

Journal Watch

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Prabhasa K



SPC Lab

Department of Electrical Communication Engineering
Indian Institute of Science

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1. Joint Power Optimization for D2D Communication in Cellular Networks With Interference Control

Authors: Ali Ramezani-Kebrya, Min Dong, Ben Liang, Gary Boudreau, and S. Hossein Seyedmehdi

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1. Joint Power Optimization for D2D Communication in Cellular Networks With Interference Control

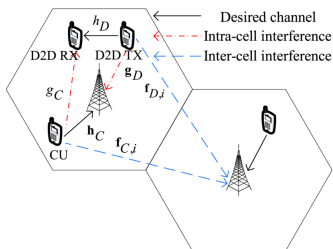
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 - Uplink resource sharing
 - Users and BS - 1, N antennas respectively
 - Orthogonal Resource Allocation among CUs
 - Perfect knowledge of all channels

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- **Assumptions:**
 - Uplink resource sharing
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- **Approach:** (i) Check for D2D admissibility under given constraints (ii) Joint power control under assumption of D2D admission.

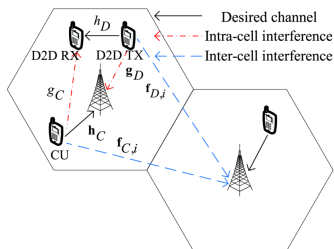


- SINR at D2D RX:

$$\gamma_D = \frac{P_D |h_D|^2}{\sigma_D^2 + P_C |g_C|^2} \quad (1)$$

- SINR at BS from CU:

$$\gamma_C = \frac{P_C |\mathbf{w}^H \mathbf{h}_C|^2}{\sigma^2 + P_D |\mathbf{w}^H \mathbf{g}_D|^2} \quad (2)$$



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- Inter-Cell Interference (ICI)

$$P_{I,i} = P_C \|\mathbf{f}_{C,i}\|^2 + P_D \|\mathbf{f}_{D,i}\|^2. \quad (3)$$

Sum rate optimization problem:

$$\begin{aligned} \mathbf{P1:} \quad & \max_{P_D, P_C, \mathbf{w}} \quad \log(1 + \gamma_C) + \log(1 + \gamma_D) \\ & \text{s.t} \quad \gamma_C \geq \tilde{\gamma}_C, & (4) \\ & \quad \gamma_D \geq \tilde{\gamma}_D, & (5) \\ & \quad P_C \leq P_C^{\max}, P_D \leq P_D^{\max}, & (6) \\ & \quad P_{I,i} \leq \tilde{I}, i = 1, \dots, b & (7) \end{aligned}$$

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Objective: Optimize (P_D, P_C, \mathbf{w})

- Admissibility - For any given (P_D, P_C) , find optimal beam vector \mathbf{w}°
- Optimality - Let $x \triangleq P_D, y \triangleq P_C$. Evaluate (P_D°, P_C°) for $b = 1$ and $b > 1$ (approximation of ICI constraints)

Sum rate optimization problem:

$$\mathbf{P2:} \quad \max_{(x,y)} \quad \log \mathcal{R}(x, y) \quad (8)$$

$$\text{s.t} \quad y \left(1 - \frac{K_1 x}{K_2 + x} \right) I \geq \tilde{\gamma}_C, \quad (9)$$

$$\frac{ax}{\sigma_D^2 + K_3 y} \geq \tilde{\gamma}_D, \quad (10)$$

$$y \leq P_C^{\max}, \quad x \leq P_D^{\max}, \quad (11)$$

$$c_1 y + c_2 x \leq 1 \quad (12)$$

$$\text{where } \mathcal{R}(x, y) \triangleq \left(1 + \frac{ax}{\sigma_D^2 + K_3 y} \right) \left(1 + y \left(1 - \frac{K_1 x}{K_2 + x} \right) I \right),$$

Lemma: The optimal power solution pair (x^o, y^o) is at the vertical, horizontal, or tilted boundary of \mathcal{A}_{xy} , given by $x = P_D^{\max}, y = P_C^{\max}$, or $c_1y + c_2x = 1$, respectively.

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Contributions:

- Proposed algorithm is optimal when ICI to a single neighboring cell is considered. For multiple neighboring cells, we provide an upper bound on the performance loss.
- Formulate an offline (non convex) optimization problem and propose an iterative algorithm to solve it.
- Consider the scenario of multiple CUs and D2D pairs, and formulate the joint power control and CU-D2D matching problem.

