40)$$

- Average energy used per frame:

$$\bar{E}_{\text{tx}} = \sum_{k=1}^K \epsilon_k K \rho E_s \prod_{i=1}^{k-1} \mathbb{E}_{\gamma} \left\{ P_e(\epsilon_i K \rho E_s, \gamma) \right\} \quad (41)$$

- Introducing the constraint ($\bar{E}_{\text{tx}} \leq \bar{E}_f$):

$$\sum_{k=1}^K \epsilon_k K \rho E_s \prod_{i=1}^{k-1} \mathbb{E}_{\gamma} \left\{ P_e(\epsilon_i K \rho E_s, \gamma) \right\} \leq K \rho E_s \quad (42)$$

- Remark 2 follows; (42) satisfied by equality:

$$\sum_{k=1}^K \epsilon_k K \rho E_s \prod_{i=1}^{k-1} \mathbb{E}_\gamma \left\{ P_e(\epsilon_i K \rho E_s, \gamma) \right\} = K \rho E_s \quad (43)$$

- Optimization problem:

- minimize:

$$P_{\text{out}} = \prod_{i=1}^K \mathbb{E}_\gamma \left\{ P_e(\epsilon_i K \rho E_s, \gamma) \right\}$$

- subject to:

$$\sum_{k=1}^K \epsilon_k K \rho E_s \prod_{i=1}^{k-1} \mathbb{E}_\gamma \left\{ P_e(\epsilon_i K \rho E_s, \gamma) \right\} = K \rho E_s$$

Dual Problem [H. Seo and B. G. Lee, 2007]

- Define $P_{e_{\text{avg}}}(E_i) = \mathbb{E}_{\gamma} \{P_e(E_i, \gamma)\}$
- Optimization problem addressed:

- minimize: $\bar{E}_{\text{tx}} = \sum_{k=1}^K E_k \prod_{i=1}^{k-1} P_{e_{\text{avg}}}(E_i)$

- subject to: $\text{FDR} = \prod_{i=1}^K P_{e_{\text{avg}}}(E_i) = c$

- Solution:

$$\frac{P_{e_{\text{avg}}}(E_{i+1}^*)}{P'_{e_{\text{avg}}}(E_{i+1}^*)} - E_{i+1}^* - \frac{1}{P'_{e_{\text{avg}}}(E_i^*)} = 0 \quad (44)$$

Remark 5:

The solution to the problem of minimizing the average energy used per frame subject to a given outage probability constraint is equivalent to that of the dual problem of minimizing the outage probability subject to a given average energy constraint.

- Policy requires $P_{e_{\text{avg}}}(E_i)$ to be analytic and strictly convex
- Alternative suboptimal policy:

Algorithm 2 To find ϵ_k^* values for a fast fading channel

Define: $\epsilon_0^* = 0$; Therefore, $P_e(\epsilon_0 K \rho E_s, \gamma) = 1$;

for $k = 1$ to K **do**

▷ Solve for ϵ_k^* using the values of $\epsilon_0^*, \dots, \epsilon_{k-1}^*$ in:

$$\epsilon_k K \rho E_s \prod_{i=1}^{k-1} \mathbb{E}_{\gamma} \left\{ P_e(\epsilon_i^* K \rho E_s, \gamma) \right\} = \rho E_s \quad (45)$$

end for

- Does not require $P_{e_{\text{avg}}}(E_i)$ to be analytic or strictly convex
- Provides insight into the optimal policy

- The solution in closed form is given as:

$$E_k^* = \epsilon_k^* K \rho E_s = \frac{\rho E_s}{\prod_{n=0}^{k-1} \mathbb{E}_\gamma \{ f_{Pe}(n, \gamma) \}} \quad (46)$$

- Where:

$$f_{Pe}(n, \gamma) = P_e \left(\frac{\rho E_s}{\prod_{i=1}^{n-1} \mathbb{E}_\gamma \{ f_{Pe}(i, \gamma) \}}, \gamma \right) \quad \forall n \geq 1 \quad (47)$$

with $f_{Pe}(0, \gamma) = 1$;

