

## Lecture 18 — October 6

Lecturer: Aditya Gopalan

Scribe: Mohammadi Zaki

## 18.1 TWO MIRROR-DESCENT ALGORITHMS FOR ON-LINE CONVEX OPTIMIZATION (O.C.O.)

Let  $K$  be a convex set and let  $R$  be a Legendre function over  $K$ .

### 18.1.1 “Lazy” Mirror Descent (“Follow the Regularized Leader FTRL”)

#### Algorithm:

**Initialize:**  $\tilde{w}_1 \in \mathbb{R}^d$ ,  $w_1 = \Pi_{R,K}(\tilde{w}_1)$

**Repeat:**  $\forall t = 2, 3, 4, \dots$ ,

$$\begin{aligned}\nabla R(\tilde{w}_t) &= \nabla R(\tilde{w}_{t-1}) - z_{t-1} \\ \Rightarrow \tilde{w}_t &= \nabla R^*(\nabla R(\tilde{w}_{t-1}) - z_{t-1}) \\ w_t &= \Pi_{R,K}(\tilde{w}_t)\end{aligned}$$

See Fig. 18.1 for a schematic description.

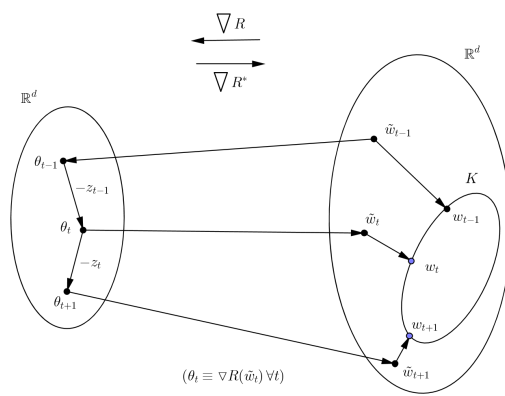


Figure 18.1. “LAZY” OMD

## 18.1.2 “Active” Mirror Descent or Proximal Point Algorithm

### Algorithm:

**Initialize:**  $w_1 \in K$

**Repeat:**  $\forall t = 2, 3, 4, \dots,$

$$\nabla R(y_t) = \nabla R(w_{t-1}) - z_{t-1}$$

$$w_t = \Pi_{R,K}(y_t).$$

See Fig. 18.2 for a schematic description.

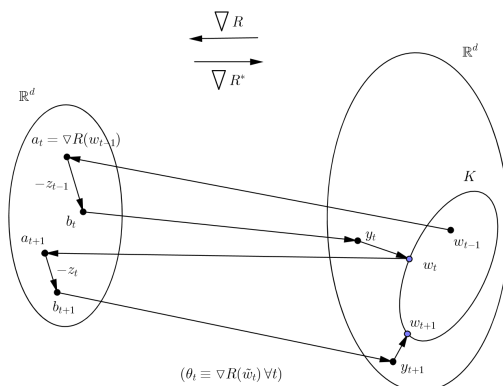


Figure 18.2. “ACTIVE” OMD

**Theorem 18.1 (Regret Bound for Active OMD).** For the active OMD algorithm with  $R$  a Legendre function,  $\forall u \in K$ ,

$$\sum_{t=1}^T [f_t(w_t) - f_t(u)] \leq D_R(u, w_1) - D_R(u, w_{T+1}) + \sum_{t=1}^T D_R(w_t, y_{t+1}).$$

[**Aside:** For the LAZY version of OMD,

$$\sum_{t=1}^T [f_t(w_t) - f_t(u)] \leq D_R(u, w_1) - D_R(u, \tilde{w}_{T+1}) + \sum_{t=1}^T D_R(w_t, \tilde{w}_{t+1}).]$$

