## भारतीया विज्ञान संस्थान <br> The Pattern Maximum Likelihood Estimation Problem

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## Overview

- Estimating properties of Markov chains and memoryless sources
- Symmetric properties and performance of plug-in estimators
- Pattern maximum likelihood (PML) estimate
- Approximating the PML estimate using a variational approach

Estimating the transition matrix of a DTMC

## An estimation problem

Suppose we have $X_{1}, X_{2}, X_{3}, \ldots, X_{n}$ from an irreducible time-homogeneous Markov chain over $\mathcal{S}=\{1,2, \ldots, k\}$ with transition kernel

$$
p_{x, y}=\operatorname{Pr}\left[X_{t+1}=y \mid X_{t}=x\right]
$$

and uniform initial distribution.
We know $k$, but we do not know $p$.

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What can we infer about $p$ ?
Regime of interest: $n \leq k^{2}$.

## Estimating graphs from random walks

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Q: What can we infer about $\mathcal{G}$ from $X_{1}, X_{2}, \ldots, X_{n}$ ?
This is important in the regime where $n$ is less than $k^{2}$.
Many parameters such as the degree distribution, eigenvalues of the adjacency matrix, etc., are of interest.

## A simpler problem: the i.i.d. case

We have a pmf $p$ over $\mathcal{S}=\{1,2, \ldots, k\}$
We observe $X_{1}, \ldots, X_{n}$, i.i.d. with each $X_{i} \sim p$.
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$$
\operatorname{Pr}\left[X^{n}=x^{n}\right]=\prod_{i=1}^{n} p_{x_{i}}=\prod_{a \in \mathcal{S}} p_{a}^{\mu_{a}}
$$

where $\mu_{a}$ is the number of times $a$ appears in $x^{n}$.

## Sequence maximum likelihood estimation

If $n$ is large enough, can find the empirical estimate of $p$ (SML estimate):
For $a \in \mathcal{S}$, let $\mu_{a}$ denote the number of times the symbol a occurs in $x_{1}, x_{2}, \ldots, x_{n}$.

$$
\left(p_{\mathrm{SML}}\right)_{a}=\frac{\mu_{a}}{\sum_{b \in \mathcal{S}} \mu_{b}}=\frac{\mu_{a}}{n}
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$$

Problem: If $n \lesssim k$, we do not get a good estimate.

- If $n<k$, some symbols will never be observed.
- The SML estimate assigns zero probability to such symbols.


## Symmetric properties of distributions

- $f(p)$ is symmetric if it is invariant to a relabeling of the alphabet.
- For every $\sigma \in S_{k}, f\left(p_{\sigma(\cdot)}\right)=f(p)$.
- Examples: Support size, entropy (Shannon, Renyi), etc.

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H(p)=-\sum_{a} p_{a} \log _{2} p_{a}
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- Want to estimate $f(p)$ from $X_{1}, \ldots, X_{n}$.
- Specifically, for $\epsilon, \delta>0$, want an estimator $\hat{f}: \mathcal{S}^{n} \rightarrow \mathbb{R}$ such that

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\operatorname{Pr}\left[\left|f(p)-\hat{f}\left(X^{n}\right)\right|>\epsilon\right]<\delta
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- Sample complexity: smallest $N$ such that the above holds for all $n \geq N$. Typically take $\delta=1 / 3$.


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## Estimating symmetric properties: A plug-in approach?

- Estimating $f(p)$ : Use ML/favourite estimator. Different estimator for each $f$. Complexity??
- Idea: Find an approximation of $p$, i.e., $\hat{p}$. Compute $f(\hat{p})$ - plug-in estimator.
- SML plug-in estimator: Choose $\hat{p}=p_{\text {SML }}$.
- Problem: If $n$ is small compared to $k$, then $p_{\text {SML }}$ is bad.
- Q: Can we do better than the SML estimate?


