#### Transmission Capacity of Wireless Ad Hoc Networks

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Workshop on Models and Protocols for Mobile Ad-Hoc Networks

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### Motivation

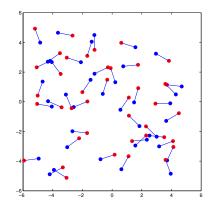
- A set of links (TX and RX) are randomly distributed in  $\mathbb{R}^2$ .
- A MAC protocol activates a subset of these links
  - ALOHA, CSMA, TDMA
- All nodes transmit at the same power
- A scheduled link succeeds if its SIR  $> \theta$

#### Questions

- 1. How to quantify the efficiency (packing) of the MAC protocol?
- 2. Better signal processing or better MAC?

### Dipole model

- $\blacktriangleright$  The links are distributed as a stationary point process  $\Phi \subset \mathbb{R}^2$  of unit
- Each transmitter has a receiver at a distance d in a random direction
- Path loss function is denoted by ℓ(x). Eg. ℓ(x) = ||x||<sup>-α</sup>
- A subset of these nodes Φ<sub>t</sub> ⊂ Φ are activated by a MAC protocol
  - Density:  $\lambda_t$



Quantify performance of MAC and its interaction with the signal processing

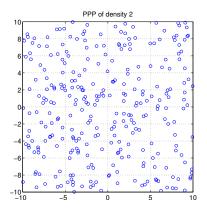
Some models for spatial locations of the links

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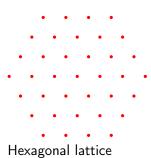
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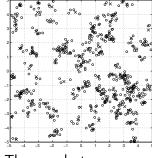
# Spatial poisson point process

- 1. The most widely used model for spatial locations of nodes
  - Most amicable to analysis
  - "Gaussian of point processes"
- 2. No dependence between node locations
- 3. Random number of nodes
- 4. Can be defined on the entire plane

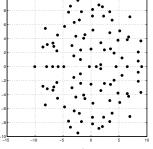


## Other interesting examples





Thomas cluster process



Determinantal point process (eigenvalues of a Gaussian matrix)

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Examples of MAC schemes

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### ALOHA

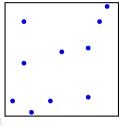
Let  $\Phi$  be a PP of density  $\lambda$ 

- 1. A node (link)  $x \in \Phi$  is scheduled with probability p independent of other nodes.
- 2. The transmitter process  $\Phi_t$  is stationary with density  $\lambda p$ .

#### 1. Begin with a PP $\Phi$ of density $\lambda$ .

- To each x ∈ Φ, associate a mark m<sub>x</sub> ~ U[0,1] independent of every other point.
- A node x ∈ Φ selected if it has the lowest mark among all the points in the ball B(x, R).

 $\Psi = \{ y : y \in \Phi, m_y \le m_x, \forall x \in B(y, R) \cap \Phi \}$ 

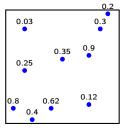


#### A minimum distance process for modelling CSMA MAC.

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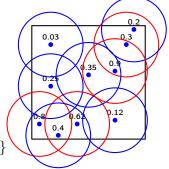


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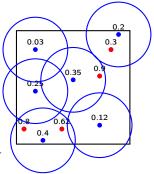
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A minimum distance process for modelling CSMA MAC.

### m-Phase TDMA

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GRK (IITM)

Transmission capacity

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# Signal processing techniques

- TX has N<sub>t</sub> transmit antenna
- RX has N<sub>r</sub> receive antenna
- MRC, spatial multiplexing....

