# Cognitive Radio Simulations in NS2

Santhosh Kumar, Nirmal Shende

October 28, 2011

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Goals

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• To study of the effect of scanning duration on the Cognitive Radio network throughput in a WLAN primary through NS2 simulations.

- Analytical models for WLAN networks already exist in literature
- We extend one such model to include the effects of scanning delay, sensing errors on the throughput performance of CR
- Bianchi's discrete time Markov Chain model for WLAN Networks is a popular and accurate analytical model
- However, Bianchi's model is for a saturated network. We extend this to model unsaturated networks

- In WLAN, the fundamental mechanism to access the medium is called Distributed Coordination Function (DCF)
- DCF is a random access scheme, based on the Carrier Sense Multiple Access scheme with Collision Avoidance (CSMA/CA)

# Random Backoff of DCF

When there is a packet to be transmitted and if the channel is found to be free, then:

- DCF chooses a backoff time uniformly distributed in the range [0,W-1], where W is called the Contention Window length
- $\bullet\,$  At the first transmission attempt, W is set to  $\,{\cal W}_0\,$
- After each unsuccessful attempt, W is doubled, going up to a maximum value of  $W_m = 2^m W_0$
- After the successful transmission, W is reset back to  $W_0$
- Once a back-off time is chosen, the back-off timer is decremented each time the channel is sensed free, else it is paused until the channel is sensed free again



Figure: Source: G. Bianchi, "Performance analysis of IEEE 802.11 distributed Coordination function", IEEE Journal on Selected Areas in Communications, Vol. 18, No. 3, pp. 535-547, March 2000

## Assumptions

- The main goal of the analysis is to obtain the stationary probability  $\tau$  that a particular transmitter transmits in any randomly chosen slot time. The rest of the parameters (throughput, delay) follow from  $\tau$
- The key approximation in the analysis is that, at each transmission attempt, regardless of the number of transmission attempts, the transmission suffers a constant and independent probability *p*
- All the nodes are assumed to be in saturation, which means a node will always have a packet to transmit

### Evaluation of $\tau$

• From the definition of  $\tau$ :

$$\tau = \sum_{i=0}^m b_{i,0}$$

• After finding the stationary distribution of the Markov Chain:

$$au = rac{2(1-2p)}{(1-2p)(W+1)+pW(1-(2p)^m)}$$

• The probability *p* with which a transmitted packet encounters a collision is the probability that at least one of the other nodes also transmit in the given slot

$$p=1-(1-\tau)^{n-1}$$

## Throughput

• The throughput (R) of the network is defined as the fraction of the time the channel is used to successfully transmit the packet

 $R = \frac{E[\text{Amount of time spent in successful transmission}]}{E[\text{Length of a slot time}]}$ 

$$R = \frac{P_{suc} T_{suc}}{P_{idle} T_{idle} + P_{col} T_{col} + P_{suc} T_{suc}}$$
$$P_{idle} = (1 - \tau)^n$$
$$P_{suc} = n\tau (1 - \tau)^{n-1}$$
$$P_{col} = 1 - n\tau (1 - \tau)^{n-1} - (1 - \tau)^n$$





# Cognitive Radio Setup

- The CR nodes are required to scan the spectrum periodically for the availability of the channel
- All the CR nodes perform the scan procedure simultaneously
- The scan procedure is performed with a period *T*, each scan procedure lasts for a time *t*
- The CR network has T t useful time for its transmission, upon the availability of the channel
- All the nodes in the network (both primary & secondary) are assumed to be in saturation

## MAC Layer issues

- The packet transmission is said to be complete if RTS/CTS and Packet/ACK are successfully exchanged between communicating nodes
- In the Cognitive Radio Network, if the remaining useful time is not sufficient for the complete transmission, we follow certain rules:
  - If the remaining useful time is not sufficient even for RTS transmission, discard the packet. The opportunity to transmit is lost and the node has to contend for the channel again
  - If RTS can fit, but not CTS, the node transmits RTS. Upon the next availability of the channel, the intended receiver has to transmit CTS after a delay of SIFS from the finish of the scan procedure

## Contd...

- If both RTS and CTS can fit, but not the data packet, resume data transmission after SIFS from scan procedure, again upon the next availability of the channel
- If RTS/CTS and the data packet can fit, but not the ACK, the receiver transmits ACK after SIFS upon the next availability of the channel
- All the remaining nodes have to pause their backoff procedure until a duration DIFS from the end of the scan procedure and upon the availability of the channel

### Evaluation of the Cognitive Radio network parameters

- The network can be in one of two possible states:
  - State-1: Secondary nodes have frozen their transmission; only primary nodes are active
  - State-2: All the nodes in the network are active; secondary nodes are also active as the result of the previous scan is *free* 
    - In this case, all nodes (primary + secondary) contend for the channel till the next scan procedure
- Therefore, the network has two sets of transmission parameters, one for each state

