

Lecture 13: Foster-Lyapunov Theorem

1 Foster's Theorem

Theorem 1.1 (Foster,1950). Let $\{X_n\}_{n \geq 0}$ be a irreducible DTMC on \mathbb{N}_0 if there exist a function $L : \mathbb{N}_0 \rightarrow \mathbb{R}_+$ with $\mathbb{E}[L(X_0)] < \infty$, such that for some $K > k \geq 0$, and $\epsilon > 0$:

1. $|\{x \in \mathbb{N}_0 : L(x) \leq k\}| < \infty$
2. $\mathbb{E}[L(X_n)|X_{n-1}] < K, \text{ when } L(X_{n-1}) \leq k.$
3. $\mathbb{E}[L(X_n) - L(X_{n-1})|X_{n-1}] < -\epsilon \text{ if } L(X_{n-1}) \geq k$

Then $\{X_n\}_{n \geq 0}$ is positive recurrent. ($L \equiv$ "potential function" or energy or lyapunov function).

Proof. As DTMC is irreducible than enough to show that some state is positive recurrent. By renewal theory, for ant DTMC, for all $x \in \mathbb{N}_0$,

$$\lim_{N \rightarrow \infty} \mathbb{E} \left[\sum_{n=1}^N \mathbf{1}_{\{X_n=x\}} \right] = \frac{1}{\mu_{xx}} \quad (1)$$

where

$$\mu_{xx} = \begin{cases} \infty & \text{if } x \text{ is transient} \\ \sum_{m \geq 0} m f_{xx}^m & \text{if } x \text{ is recurrent} \end{cases}$$

consider the RHS of equation (1)

$$\lim_{N \rightarrow \infty} \mathbb{E} \left[\sum_{n=1}^N \mathbf{1}_{\{X_n=x\}} \right] > 0 \iff x \text{ is positive recurrent}$$

consider

$$\begin{aligned}
0 &\leq \mathbb{E}[L(X_n)] = \mathbb{E}[L(X_0)] + \sum_{n=1}^N \mathbb{E}[L(X_n) - L(X_{n-1})] \\
&= \mathbb{E}[L(X_0)] + \sum_{n=1}^N \mathbb{E}[L(X_n) - L(X_{n-1}) \mathbf{1}_{\{L(X_{n-1}) > k\}}] + \sum_{n=1}^N \mathbb{E}[L(X_n) - L(X_{n-1}) \mathbf{1}_{\{L(X_{n-1}) \leq k\}}] \\
&\leq \mathbb{E}[L(X_0)] + \sum_{n=1}^N \mathbb{E}[-\epsilon \mathbf{1}_{\{L(X_{n-1}) > k\}}] + \sum_{n=1}^N \mathbb{E}[K \mathbf{1}_{\{L(X_{n-1}) \leq k\}}] \\
&\Rightarrow \left[\mathbb{E} \sum_{n=1}^N \mathbf{1}_{\{L(X_{n-1}) \leq k\}} \right] (K + \epsilon) \geq -\mathbb{E}[L(X_0)] + \epsilon N \\
&\Rightarrow \frac{1}{N} \left[\mathbb{E} \sum_{n=1}^N \mathbf{1}_{\{L(X_{n-1}) \leq k\}} \right] \geq -\frac{\mathbb{E}[L(X_0)]}{K + \epsilon} + \frac{\epsilon}{K + \epsilon} \\
&\Rightarrow \limsup_{N \rightarrow \infty} \frac{1}{N} \left[\mathbb{E} \sum_{n=1}^N \mathbf{1}_{\{L(X_{n-1}) \leq k\}} \right] \geq \frac{\epsilon}{K + \epsilon}
\end{aligned}$$

Let $F = \{x \in \mathbb{N}_0 : L(x) \leq k\}$, $|F| < \infty$. Now,

$$\begin{aligned}
&\limsup_{N \rightarrow \infty} \frac{1}{N} \left[\mathbb{E} \sum_{n=1}^N \mathbf{1}_{\{L(X_{n-1}) \leq k\}} \right] \\
&= \limsup_{N \rightarrow \infty} \frac{1}{N} \left[\mathbb{E} \sum_{n=1}^N \sum_{x \in F} \mathbf{1}_{\{X_{n-1} \leq k\}} \right] \leq \sum_{x \in F} \left[\limsup_{N \rightarrow \infty} \frac{1}{N} \mathbb{E} \sum_{n=1}^N \mathbf{1}_{\{X_{n-1} \leq k\}} \right] \\
&= \sum_{x \in F} \left[\limsup_{N \rightarrow \infty} \frac{1}{N} \mathbb{E} \sum_{n=1}^N \mathbf{1}_{\{X_{n-1} \leq k\}} \right] \geq \frac{\epsilon}{K + \epsilon} > 0
\end{aligned}$$