Multiple CUs and D2D pairs:

$$\mathbf{P3:} \quad \max_{\mathbf{P}, \mathbf{w}, \mathbf{x}} \quad \sum_{k \in \mathcal{D}} \sum_{j \in \mathcal{C}} \log(1 + \gamma_{C,j}) + x_{k,j} \log(1 + \gamma_{D,k})$$

$$\text{s.t.} \quad \frac{P_{C,j} |\mathbf{w}_j^H \mathbf{h}_{C,j}|^2}{\sigma^2 + x_{k,j} P_{D,k} |\mathbf{w}_j^H \mathbf{g}_{D,k}|^2} \geq \tilde{\gamma}_C, \quad \forall j \in \mathcal{C}$$

$$\frac{P_{D,k} |h_{D,k}|^2}{\sigma_{D,k}^2 + x_{k,j} P_{C,j} |g_{j,k}|^2} \geq \tilde{\gamma}_D, \quad \forall k \in \mathcal{D}$$

$$P_{C,j} \leq P_C^{\max}, \quad P_{D,k} \leq P_D^{\max}, \quad \forall j \in \mathcal{C}, \quad k \in \mathcal{D}$$

$$P_{\mathcal{I},i,j} \leq \tilde{\mathcal{I}}, \quad \forall j \in \mathcal{C}, \quad i = 1, \dots, b$$

$$\sum_{k \in \mathcal{D}} x_{k,j} \leq 1, \quad \sum_{j \in \mathcal{C}} x_{k,j} \leq 1, \quad \forall j \in \mathcal{C}, \quad k \in \mathcal{D}$$

$$x_{k,j} \in \{0, 1\}, \quad \forall j \in \mathcal{C}, \quad k \in \mathcal{D}$$

2. To Motivate Social Grouping in Wireless Networks

Authors: Yu-Pin Hsu, and Lingjie Duan

- **Goal:** Employ social grouping (users with *common interest*) and obtain optimal equal-reciprocal scheme that maximizes the shared content and per-user utility.

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- **Motivation:** Exploiting the hybrid B2D and D2D networks for local content sharing by proposing a non-monetary incentive, assuming selfish users. Traditional incentive schemes are unicast-based.
- **Assumptions:**
 - Users with heterogeneous B2D channels (ON/OFF)
 - No user mobility
 - BS knows channel states
 - BS TX at most 1 packet/slot. Users cannot TX, RX simultaneously
 - B2D, D2D operate on different channels (half-duplex)
 - Users are rational, not malicious

N	Total number of users in consideration
$\mathbf{s}(t) = (s_1(t), \dots, s_N(t))$	Channel state vector at time slot t , where $s_i(t)$ is the channel state for user c_i .
$p_{e,i}$	B2D channel error probability for user c_i , while i is ignored for the symmetric networks.
$T_{\neq}^{(N)}$	Completion time for delivering a common packet employing the unicast communications, while N is ignored for $N = 2$.
$T_{=}^{(N)}$	Completion time for delivering a common packet employing the broadcast communications, while N is ignored for $N = 2$.
$T_{\cup}^{(N)}$	Completion time for delivering a common packet employing the broadcast communications along with social grouping, while N is ignored for $N = 2$.
$p_{i \rightarrow \mathbf{R}}$	Probability that user c_i shares the common packet with the users in \mathbf{R} .
λ	Arrival rate of common interests at the BS.
Λ	Equal-reciprocal stability region.

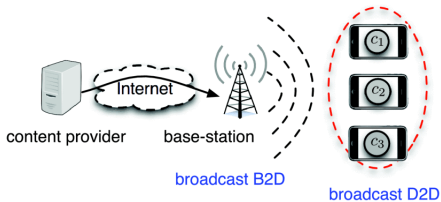


Figure : Hybrid Network Model

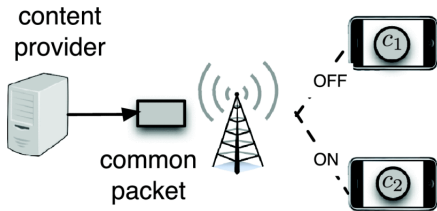


Figure : 2 users with common interest

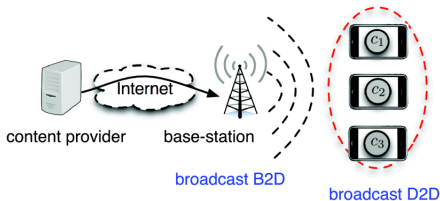


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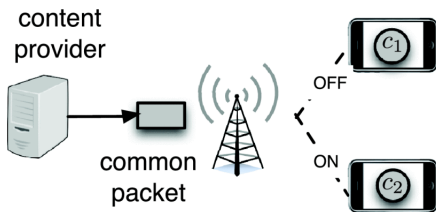