Finite Battery Capacity

- In the case of infinite energy buffer:
 - No energy lost because of the battery being full
 - Randomness in harvesting is completely absorbed
- Infinite battery capacity assumption for a battery with finite energy buffer is valid when:

$$\left(\sum_{k=1}^K \epsilon_k K \rho E_s \right) \ll B_{\max} \quad (48)$$

- What if the battery is not that large?
 - But large enough to allow all K transmission attempts:

$$\left(\sum_{k=1}^K \epsilon_k K \rho E_s \right) < B_{\max} \quad (49)$$

- Number of instances when the battery hits its upper limit
 - Energy arriving at such instances can not be harvested
 - Average energy harvested:

$$\bar{E}'_f = \bar{E}_f - \delta = K \rho E_s - \delta \quad (50)$$

Exact Modification

- Optimal transmit energy vector should satisfy:

$$\bar{E}_{\text{tx}} = \bar{E}'_f = K\rho E_s - \delta \quad (51)$$

- Where $\delta = \Pr(B_i = B_{\max}) \times \rho \times E_s$
 - $\Pr(B_i = B_{\max})$: steady state probability of the battery being full.¹
 - $\Pr(B_i = B_{\max})$ and \bar{E}'_{tx} are interrelated
- Transmit energy vector for a finite battery buffer can be obtained as:

Algorithm 3 To find ϵ_k^* values for a finite battery capacity

Initialize: $\Pr(B_i = B_{\max} | \bar{E}'_f) = 0$

do

▷ $\bar{E}'_f = K\rho E_s - \Pr(B_i = B_{\max} | \bar{E}'_f) \times \rho E_s$

▷ Evaluate $\epsilon_1^*, \dots, \epsilon_K^*$ for the given \bar{E}'_f

▷ Evaluate corresponding $\Pr(B_i = B_{\max} | \bar{E}'_f)$

while ($\bar{E}'_{\text{tx}} \neq K\rho E_s - \Pr(B_i = B_{\max} | \bar{E}'_f) \times \rho E_s$)

¹2012, A. Aprem and C. R. Murthy

Heuristic Modification

- Exact modification: does not take into account the sporadicity of harvesting
- Following heuristic modification can be used:
 - Transmit energy at the k^{th} slot and i^{th} instant:

$$E_{\text{tx}} = \min(B_i, \epsilon_k K \rho E_s) \quad (52)$$

- Heuristically takes into account the loss of energy
- Does not allow the battery to reach its full state too often

Simulations and Results

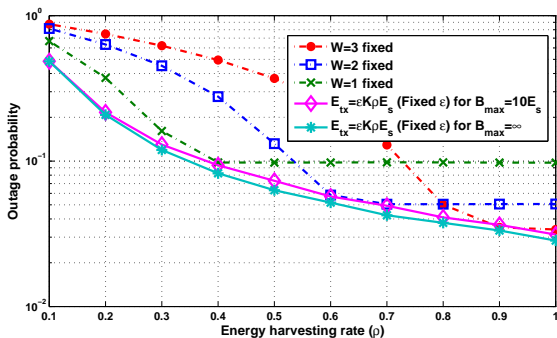


Figure : Comparison of the performance of the harvesting rate optimized fixed energy transmission scheme with the standard fixed energy transmission schemes for $B_{\max} = 10E_s$ for a slow fading channel. An additional plot of outage probability obtained using the same policy but for an infinite energy buffer is also plotted.

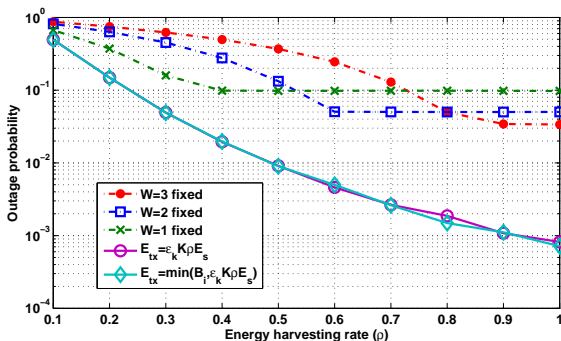


Figure : Performance evaluation of the harvesting rate optimized variable transmit energy policy for a slow fading channel and a transmitter with infinite energy buffer. Here, $E_s = 12\text{dB}$, $D = 50$ and $K = 4$.

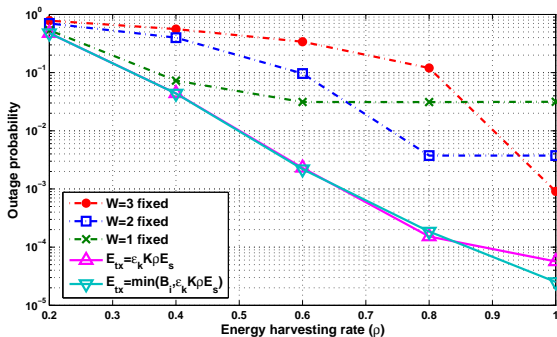


Figure : Performance evaluation of the harvesting rate optimized variable transmit energy policy for a fast fading channel and a transmitter with infinite energy buffer. Here, $E_s = 5\text{dB}$, $D = 20$ and $K = 4$.