**Proof:** Let us look at a single term of L.H.S,

$$\begin{aligned} f_t(w_t) - f_t(u) &\leq \langle \nabla f_t(w_t), w_t - u \rangle && \text{because } f \text{ is convex.} \\ &= \langle \nabla R(w_t) - \nabla R(y_{t+1}), w_t - u \rangle \\ &= \langle w_t - u, \nabla R(w_t) - \nabla R(y_{t+1}) \rangle \\ &= D_R(u, w_t) - D_R(u, y_{t+1}) + D_R(w_t, y_{t+1}) && \text{By 3-point inequality} \\ &\leq D_R(u, w_t) - D_R(u, w_{t+1}) + D_R(w_t, y_{t+1}) \\ & \text{[\"Generalized Pythagoras\"} \\ & \forall u \in K, \forall w : \\ & D_R(u, w) \leq D_R(u, \Pi_{R,K}(w)) + D_R(\Pi_{R,K}(w), w) \end{aligned}$$

Adding over  $t = 1, 2, \dots, T$ ,

$$f_t(w_t) - f_t(u) \leq D_R(u, w_1) - D_R(u, w_{T+1}) + \sum_{t=1}^T D_R(w_t, y_{t+1}).$$

□

**Note:** This can be specialized to give concrete regret bounds, e.g., [1, Zinkevich, '03]. Consider active OMD with  $R(x) = \frac{1}{2\eta} \|x\|_2^2$ .

$$y_t = w_{t-1} - \eta \nabla f_{t-1}(w_{t-1})$$

$$w_t = \Pi_{R,K}(y_t) \quad [\equiv \text{POGD}.]$$

$$\text{Regret}_T(\text{POGD}) \leq \frac{1}{2\eta} D^2 + \frac{\eta}{2} T G^2.$$

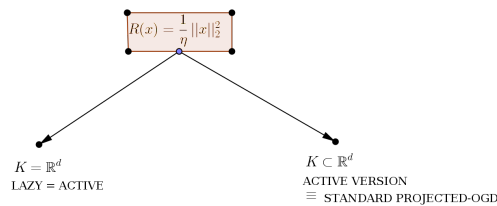
where,

$$D = \max_{x,y \in K} \|x - y\|_2^2 \quad (\text{“diameter”}).$$

$$G = \sup_{t \leq T, x \in K} \|\nabla f_t(x)\|_2^2. \quad (\text{“Largest possible gradient”})$$

### 18.1.3 Examples:

1.  $R(x) = \frac{1}{2\eta} \|x\|_2^2$ .



## 18.2 Online Learning with limited/partial information-BANDITS

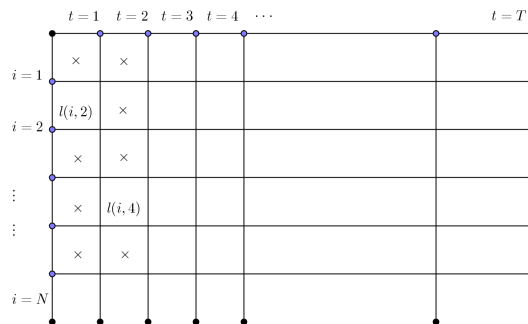
Recall: The problem of competing with the set of  $N$  actions ( $\equiv$  experts problem or O.C.O. with linear losses).

At each time  $t = 1, 2, \dots$ ,

1. Environment sets a loss  $l(i, t), \forall i = 1, 2, \dots, N$ .
2. Learner/algorithm plays action  $I_t \in [N]$ .
3. Algorithm suffers loss  $l(I_t, t)$ .
4. Algorithm gets to “see” only  $l(I_t)$  and NOT  $(l(i, t))_{i=1}^N$ .

### 18.2.1 Assumptions on $\{l(i, t)\}_{i,t}$

We will assume “Oblivious adversary”, i.e.,  $l(i, t)$  is deterministic but arbitrarily chosen at  $t = 0$ . ( $l(i, t)$  chosen is “independent” of the algorithm chosen).



#### Notes:

1.  $Regret(T) = \sum_{t=1}^T l(I_t, t) - \min_{i \in [N]} \sum_{t=1}^T l(i, t)$ .

2. Randomization is absolutely necessary here to avoid linear regret as for any deterministic algorithm the adversary can set worst loss for each round forcing the algorithm to incur regret of  $\mathcal{O}(T)$ .

# Bibliography

- [1] Martin Zinkevich, "Online convex programming and generalized infinitesimal gradient ascent," in *Proc. of the Twenty International Conference on Machine Learning (ICML'03)*, pp. 928-936, AAAI Press, 2003.