## The Pattern Maximum Likelihood Estimate

## An alternative to $p_{\text {SML }}$

## Pattern:

- Given $\mathbf{x}=x_{1}, x_{2}, \ldots, x_{n}$, the index of symbol $a$ in $\mathbf{x}$ is 1 plus the number of distinct symbols occurring before the first occurrence of $a$ in x .
- The pattern of x is the string obtained by replacing $x_{i}$ by the index of $x_{i}$.


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Example: Consider
$\mathbf{x}=$ abracadabra.
The pattern of $\mathbf{x}$,

$$
\psi(\mathbf{x})=12314151231
$$

| Symbol | Index |
| :---: | :---: |
| $a$ | 1 |
| $b$ | 2 |
| $r$ | 3 |
| $c$ | 4 |
| $d$ | 5 |

## Profile:

multiset of number of occurrences of different symbols

$$
\left\{\mu_{1}, \mu_{2}, \ldots, \mu_{n}\right\}
$$

Profile of abracadabra: $\{5,2,2,1,1\}$

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Pattern probability:

$$
\mathbb{P}(\boldsymbol{\psi} \mid p) \triangleq \sum_{\sigma} \prod_{i=1}^{k} p_{\sigma(i)}^{\mu_{i}}
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$$
\mathbb{P}(12314151231 \mid p)=\operatorname{Pr}[\text { abcadaeabca }]+\cdots+\operatorname{Pr}[\text { abracadabra }]+\cdots
$$

## An alternative to $p_{\mathrm{SML}}$ : The PML estimate

## SML and PML estimates

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For convenience, maximize over ordered pmfs, i.e., $p_{1} \geq p_{2} \geq \ldots \geq p_{k}$.

$$
\begin{align*}
p_{\mathrm{PML}}^{(\psi)} & =\arg \max _{p \in \mathcal{P}_{k}} \mathbb{P}(\psi \mid p) \\
& =\arg \max _{p \in \mathcal{P}_{k}} \sum_{\sigma} \prod_{i=1}^{k} p_{\sigma(i)}^{\mu_{i}} \tag{1}
\end{align*}
$$

## Origins of PML

- Origins of PML: Universal compression of memoryless sources over unknown alphabets by Orlitsky et al. ${ }^{1}$
${ }^{1}$ A. Orlitsky, N. Santhanam, and J. Zhang, "Universal compression of memoryless sources over unknown alphabets," IEEE Trans. Inf. Theory, Jul. 2004


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- Universal compression: block redundancy (average number of additional bits required compared to the case when distribution is known)

$$
R(\mathcal{P})=\inf _{q} \sup _{p} \sup _{x \in \mathcal{S}} \log \frac{p(x)}{q(x)}
$$

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- For sequences, block redundancy

$$
R\left(I_{k}^{n}\right)=\frac{k-1}{2} \log \frac{n}{2 \pi}+\log \left(\frac{\Gamma(1 / 2)^{k}}{\Gamma(k / 2)}\right)+o_{k}(1)
$$

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- (Orlitsky et al.) For compressing patterns, block redundancy

$$
(1.5 \log e) n^{1 / 3}(1+o(1)) \leq R\left(I_{\psi}^{n}\right) \leq \pi \sqrt{2 / 3}(\log e) \sqrt{n}
$$

[^2]
## Estimating symmetric properties using $p_{\mathrm{PML}}$

PML plug-in estimator: Compute $p_{\text {PML }}$, and find $f\left(p_{\text {PML }}\right)$.

[^3]
## Estimating symmetric properties using $p_{\mathrm{PML}}$

PML plug-in estimator: Compute $p_{\text {PML }}$, and find $f\left(p_{\text {PML }}\right)$.
Let $\mathcal{Z}^{(n)}$ denote the set of all length- $n$ patterns.

