Optimal Transmit Power Control Policies for Device-to-device Networks with Energy Harvesting

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Outline

- Network assisted device discovery
- Throughput optimal policies for dual EH links

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Device Discovery: Motivation¹



- Achieved utility depends on number of discovered devices
- In EH networks exploration vs exploitation trade-off does not exist

¹Zou et al., Proximity Discovery for Device-to-Device Communications over a Cellular Network, IEEE Comm. Mag., June 2014

System Model

- The devices are assumed to be distributed according to a homogeneous Poisson point process (PPP) with density \u03b1
- Slotted ALOHA protocol with parameter p
- Discovery phase followed by a communication phase
- Network infrastructure assists in perfect synchronization of the slot boundaries across the nodes
- Also notifies the optimal transmission probability parameter p at the beginning of the discovery phase
- Spatial distribution of the nodes changes in an i.i.d. fashion at the beginning of each discovery phase
- Device choose a resource block uniformly randomly from M FDM RBs

System Model: Contd

- In each slot, a node harvests energy E_s with probability ρ
- A node participates in the discovery process if it has E energy in the battery
- ► Energy used for transmission and reception are E = PT_p and R, respectively (R ≤ E)
- Operation of an EHN is constrained by energy neutrality constraint
- Battery evolves as

$$B_{n+1} = \max\{(B_n + \mathbb{1}_{\{e_n \neq 0\}} - \mathbb{1}_{\{T \neq 0\}} PT_{\rho} + \mathbb{1}_{\{Rx \neq 0\}} R)^+, B_{\max}\}$$

Transition probabilities

$$B_{n+1} = \begin{cases} \min\{B_n + 1 - E, B_{\max}\} \\ B_n - E \\ \min\{B_n + 1 - R, B_{\max}\} \\ B_n - R \\ B_n + 1 \\ B_n \\ \min\{B_n + 1 - R, B_{\max}\} \\ B_n - R \\ \min\{B_n + 1, B_{\max}\} \\ B_n \\ B_n \\ B_n \\ B_n + 1 \end{cases}$$

w.p.
$$p\rho$$
 if $B_n \ge E$
w.p. $p(1-\rho)$ if $B_n \ge E$
w.p. $q\rho$ if $B_n \ge E$
w.p. $q(1-\rho)$ if $B_n \ge E$
w.p. $(1-p-q)(\rho)$ if $B_n \ge E$
w.p. $(1-p-q)(1-\rho)$ if $B_n \ge E$
w.p. $r\rho$ if $E > B_n \ge R$
w.p. $r(1-\rho)$ if $E > B_n \ge R$
w.p. $(1-r)\rho$ if $E > B_n \ge R$
w.p. $(1-r)\rho$ if $E > B_n \ge R$
w.p. $(1-\rho)$ if $B_n < R$.
w.p. $(1-\rho)$ if $B_n < R$.
w.p. ρ if $B_n < R$.

- Given: λ, ρ and R
- Design p and P to maximize the number of discovered devices

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Metric: mean number of discovered devices

Mean number of discovered devices, $\mathbb{E}(D)$

Lemma

For an energy harvesting device-to-device network with Rayleigh fading links and path-loss exponent $\alpha = 4$, $\mathbb{E}(D)$, is given as

$$\mathbb{E}(D) = \frac{\lambda_a \pi^{\frac{3}{2}} p(1-p)}{2\sqrt{\gamma_0 \sigma^2}} \exp\left(\left(\frac{\lambda_a \pi^2 p}{4M\sigma}\right)^2\right) \operatorname{erfc}\left(\frac{\lambda_a \pi^2 p}{4M\sigma}\right),$$

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where $\operatorname{erfc}(t) = \frac{2}{\sqrt{\pi}} \int_t^\infty \exp(-y^2) dy$, and $\lambda_a = \lambda \pi_a$. π_a : the stationary probability that the node is active.