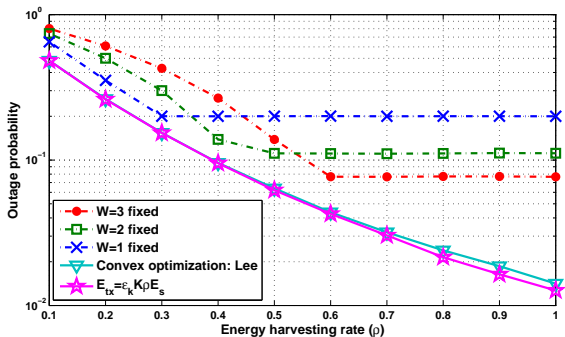


Figure : Comparison of $E_{tx} = \epsilon_k K \rho E_s$ policy with policy obtained in [H. Seo and B. G. Lee, 2007] for a slow fading channel. Here, $E_s = 1.2\text{dB}$ and $K = 4$.

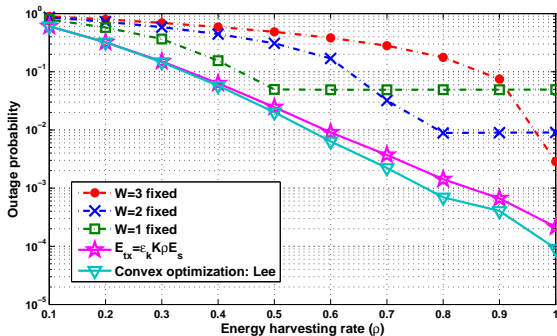


Figure : Comparison of $E_{tx} = \epsilon_k K \rho E_s$ policy with policy obtained in [H. Seo and B. G. Lee, 2007] for a fast fading channel. Here, $E_s = 0.5\text{dB}$ and $K = 4$.

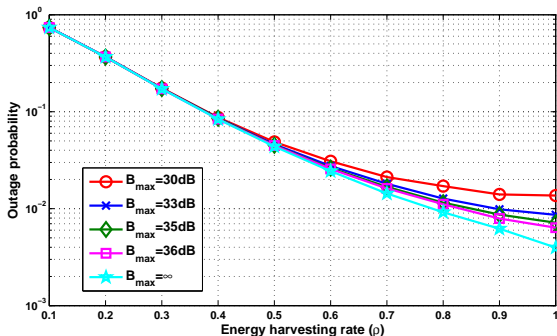


Figure : Comparison of the performance of the harvesting rate optimized variable transmission energy scheme for different finite battery capacities when the channel is slow fading. Here, $E_s = 10\text{dB}$, $D = 50$, $K = 4$ and $\sum_{k=1}^K \epsilon_k K \rho E_s \approx 30\text{dB}$.

Here, $E_s = 10\text{dB}$, $D = 50$, $K = 4$ and $\sum_{k=1}^K \epsilon_k K \rho E_s \approx 30\text{dB}$.

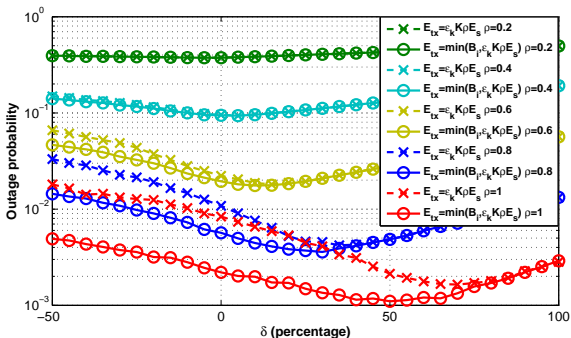


Figure : $\bar{E}'_f = K\rho E_s - \delta \times \rho E_s$ for $B_{\max} = 80E_s$ and $E_s = 0$ dB; δ varying from -50% to +100%. Converged δ values obtained using Algorithm 3: $\rho = 0.2 : 0.2\%$, $\rho = 0.4 : 0.71\%$, $\rho = 0.6 : 2.02\%$, $\rho = 0.8 : 5.75\%$, $\rho = 1 : 13.95\%$.

Thank you!