# Link formation probability

Signal-to-interference ratio

$$SIR(x) = \frac{h_{xr(x)}\ell(d)}{\sum_{y \in \Phi_t} h_{yr(x)}\ell(y - r(x))}$$

The success probability of a typical link

$$P_{c}(\theta,\lambda) = \mathbb{P}^{!o}(\mathsf{SIR}(o) > \theta)$$

Reduced Palm probability:

- picking a link at random
- average over all the links
- conditioning on a point at the origin (but not counting it)

# Metric: Transmission capacity(TC)

Let  $\epsilon \in (0, 1)$ . TC is defined as

$$TC(\epsilon) = (1 - \epsilon) \max\{\lambda : p_s(\theta, \lambda) > 1 - \epsilon\}$$

- 1. TC measures the maximum spatial intensity of successful transmissions per unit area for a given success probability.
- 2. Can be related to area spectral efficiency (ASE) by multiplying with  $\log_2(1+\theta)$ .

We are interested in the behaviour (scaling) of  $TC(\epsilon)$  when  $\epsilon \to 0$ 

### ALOHA MAC

Node distribution:  $\Phi \sim PP(1)$ MAC: ALOHA with parameter *p*.  $\Phi_t$  is a stationary PP with density *p*.

$$p_{s}(\theta, \lambda) = \mathbb{P}^{!o}(\mathsf{SIR}(o, r(o)) > \theta)$$
$$= \mathbb{P}^{!o}\left(\frac{S\ell(d)}{I_{t}} > \theta\right) = \mathbb{E}^{!o}\left[F_{s}\left(\theta\frac{I_{t}}{\ell(d)}\right)\right]$$

Here  $I_t = \sum_{x \in \Phi} \mathbf{1}(x \text{ is selected}) h_x \ell(x - r(o))$ 

### Basic idea: 2 node network

Let 
$$Y_i = \frac{\theta h_{x_i}}{\ell(d)} l(x_i - r(o)), i = 1, 2$$
. Then  
 $\hat{l}_t = Y_1 \mathbf{1}(1 \text{ is on}) + Y_2 \mathbf{1}(2 \text{ is on}).$ 

Hence averaging over the ALOHA selection

$$\mathbb{E}_{p}[F_{s}(\hat{I}_{t})] = (1-p)^{2}F_{s}(0) + p(1-p)(F_{s}(Y_{1}) + F_{s}(Y_{2})) + p^{2}F_{s}(Y_{1} + Y_{2})$$
$$\sim 1 - p\sum_{i=1}^{2}(1 - F_{s}(Y_{i})) + o(p)$$

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### ALOHA: Asymptotic success probability

#### Theorem (RK-MH-JA)

When ALOHA is used as the MAC protocol, the density of  $\Phi_t$  is  $\lambda_t = p$  and

$$P_s(\theta, \lambda_t) \sim 1 - \lambda_t \gamma_{ALOHA}, \ \lambda_t \to 0,$$

where

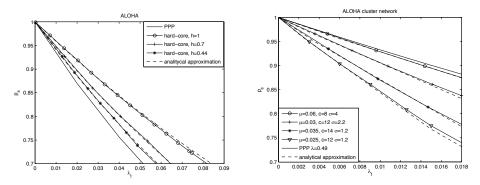
$$\gamma_{ALOHA} = \int_{\mathbb{R}^2} \left[ 1 - E_h F_s \left( h \frac{\theta \ell(x - r(o))}{\ell(d)} \right) \right] \rho^{(2)}(x) \mathrm{d}x$$

$$egin{split} 
olimits p_s( heta,\lambda) &= \mathbb{E}^{!o}\left[F_s\left(rac{ heta}{\ell(d)}I_t
ight)
ight] \ &\sim 1 - p\mathbb{E}^{!o}\left[\sum_{x\in \Phi}1 - F_s\left(rac{ heta}{\ell(d)}h_x\ell(x-r(o))
ight)
ight] + o(p) \end{split}$$

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### How good is the approximation



ALOHA on a hard-core network; Right: ALOHA on a clustered network

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Transmission capacity

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# ALOHA: Transmission capacity

$$\mathcal{TC}_{\mathsf{ALOHA}}(\epsilon)\simrac{\epsilon}{\gamma_{\mathsf{ALOHA}}},\quad\epsilon
ightarrow \mathsf{0}$$

- $\blacktriangleright$  Scales linearly with  $\epsilon$
- Only  $\gamma_{ALOHA}$  depends on  $F_s(.)$ , h and  $\Phi$