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$$T_{=} = p_e^2(1 + T_{=}) + (1 - p_e)^2 \cdot 1 + 2p_e(1 - p_e) \left(1 + \frac{1}{1 - p_e}\right) = \frac{2p_e + 1}{1 - p_e^2},$$

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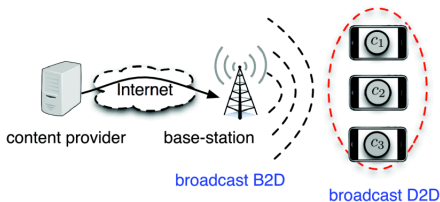


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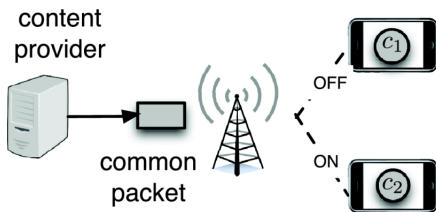


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$$R_{\neq/=} := \frac{T_{\neq}}{T_{=}} = \frac{p_e + 2}{2p_e + 1}.$$

$$R_{=/U} := \frac{T_{=}}{T_U^*} = \frac{2p_e + 1}{-2p_e^2 + 2p_e + 1}.$$

$$T_{\equiv}^{(n)} = \sum_{i=0}^n \binom{n}{i} p_e^i (1 - p_e)^{N-i} (1 + T_{\equiv}^{(i)}),$$

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Lemma: The optimal equal reciprocal scheme for the large symmetric network is to pick a user for sharing with the equal probability, i.e., for user c_i , the sharing probability is

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Lemma: When $p_e < 0.5$, T_{\cup}^* increases in N . When $p_e > 0.5$, T_{\cup}^* decreases in N . When $p_e = 0.5$, $T_{\cup}^* = 2$ for all N . Moreover, $T_{\cup}^* \rightarrow 2$ when $N \rightarrow \infty$, independent of p_e .

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Lemma: The incentivized social group has a performance loss compared with the full cooperation

$$T_{\cup}^* - T_f = \frac{\Delta}{1 - (p_{e,2} - \Delta)p_{e,2}} \left(\frac{1}{1 - p_{e,2}} - 1 \right),$$

where $\Delta = p_{e,2} - p_{e,1}$

Approach:

- Evaluate completion time for different scenarios. Compute improvement ratios.
- Employ equal-reciprocal incentives and illustrate performance for symmetric and asymmetric cases.

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Contributions:

- Consider a practical *MicroCast* network scenario and propose the equal-reciprocal incentive scheme to motivate social grouping, which improved network performance.
- The optimal equal-reciprocal mechanism is a win-win policy that improves the performance of both BSs and local users.
- Propose on-line scheduling algorithms that dynamically select a user to share content (additional 1-bit information).

3. Analytical Characterization of Device-to-Device and Cellular Networks Coexistence

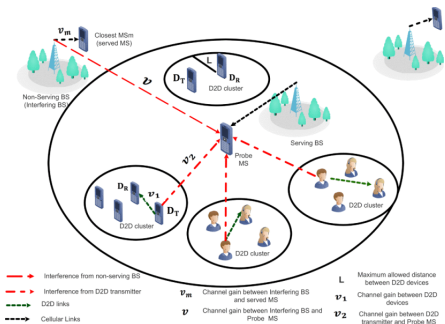
Authors: Ashraf Al-Rimawi and Davide Dardari

Goal: Evaluate the amount of (downlink) traffic offloaded through D2D communications while considering the effect of power control, users spatial distribution, shadowing and random base station deployment.