## Proposition (Acharya et al. ${ }^{1}$ )

Consider any estimator $\hat{f}$ for $f$ that takes as input ${ }^{2} \psi\left(\mathbf{X}^{(n)}\right)$. Suppose that for every $\epsilon>0, \delta>0$ and transition probability distribution $p$, there exists $N$ such that

$$
\operatorname{Pr}\left[\left|f(p)-\hat{f}\left(\psi\left(\mathbf{X}^{(n)}\right)\right)\right| \geq \epsilon\right]<\delta
$$

for all $n \geq N$. Then,

$$
\operatorname{Pr}\left[\left|f\left(p_{\mathrm{PML}}^{\left(\psi\left(\mathbf{X}^{(n)}\right)\right)}\right)-f(p)\right| \geq 2 \epsilon\right]<\delta \cdot\left|\mathcal{Z}^{(n)}\right|
$$

for all $n \geq N$.

$$
\left|\mathcal{Z}^{(n)}\right| \leq \min \left\{e^{3 \sqrt{n}},\binom{n+k-1}{k-1}\right\}
$$

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| Property | SML | Optimal |
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| Entropy | $O(k / \epsilon)$ | $O\left(\frac{k}{\epsilon \log k}\right)$ |
| Support size* | $O(k \log (1 / \epsilon))$ | $O\left(\frac{k}{\log k} \log ^{2}(1 / \epsilon)\right)$ |

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- Basic idea: There exist optimal estimators that give bias $\epsilon$, error probability $1 / 3$, and satisfy a "bounded difference property"
- Have sample complexity $N_{s}$
- Use McDiarmid's inequality to show that probability of error $e^{-\Omega(\sqrt{n})}$ can be achieved using $O\left(N_{s}\right)$ samples
- Then, use previous result


## Efficiently approximating $p_{\text {PML }}$

## Computing $p_{\mathrm{PML}}$

$$
p_{\mathrm{PML}}^{(\psi)}=\arg \max _{p \in \mathcal{P}} \sum_{\sigma} \prod_{a=1}^{k} p_{\sigma(a)}^{\mu_{a}}
$$

## Computing $p_{\mathrm{PML}}$

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Permanent: Given $k \times k$ matrix $M=\left(m_{i, j}\right)$,

$$
\operatorname{perm}(M)=\sum_{\sigma \in S_{k}} \prod_{i=1}^{k} a_{i, \sigma(i)}
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$$

Determinant: Given $k \times k$ matrix $M=\left(m_{i, j}\right)$,

$$
\operatorname{det}(M)=\sum_{\sigma \in S_{k}}(-1)^{\operatorname{sgn}(\sigma)} \prod_{i=1}^{k} a_{i, \sigma(i)}
$$

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$$

Pattern probability $=\operatorname{perm}\left(\left(p_{i}^{\mu_{j}}\right)\right)$
Computing permanent is hard!
For $0-1$ matrix, best known Ryser's algorithm requires $O\left(k 2^{k}\right)$ operations. We use a variational approach as done by Vontobel ${ }^{a}$.

[^5]
## A variational approach: Reformulating $p_{\mathrm{PML}}$

$$
Z \triangleq \mathbb{P}(\psi \mid p)=\sum_{\sigma \in K} \prod_{i, j} p_{i}^{\mu_{j} \sigma_{i j}}
$$

where $\sigma_{i j}$ is the $(i, j)$ th entry of the permutation matrix $\sigma$
Objective: Express this as the minimum of a certain free energy function.

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$$

where $\sigma_{i j}$ is the $(i, j)$ th entry of the permutation matrix $\sigma$
Introduce a "trial" distribution $\beta$ on all permutations on $\{1,2, \ldots, k\}$.
Define the Gibbs average energy function

$$
\begin{aligned}
U_{\mathrm{G}}(\beta ; p, \boldsymbol{\psi}) & \triangleq-\sum_{\boldsymbol{\sigma} \in \mathcal{K}} \beta(\boldsymbol{\sigma}) \log \left(\prod_{i, j} p_{i}^{\mu_{j} \sigma_{i j}}\right) \\
& =-\sum_{\boldsymbol{\sigma} \in \mathcal{K}} \sum_{i, j} \beta(\boldsymbol{\sigma}) \sigma_{i j} \log \left(p_{i}^{\mu_{j}}\right),
\end{aligned}
$$

and the Gibbs entropy function

$$
H_{\mathrm{G}}(\beta) \triangleq-\sum_{\boldsymbol{\sigma} \in \mathcal{K}} \beta(\boldsymbol{\sigma}) \log \beta(\boldsymbol{\sigma}) .
$$

