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# Beamforming Design for Full-Duplex Two-Way Amplify-and-Forward MIMO Relay

(Yeonggyu Shim, Wan Choi, Hyuncheol Park)

- Contributions
  - Optimized Relay Beamforming matrix
  - Optimized Receive Beamforming matrices at sources
- System Model
  - Consider a two-way FD AF relaying system consisting of two source nodes s<sub>1</sub> and s<sub>2</sub> with N<sub>s</sub> antennas, and one relay node r with N<sub>r</sub> antennas, where all nodes operate in the FD mode.
  - Channel links are modelled as Frequency-flat fading channels and assumed to be static in Time slot.
  - CSI between two nodes is assumed to be perfectly known.
  - CSI of the loopback channels is assumed to be imperfect.
  - focuses on loopback Self Interference.





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Received signals with imperfect loopback SI cancellation:

• At relay in time slot 
$$t - 1(t \ge 2)$$
 is given by,

$$\hat{\mathbf{y}}_{r}^{(t-1)} = \mathbf{H}_{1,r}^{(t-1)} \mathbf{x}_{1}^{(t-1)} + \mathbf{H}_{2,r}^{(t-1)} \mathbf{x}_{2}^{(t-1)} + \mathbf{n}_{r}^{(t-1)} \\
+ \sum_{i=0}^{t-2} \left\{ \prod_{j=1}^{t-1-i} (\mathbf{\Delta}_{r,r}^{(t-j)} \mathbf{F}^{(t-j)}) \\
\times (\mathbf{H}_{1,r}^{(i)} \mathbf{x}_{1}^{(i)} + \mathbf{H}_{2,r}^{(i)} \mathbf{x}_{2}^{(i)} + \mathbf{n}_{r}^{(i)}) \right\}$$
(1)

At source / is,

$$\hat{\mathbf{y}}_{l}^{(t)} = \mathbf{H}_{r,l}^{(t)} \mathbf{F}^{(t)} \mathbf{H}_{\bar{l},r}^{(t-1)} \mathbf{x}_{\bar{l}}^{(t-1)} + \mathbf{H}_{r,l}^{(t)} \mathbf{F}^{(t)} \\
\sum_{i=0}^{t-2} \left\{ \prod_{j=1}^{t-1-i} (\mathbf{\Delta}_{r,r}^{(t-j)} \mathbf{F}^{(t-j)}) (\mathbf{H}_{l,r}^{(i)} \mathbf{x}_{l}^{(i)} + \mathbf{H}_{\bar{l},r}^{(i)} \mathbf{x}_{\bar{l}}^{(i)} + \mathbf{n}_{r}^{(i)}) \right\} \\
+ \mathbf{H}_{r,l}^{(t)} \mathbf{F}^{(t)} \mathbf{n}_{r}^{(t-1)} + \mathbf{\Delta}_{l,l}^{(t)} \mathbf{x}_{l}^{(t)} + \mathbf{n}_{l}^{(t)}.$$
(2)

 $\mathbf{F}^{(t)}$ ,  $\mathbf{R}^{(t)}_{l}$  are the relay Beamforming matrix and receive Beamforming matrix respectively.

### MMSE Based Beamforming design

Proposes an iterative algorithm which decouples the primal problem into two subproblems and solve them alternately; one is for relay beamforming design and the other is for receive beamforming design at sources.

## Results:

 Closed form expressions for Relay Beamforming and Receive Beamforming matrices are derived.

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 An Iterative Algorithm is proposed for joint Beamforming Design.

# Resource Allocation for D2D Communication Underlaid Cellular Networks Using Graph-Based Approach (Tuong Duc Hoang,Long Bao Le,Tho Le-Ngoc)

#### Contributions

focuses on the radio resource allocation for D2D communications in cellular networks for the first scenario and the developed algorithm for scenario I is employed to tackle the resource allocation for scenario II

- Novel Iterative Rounding algorithm is proposed solve the subband assignment problem based on the combination of linear programming and efficient rounding techniques. Specifically, each iteration solves a relaxed version of the subband assignment problem for unallocated subbands and network links.
- Scenario I: Each active (admitted) D2D link is assigned one subband and each subband is exploited by at most one D2D link.
- Scenario II:Each active (admitted) D2D link can be assigned multiple subbands and each subband is exploited by at most one D2D link.

 Scenario III:Each active (admitted) D2D link can be assigned multiple subbands and each subband can be exploited by multiple D2D links.

# System model:

- The spectrum sharing problem among multiple D2D and cellular links in the uplink direction.
- Let N = {1, 2, ..., N} with size |N| = N be the set of subbands in the system.
- We denote  $\mathcal{K}_c = \{1, 2, ..., \mathcal{K}_c\}$  as the set of cellular links,  $\mathcal{K}_d = \{\mathcal{K}_c + 1..., \mathcal{K}_c + \mathcal{K}_d\}$ as the set of D2D links, and  $\mathcal{K} = \mathcal{K}_c \cup \mathcal{K}_d$  as the set of all communications links with size  $|\mathcal{K}| = \mathcal{K}_c + \mathcal{K}_d = \mathcal{K}.$
- Assumption: Each subband can be allocated to at most one cellular and one D2D link.