Optimization Problem

 $\begin{array}{ll} \max_{p,P} & \mathbb{E}(D) \\ \text{subject to} & P \leq P_{\max} \end{array}$

and energy neutrality constraint

For a node operating under average power constraint

$$\pi_{a} = 1 - \Theta(e^{r * B_{\max}})$$

Reformulated problem:

$$\begin{array}{ll} \max_{p,P} & \mathbb{E}(D) \\ \text{subject to} & P \leq P_{\max}, \\ & pE + (1-p)R \leq \rho E_s. \end{array}$$

Proof

We work out

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Throughput optimal policies for dual EH links

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System Model



ACK/NACK signal provides the perfect synchronization

Goal

Goal: Maximize the long-term time averaged utility

$$\max_{\boldsymbol{e}_n^t, \boldsymbol{e}_n^r} \quad \mathcal{U} = \max_{\boldsymbol{e}_n^t, \boldsymbol{e}_n^r} \left(\liminf_{T \to \infty} \sum_{n=1}^T \mathbb{1}_{\{\boldsymbol{e}_n^r \neq 0\}} U(\boldsymbol{e}_n^t) \right)$$

where U(.) is a concave non-decreasing function.

Lemma

For an uncoordinated dual EH link the time-averaged utility is upper bounded as

$$\mathcal{U} \leq \min\left(\frac{U(\mu_t)}{U(\mathcal{B}_{\max}^t)}, \mu_r\right) U(\mathcal{B}_{\max}^t).$$

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Proof

Using ergodicity, the time-averaged utility is written as

$$\begin{aligned} \mathcal{U} &= \liminf_{T \to \infty} \sum_{n=1}^{T} \mathbb{1}_{\{e_n^r \neq 0\}} U(e_n^t), \\ \stackrel{(a)}{=} & \mathbb{E} \left[\mathbb{1}_{\{e_n^r \neq 0\}} U(e_n^t) \right], \\ \stackrel{(b)}{=} & \min \left(\frac{\mathbb{E} \left[U(e_n^t) \right]}{U(B_{\max}^t)}, \mathbb{E} \left[\mathbb{1}_{\{e_n^r \neq 0\}} \right] \right) U(B_{\max}^t), \\ \stackrel{(c)}{\leq} & \min \left(\frac{U \left[\mathbb{E}(e_n^t) \right]}{U(B_{\max}^t)}, \mathbb{E} \left[\mathbb{1}_{\{e_n^r \neq 0\}} \right] \right) U(B_{\max}^t), \\ \stackrel{(d)}{\leq} & \min \left(\frac{U(\mu_t)}{U(B_{\max}^t)}, \mu_r \right) U(B_{\max}^t). \end{aligned}$$

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Optimal Policy

• CASE I:
$$\frac{\mu_r}{R} \ge 1$$

$$\boldsymbol{e}_{n}^{t} = \begin{cases} \mu_{t} + \delta, & \text{if } B_{n}^{t} \geq \frac{B_{\max}^{t}}{2}, \\ \min\left\{\mu_{t} - \delta, B_{n}^{t}\right\}, & \text{if } B_{n}^{t} < \frac{B_{\max}^{t}}{2}. \end{cases}$$
(1)

► CASE II(^µ_R > 1): The receiver employs a policy where it turns on after every N_r slot which is given as follows

$$N_r = \begin{cases} [N], & \text{if } B_n^r \leq \frac{B_{\max}^r}{2}, \\ \lfloor N \rfloor, & \text{if } B_n^r < \frac{B_{\max}^r}{2}. \end{cases}$$

In each slot transmitter allocates the energy according to (1), and transmits the accumulated energy in the next slot when the receiver is on.

Performance: CASE I



Figure: Comparison of upper bound with policy in (1). The parameters chosen are $B_{max}^t = B_{max}^r = 50$

Performance: CASE II



Figure: Case II: Comparison of upper bound and policy in (15) and (15). The parameters chosen are $B_{max}^t = B_{max}^r = 50$