## A variational approach: Reformulating $p_{\mathrm{PML}}$

$$
\begin{gathered}
U_{\mathrm{G}}(\beta ; p, \boldsymbol{\psi})=\log k-\sum_{\boldsymbol{\sigma} \in \mathcal{K}} \sum_{i, j, l, m} \beta(\boldsymbol{\sigma}) \log \left(p_{l, m}^{\mu_{i, m} \sigma_{i j} \sigma_{j m}}\right), \\
H_{\mathrm{G}}(\beta)=-\sum_{\boldsymbol{\sigma} \in \mathcal{K}} \beta(\boldsymbol{\sigma}) \log \beta(\boldsymbol{\sigma}) .
\end{gathered}
$$

We define the Gibbs free energy function

$$
F_{\mathrm{G}}(\beta ; p, \boldsymbol{\psi}) \triangleq U_{\mathrm{G}}(\beta ; p, \boldsymbol{\psi})-H_{\mathrm{G}}(\beta),
$$

It is a fact that ${ }^{2}$

$$
\min _{\beta} F_{\mathrm{G}}(\beta ; p, \boldsymbol{\psi})=-\log Z=-\log \mathbb{P}(\psi \mid p)
$$

Therefore,

$$
p_{\mathrm{PML}}^{(\psi)}=\arg \min _{p \in \mathcal{C}} \min _{\beta \in \mathcal{P}} F_{\mathrm{G}}(\beta ; p, \boldsymbol{\psi}) .
$$

[^6]
## Interpretation

Thermodynamic system with state space $\mathcal{K}$. Probability that the system is in state $\sigma$ :

$$
\gamma(\sigma)=\frac{e^{-E(\sigma) / T}}{Z}
$$

where

- $E: \mathcal{K} \rightarrow \mathbb{R}$ is the energy function (Hamiltonian)
- $T$ is the temperature, $\kappa$ is Boltzmann's constant $\left(1.38 \times 10^{-23} \mathrm{JK}^{-1}\right)$
- $Z=\sum_{\sigma} e^{-E(\sigma) /(\kappa T)}$ is the Helmholtz free energy


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In our case,

- $E(\sigma)=\sum_{i, j} \sigma_{i, j} \log p_{i}^{\mu_{j}}$
- $\kappa T=1$
- $Z=\mathbb{P}(\boldsymbol{\psi} \mid p)$


## Interpretation

The Helmholtz average energy function

$$
U_{\mathrm{H}}(\gamma ; E) \triangleq \sum_{\sigma} \gamma(\sigma) E(\sigma)
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and the Helmholtz entropy function

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H_{\mathrm{H}}(\gamma) \triangleq-\sum_{\boldsymbol{\sigma} \in \mathcal{K}} \gamma(\boldsymbol{\sigma}) \log \gamma(\boldsymbol{\sigma}) .
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and $F_{G}=U_{G}-T H_{G}$

## A variational approach to approximating $p_{\mathrm{PML}}$

$$
p_{\mathrm{PML}}^{(\psi)}=\arg \min _{p} \min _{\beta} F_{\mathrm{G}}(\beta ; p, \psi)
$$

But how do we compute this?

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Idea: Use approximations that are easy to compute ${ }^{3}$. Specifically, perform minimization w.r.t. $\beta$ over an easier set.

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Idea: Use approximations that are easy to compute ${ }^{3}$.
Specifically, perform minimization w.r.t. $\beta$ over an easier set.

Mean field approximation: Choose $\beta$ to be a product distribution. Easy to compute.

## Bethe approximation:

Typically use low-complexity belief propagation algorithms.