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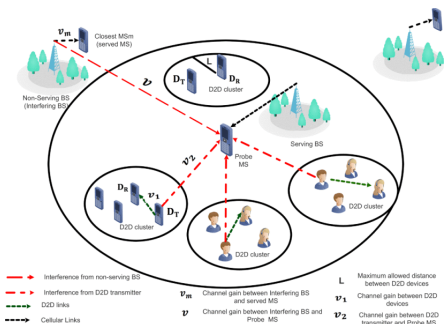
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• Motivation:

- HetNets Interference Mgmt
- Hex cells \rightarrow Sto-Geo
- HPPP unrealistic for D2D

• Assumptions:

- One D2D link active at a time, without violating $SINR_{min}$ constraint
- D2D groups \rightarrow HPPP. D2D users \rightarrow arbitrary spatial model

- **SINR at probe MS:**

$$\text{SINR} = \frac{P_u}{\sum_{i=0}^n P_i + \sum_{j=0}^m P_j^{(d)} + \sigma_0^2},$$

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- **Coverage Probability:**

$$\begin{aligned} P_c &= \text{Prob}(\text{SINR} > \eta) = \text{Prob}\left(\sum_{i=0}^n P_i + \sum_{j=0}^m P_j^{(d)} < \gamma\right) \\ &= \text{Prob}(10 \log_{10}(P_I) < 10 \log_{10} \gamma) \\ &\approx \text{Prob}\left(\max_{i,j} (P_i \text{ (dBm)}, P_j^{(d)} \text{ (dBm)}) < 10 \log_{10} \gamma\right), \end{aligned}$$

- Coverage Probability averaged over n, m :

$$\begin{aligned} P_{c_0} &= \mathbb{E}[P_{c_0}|n,m] \\ &= \sum_{m=0}^{\infty} \sum_{n=0}^{\infty} P_{c_0|n,m} \frac{(\rho_{BS}\pi R_0^2)^n}{n!} \frac{(\rho_D\pi R_0^2)^m}{m!} \\ &\quad \times \exp(-(\rho_{BS} + \rho_D)\pi R_0^2) \\ &= \exp(-\rho_D\pi R_0^2(1 - F_D(\gamma)) - \rho_{BS}\pi R_0^2(1 - F_I(\gamma))), \end{aligned}$$

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- Coverage Probability over the entire area:

$$P_c = \lim_{R_0 \rightarrow \infty} P_{c_0} = \exp(-\lambda_{BS}(\gamma) - \lambda_D(\gamma)).$$

Contributions:

- Presented a new analytical framework for analyzing the coverage probability in coexisting cellular and D2D networks.
- Characterizing channel gain statistics of serving, non-serving BSs, and D2D links.
- The reciprocal impact of D2D and cellular communications on the downlink coverage is investigated as a function of the D2D links maximum range and density.

4. Phase Retrieval Motivated Nonlinear MIMO Communication With Magnitude Measurements

Authors: Shengchu Wang, Lin Zhang, and Xiaojun Jing

- **Goal:** Solve the quantized PR problem (inherently - CE and MUD) for a NL-MIMO scheme.

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 - In the uplink TX, MO-MIMO must complete CE and MUD based on the quantized magnitude measurements at the BS
 - Developed CE and MUDs tested for various noise conditions, ADC resolutions, lengths of channel training sequences, and numbers of BS antennas

Contributions:

- Propose a low-power and low-cost NL-MIMO scheme (MO-MIMO).








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- Propose a low-power and low-cost NL-MIMO scheme (MO-MIMO).
- Practical CEs and MUDs were developed for MO-MIMO by categorizing the two problems as a quantized phase retrieval (PR) problem.
- Solved the PR problem by constructing two algorithms under the framework of generalized approximate message passing (GAMP).

Other Interesting Papers

-  Mobility-Aware Caching in D2D Networks.
-  Full-Duplex Massive MIMO Relaying Systems With Low-Resolution ADCs.
-  Distributionally Robust Collaborative Beamforming in D2D Relay Networks With Interference Constraints.
-  On the Performance of Beam Division Nonorthogonal Multiple Access for FDD-Based Large-Scale Multi-User MIMO Systems.
-  Feedback Mechanisms for FDD Massive MIMO With D2D-Based Limited CSI Sharing.
-  Secret Key Generation Based on Estimated Channel State Information for TDD-OFDM Systems Over Fading Channels
-  Millimeter-Wave Channel Estimation Based on 2-D Beamspace MUSIC Method