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Transmission capacity

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 $\gamma_{\mathsf{ALOHA}}$  for a PPP with  $\ell(x) = \|x\|^{-lpha}$ ,  $\delta \triangleq 2/lpha$ 

1. Nakagami-*m* fading with  $N_t = N_r = 1$ 

$$\gamma_{ ext{ALOHA}} = d^2 rac{\pi heta^\delta \Gamma(m-\delta) \Gamma(m+\delta)}{\Gamma(m)^2} \sim \Theta(1)$$

2. Receiver beamforming:  $N_t = 1$ ,  $N_r = m$ 

$$\gamma_{\mathsf{ALOHA}} = d^2 rac{\pi heta^\delta \Gamma(m-\delta) \Gamma(1+\delta)}{\Gamma(m)} \sim \Theta(m^{-\delta})$$

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# Sphere packing interpretation of TC

$$TC(\epsilon) \sim rac{\epsilon}{d^2 C(\delta)} = rac{1}{\pi \left( d \sqrt{rac{C(\delta)}{\pi \epsilon}} 
ight)^2}.$$

#### Interpretation (heuristic)

Hence each transmission approximately requires an interference free disc of radius

$$R = d\sqrt{\frac{\mathcal{C}(\delta)}{\pi\epsilon}}$$

- The disc radius increases as  $\frac{1}{\sqrt{\epsilon}}$ .
- $\blacktriangleright$  The disc radius decreases with increasing  $\alpha$ 
  - ► Higher path loss exponent → better packing.

### CSMA MAC

Node distribution:  $\Phi \sim \text{PPP}(1)$ MAC: CSMA with radius R.  $\Phi_t$  is a stationary PP with density  $\lambda_t = \frac{1 - \exp(-\pi R^2)}{\pi R^2}$ .

$$p_{s}(\theta, \lambda) = \mathbb{P}^{!o}(\mathsf{SIR}(o, r(o)) > \theta)$$
$$= \mathbb{P}^{!o}\left(\frac{S\ell(d)}{I_{t}} > \theta\right) = \mathbb{E}^{!o}\left[F_{s}\left(\theta\frac{I_{t}}{\ell(d)}\right)\right]$$

Here  $I_t = \sum_{x \in \Phi} \mathbf{1}(x \text{ is selected}) h_x \ell(x - r(o))$ 

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### Basic idea

- ▶ Minimum distance between nodes  $\mathbb{E}[I] < \infty$ . Not true for ALOHA
- Taylor series of CCDF of the desired signal

$$F_s(x) = 1 - c_0 x^{\nu} + o(x^{\nu}), x \to 0.$$

• Average nearest interferer distance  $\approx 1/\sqrt{\lambda_t}$ .

$$F_s(I) \approx 1 - c_0 \lambda^{\alpha \nu/2}.$$

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## CSMA: Asymptotic success probability and TC

#### Theorem

For a Matern CSMA process,

$$P_s( heta, \lambda_t) \sim 1 - c_0 \pi \left(rac{ heta}{\ell(d)}
ight)^{
u} \lambda_t^{
u/\delta} A_I, \lambda_t 
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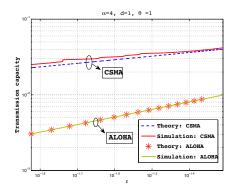
- $A_I$  is known and depends on the higher order moment densities ( $\nu + 1$  of them) of the Matern process
- Difficult to evaluate except for  $\nu = 1$

 $\begin{aligned} & \text{Transmission capacity} \\ & \mathcal{T}C_{\mathsf{CSMA}}(\epsilon) \sim \frac{\epsilon^{\delta/\nu}}{(c_o \pi A_I)^{\delta/\nu} \theta^{\delta} \ell(d)^{-\delta}}, \quad \epsilon \to 0 \end{aligned}$ 

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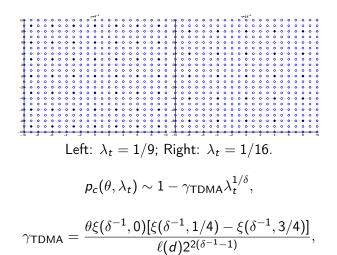
### Observations