[^9]
## Generalization to DTMCs

## An alternative to $p_{\mathrm{SML}}$ : The PML estimate

## SML and PML estimates

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## Pattern probability:

$$
\mathbb{P}(\boldsymbol{\psi} \mid p) \triangleq \frac{1}{k} \sum_{\sigma} \prod_{i=1}^{k} \prod_{j=1}^{k} p_{\sigma(i), \sigma(j)}^{\mu_{i j}}
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PML estimate:

$$
p_{\mathrm{PML}}^{(\psi)}=\arg \max _{p} \mathbb{P}(\boldsymbol{\psi} \mid p) .
$$

## The traditional mean field approximation

$$
p_{\mathrm{PML}}^{(\psi)}=\arg \min _{p \in \mathcal{P}} \min _{\beta \in \mathcal{P}^{\prime}} F_{\mathrm{G}}(\beta ; p, \psi)
$$

Choose $\beta$ to be a product distribution on $k \times k$ binary matrices, i.e., $\beta(\sigma)=\prod_{i, l} \beta_{i l}\left(\sigma_{i l}\right)$.

$$
\begin{gathered}
F_{\mathrm{TMF}}(\beta ; p, \boldsymbol{\psi})=-\sum_{\boldsymbol{\sigma} \in\{0,1\}^{k \times k}}\left(\left(\prod_{i, l} \beta_{i l}\left(\sigma_{i l}\right)\right) \log \left(1_{\mathcal{K}}(\boldsymbol{\sigma}) \prod_{i, j, l, m} p_{l, m}^{\mu_{i j} \sigma_{i l} \sigma_{j m}}\right)\right) \\
+\sum_{i, l} \sum_{\sigma_{i l}=0}^{1} \beta_{i l}\left(\sigma_{i l}\right) \log \beta_{i l}\left(\sigma_{i l}\right)+\log k .
\end{gathered}
$$

The traditional mean-field PML estimate is

$$
p_{\mathrm{TMFPML}}^{(\psi)}=\arg \min _{p \in \mathcal{C}} \min _{\beta} F_{\mathrm{TMF}}(\beta ; p, \psi)
$$

However, we show that this actually reduces to the SML estimate.

## A modified mean field estimate

Inspired by mean field approach used by Chertkov and Yedidia ${ }^{4}$ for approximating permanent of a nonnegative matrix.

In the MF approximation, impose constraint that $\sum_{l} \beta_{i /}(1)=\sum_{i} \beta_{i /}(1)=1$. Define $b_{i l} \triangleq \beta_{i l}(1)$.

$$
\begin{align*}
& F_{\mathrm{MF}}(\cdot ; p, \boldsymbol{\psi}): \mathcal{D} \rightarrow \mathbb{R} \\
& F_{\mathrm{MF}}(\mathbf{b} ; p, \psi)=-\sum_{\substack{i, j, l, m \\
j \neq i \\
m \neq l}} b_{i l} b_{j m} \log p_{l m}^{\mu_{i j}}-\sum_{i, l} b_{i l} \log p_{l l}^{\mu_{i i}} \\
& \quad+\sum_{i, l}\left(b_{i l} \log b_{i l}+\left(1-b_{i l}\right) \log \left(1-b_{i l}\right)\right)+\log k \tag{2}
\end{align*}
$$

The mean-field PML (MFPML) estimate is defined as

$$
p_{\mathrm{MFPML}}^{(\psi)} \triangleq \arg \min _{p \in \mathcal{C}} \min _{\left(b_{j}\right) \in \mathcal{D}} F_{\mathrm{MF}}(\mathbf{b} ; p, \boldsymbol{\psi}) .
$$

[^10]
## Empirical results

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We have a low-complexity algorithm to compute MFPML estimate.


Figure: Histogram of estimation error of absolute second largest eigenvalue of transition matrix for $k=20$ and $n=400$.

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Figure: Histogram of estimation error of entropy rate for $k=20$ and $n=400$.

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Figure: Histogram of estimation error of absolute second largest eigenvalue of transition matrix for $k=50$ and $n=2000$.

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Figure: Histogram of estimation error of entropy rate for $k=50$ and $n=2000$.

## Points to ponder on

- Good reasons to study PML estimates for Markov chains.
- Obtaining efficient approximations is hard.
- Bethe approximation: Complexity blows up very quickly.
- Ideally want algorithms to work for large $k$.
- Even the mean field PML estimate becomes difficult to implement for very large $k$.


## Thank you!


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