#### State-1

 The network has only primary nodes being active. Hence, the transmission (τ<sub>p,1</sub>) and the collision (p<sub>p,1</sub>) probabilities are given by:

$$au_{p,1} = rac{2(1-2p_{p,1})}{(1-2p_{p,1})(1+W_p)+p_{p,1}W_p(1-(2p_{p,1})^{M_p})} 
onumber \ p_{p,1} = 1-(1- au_{p,1})^{N_p-1}$$

- $N_p$  is the number of primary nodes,  $M_p$  is the number of back-off stages of each primary node,  $W_p$  is the initial contention window length of each primary node
- As the secondary nodes are inactive in this state, their transmission probability is zero

#### State-2

• In this state transmission and collision probabilities are given by:

$$\begin{aligned} \tau_{p,2} &= \frac{2(1-2p_{p,2})}{(1-2p_{p,2})(1+W_p) + p_{p,2}W_p(1-(2p_{p,2})^{M_p})} \\ \tau_{s,2} &= \frac{2(1-2p_{s,2})}{(1-2p_{s,2})(1+W_s) + p_{s,2}W_s(1-(2p_{s,2})^{M_s})} \\ p_{p,2} &= 1-(1-\tau_{p,2})^{N_p-1}(1-\tau_{s,2})^{N_s} \\ p_{s,2} &= 1-(1-\tau_{p,2})^{N_p}(1-\tau_{s,2})^{N_s-1} \end{aligned}$$

• *N<sub>s</sub>* is the number of secondary nodes, *M<sub>s</sub>* is the number of back-off stages of each secondary node, *W<sub>s</sub>* is the initial contention window length of each secondary node

# Definition of $\alpha_c$ , $\alpha_{busy}$ , $\alpha_{idle}$

- $\alpha_c \triangleq$  probability with which the network is in State-1
- $\alpha_{busy} \triangleq$  probability with which the network will be in State-1 given that the previous state of network is State-1
- $\alpha_{idle} \triangleq$  probability with which the network will be in State-1 given that the previous state of network is State-2
- $\alpha_c$  can be recursively written as:

$$\alpha_{c} = \alpha_{c}\alpha_{busy} + (1 - \alpha_{c})\alpha_{idle}$$

$$\Rightarrow \alpha_{c} = \frac{\alpha_{idle}}{1 + \alpha_{idle} - \alpha_{busy}}$$

## Importance of $\alpha_c$

- The availability of the channel for the secondary network is determined by the parameter  $\alpha_c$
- $\alpha_c$  reflects the effect of sensing duration and sensing errors, which therefore makes it an important network parameter

## Simulations

 $\alpha_{busy}$  and  $\alpha_{idle}$ :



## Contd...

 $\alpha_{\textit{c}} \text{ and } \tau_{\textit{p},1}$ 



## Contd...

 $\tau_{p,2}$  and  $\tau_{s,2}$ 



## Comparison of Theoretical and Simulated Throughput

- The plots below show the simulated and theoretical achievable throughputs for both the primary and secondary network
- The primary throughput is the fraction of the time spent in transmitting packets successfully, averaged over different states of the network
- As the secondary network is not active in the State-1, its throughput is calculated as the fraction of the time spent in transmitting packets successfully in the useful time T t, when the channel is available





- In the previous section, analysis was done for saturated networks. However it is of practical interest to extend the study to unsaturated networks (particularly when the primary is unsaturated)
- As a first step, a network of only WLAN nodes (the nodes are unsaturated) is modeled with a modified Bianch's Markov model and its compared performance with the simulations.
- The future work includes extending this model to include non-saturated cognitive nodes and study the throughput performance



Figure: Modified Markov model. The parameter  $\lambda$  quantifies the saturation of the network. The greater the  $\lambda$ , the busier the network.

### Transmission Parameters

• The transmission ( $\tau$ ) and the collision probability (p) for the above unsaturated model is given by:

$$\tau = \frac{2(1-2p)}{(1-2p)(CW_o+1) + pCW_o(1-(2p)^M) + 2(1-2p)(1-p)\frac{1-\lambda}{\lambda}}$$
$$p = 1 - (1-\tau)^{n-1}$$

• *CW*<sub>o</sub> is the initial contention window length, *M* is the number of back-off stages and *N* is the number of nodes in the network



Figure: For smaller  $\lambda$ , due to an approximation made in the Markov model (which itself came from Bianchi's Markov model), the theoretical predictions are accurate only for a large number of nodes ( $\approx$  100 nodes). However, even in other regimes, the performance is satisfactory.

## Thank You

Questions?