- Better TC  $\Theta(\epsilon^{\delta/\nu})$
- Interacts well with signal processing techniques
  - Nakagami-m: ν = m
  - MRC:  $\nu = N_r$



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# m-phase TDMA with $F_s(x) = \exp(-x)$



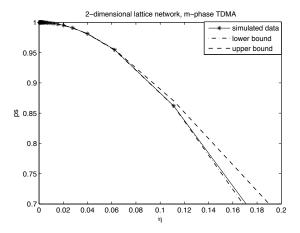
and  $\xi(s, b)$  is the generalized Riemann zeta function.

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where

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### m-phase TDMA



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Transmission capacity

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### Reasonable MAC

A.1 lim<sub>λ→0</sub> E<sup>lo</sup>(Points of Φ in B(o, λ<sup>-a</sup>)) = 0 for all a < 1/2.</li>
A.2 There exists some constant R such that lim<sub>λ→0</sub> E<sup>lo</sup>(Points of Φ in B(o, Rλ<sup>-1/2</sup>)) > 0.

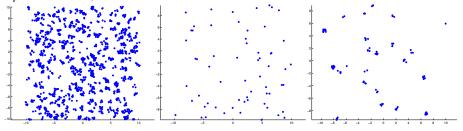
#### Definition

Any MAC scheme which satisfies A.1 and A.2 is a reasonable MAC scheme

Example: PPP with ALOHA is a reasonable MAC scheme.

### An unreasonable MAC scheme

Consider a cluster process with average number of clusters per unit area as  $\lambda_p$  and nodes per cluster as  $\bar{c}$ .



$$\lambda_p = 1, \bar{c} = 15, n = 25$$

- ► Using ALOHA select a node with probability  $n^{-1}$ . Nodes per cluster:  $\bar{c}n^{-1}$ , Density:  $\bar{c}\lambda_p n^{-1}$ .
- Allow a cluster to transmit with probability  $n^{-1}$ . Cluster density:  $\lambda_p n^{-1}$ , Density:  $\bar{c}\lambda_p n^{-1}$ .

### m-Phase TDMA

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### Bounds on scaling exponent

For a reasonable MAC

$$p_s(\Theta, \lambda_t) \sim 1 - \gamma \lambda_t^{\kappa}, \quad \lambda_t \to 0$$

 $\kappa$ : Interference scaling exponent.  $\gamma$ : Spatial contention parameter

Theorem

For any reasonable MAC protocol,

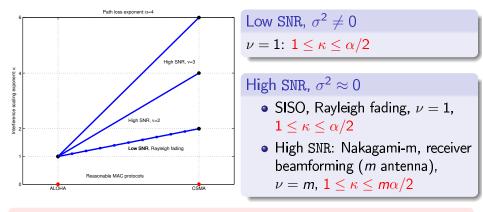
$$1 \le \kappa \le \frac{\nu}{\delta},$$

- $\kappa = 1$ : ALOHA
- $\kappa = \frac{\nu}{\delta}$ : CSMA

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### Interference scaling exponent $\kappa$

$$\mathbf{F}_{s}\left(\sigma^{2}+x\right)=\mathbf{F}_{s}\left(\sigma^{2}\right)+C_{0}x^{\nu}+o(x^{\nu}),x\rightarrow0$$



 $1 \leq \kappa \leq rac{
ulpha}{2}$ , exponent  $\kappa$  depends on the diversity order of the receiver

GRK (UTA)

### Key results

- 1. For any reasonable MAC scheme, the interference scaling exponent  $1 \leq \kappa \leq \nu/\delta$
- 2. ALOHA always leads to  $\kappa = 1$
- 3. CSMA leads to  $\kappa = \nu/\delta$
- 4. Interference cancellation (CSMA) is important to take advantage of advanced signal processing techniques
- 5. Transmission capacity provides a clean framework for the analysis of networks