► The signal to interference plus noise ratio (SINR) achieved by link k ∈ K on subband n can be expressed as

$$\Gamma_{k}^{[n]}(\mathbf{p}^{[n]}, \boldsymbol{\rho}^{[n]}) = \frac{\rho_{k}^{[n]} p_{k}^{[n]} h_{kk}^{[n]}}{\sigma_{k}^{[n]} + \sum_{l \in \mathcal{K} \setminus \mathcal{K}_{k}} \rho_{l}^{[n]} p_{l}^{[n]} h_{kl}^{[n]}},$$
(3)

► The achievable rates of link k ∈ K on subband n and all the subbands can be expressed as

$$r_{k}^{[n]}(\mathbf{p}^{[n]}, \boldsymbol{\rho}^{[n]}) = \log_{2} \left( 1 + \Gamma_{k}^{[n]}(\mathbf{p}^{[n]}, \boldsymbol{\rho}^{[n]}) \right),$$
$$r_{k}(\mathbf{p}, \boldsymbol{\rho}) = \sum_{n \in \mathcal{N}} r_{k}^{[n]}(\mathbf{p}^{[n]}, \boldsymbol{\rho}^{[n]}),$$

 The considered resource allocation problem can now be formulated as

$$\max_{\mathbf{p},\boldsymbol{\rho}} \ \mathcal{R} = \sum_{k \in \mathcal{K}_c} \alpha r_k \left(\mathbf{p}, \boldsymbol{\rho}\right) + \sum_{k \in \mathcal{K}_d} \left(1 - \alpha\right) r_k \left(\mathbf{p}, \boldsymbol{\rho}\right)$$

- Solution:
  - we first characterize the optimal power allocation solution for a given subband assignment.
  - Based on this result, we formulate the subband assignment problem by using the graph-based approach, in which each link corresponds to a vertex and each subband assignment is represented by a hyper-edge.

 We then propose an iterative rounding algorithm and an optimal branch-and-bound (BnB) algorithm to solve the resulting graph-based problem.

### Secure Transmission in Cooperative Relaying Networks With Multiple Antennas

(Yuzhen Huang; Jinlong Wang; Caijun Zhong; Trung Q. Duong; George K. Karagiannidis)

#### Contributions

- For ZF/MRC and ZF/SC, we present novel closed-form lower and upper bounds for the secrecy outage probability and the probability of non-zero secrecy capacity, respectively, as well as a simple high SNR secrecy outage analysis.
- For MRT/MRC and MRT/SC, closed-form approximations for the secrecy outage probability and the probability of non-zero secrecy capacity are provided, respectively.
- For the CJ/ZF scheme, new exact closed-form expressions for the secrecy outage probability and the probability of non-zero secrecy capacity are derived.

The analytical results suggest that the ZF/MRC (MRT/MRC) scheme always achieves better performance than that of the corresponding ZF/SC (MRT/SC) scheme. In addition, the ZF/MRC (ZF/SC) scheme outperforms the corresponding MRT/MRC (MRT/SC) scheme in the low SNR regime, while in the high SNR regime, the MRT/MRC (MRT/SC) scheme attains better secrecy performance than the corresponding ZF/MRC (ZF/SC) scheme.

**System model:** We consider a dual-hop multiple antenna AF relaying network, where both Alice (A), Bob (B), and Eve (E) are equipped with a single antenna, while the relay (R) is equipped with M antennas.

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SNR of  $A \rightarrow R \rightarrow B$  link is given by,

$$\gamma_{\rm ARB} = \frac{P_s}{\sigma^2} \frac{\left| \mathbf{h}_{\rm RB}^{\dagger} \mathbf{W} \mathbf{h}_{\rm AR} \right|^2}{1 + \left\| \mathbf{h}_{\rm RB}^{\dagger} \mathbf{W} \right\|_F^2}.$$
 (4)

SNR of  $A \rightarrow R \rightarrow E$  link can be derived as,

$$\gamma_{\text{ARE}} = \frac{P_s}{\sigma^2} \frac{\left| \mathbf{h}_{\text{RE}}^{\dagger} \mathbf{W} \mathbf{h}_{\text{AR}} \right|^2}{1 + \left\| \mathbf{h}_{\text{RE}}^{\dagger} \mathbf{W} \right\|_F^2}.$$
(5)

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instantaneous SNRs of the main channel under MRC,

$$\gamma_{B_{MRC}} = \gamma_{AB} + \gamma_{ARB}$$
$$\gamma_{E_{MRC}} = \gamma_{AE} + \gamma_{ARE}.$$

The achievable secrecy capacity of the relaying wiretap channels is defined as,

$$C_{\rm S} \stackrel{\Delta}{=} \frac{1}{2} \left[ \log_2 \left( 1 + \gamma_{\rm B_i} \right) - \log \left( 1 + \gamma_{\rm E_{\rm MRC}} \right) \right]^+,$$

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The SNR at B and E is maximized using different schemes.