Therefore there exist some $x \in F$ such that ,

$$\limsup_{N \rightarrow \infty} \frac{1}{N} \mathbb{E} \left[\sum_{n=1}^N \mathbf{1}_{\{X_{n-1} \leq k\}} \right] \geq \frac{\epsilon}{(K + \epsilon)|F|} > 0$$

Therefore there exist some $x \in F$ such that ,

$$\lim_{N \rightarrow \infty} \frac{1}{N} \mathbb{E} \left[\sum_{n=1}^N \mathbf{1}_{\{X_{n-1} \leq k\}} \right] > 0$$

□

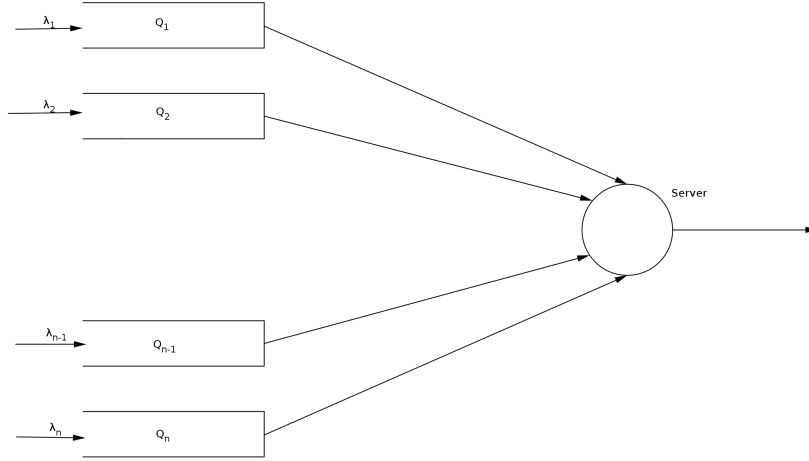


Figure 1: N queues single server, server chooses one queue and serve upto one packet

2 Applications of Foster's theorem: Queue scheduling/Max-weight scheduling

Consider N queue served by a single server in discrete time(Figure 1). At time slot $t = 1, 2, 3, \dots$ $A_i(t) \in \mathbb{N}_0$ packets arrive to each queue $i \in [N]$ independently.

1. $\mathbb{E}[A_i(t)] = \lambda_i$
2. $\mathbb{P}[A_i(t) = 0] > 0$
3. $\mathbb{E}[A_i(t)^2] \leq C$

Server picks one queue $Q(t) \in [N]$ for service. Let $R_i(t) = \mathbf{1}\{Q(t) = i\}$. One packet is served from $Q(t)$ if it is not empty. Let $X_i(t) =$ number of packets in queue i just before time slot t .

$$X_i(t+1) = (X_i(t) + A_i(t) - R_i(t))_+$$

Where

$$a_+ = \max(0, a)$$

aslo

$$X_i(t+1) = X_i(t) + A_i(t) - R_i(t) + L_i(t)$$

Where

$$L_i(t) = \begin{cases} 1 & \text{service attempted when } i \text{ is empty} \\ 0 & \text{otherwise} \end{cases}$$

Mean rate of arrivals to system : $= \sum_{i=1}^N \lambda_i$

Maximum rate of departure = 1, we will assume $\sum_{i=1}^N \lambda_i < 1$.

Theorem 2.1 (Max-weight scheduling algorithm, 1992).

$$Q(t) = \arg \max_{i \in N} X_i(t)$$

that is serve the longest queue. Under MAX-WT, $X(t) = (X_i(t))_{i=1}^N$ is a DTMC which is irreducible and aperiodic on state space \mathbb{N}_0^N . As long as $\sum_{i=1}^N \lambda_i < 1$, $\{X_n\}$ is positive recurrent.

Proof. By Foster's theorem, define the Lyapunov function:

$$L(x) = \frac{1}{2} \sum_{i=1}^N x_i^2$$

consider

$$\begin{aligned} & L(X(t)) - L(X(t-1)) \\ &= \frac{1}{2} \sum_{i=1}^N [(X_i(t))^2 - (X_i(t-1))^2] \\ &= \frac{1}{2} \sum_{i=1}^N [\{X_i(t-1) + A_i(t-1) - R_i(t-1) + L_i(t-1)\}^2 - (X_i(t-1))^2] \\ &\leq \frac{1}{2} \sum_{i=1}^N [\{X_i(t-1) + A_i(t-1) - R_i(t-1)\}^2 - (X_i(t-1))^2] \end{aligned}$$

Therefore

$$\begin{aligned} \mathbb{E}[L(X(t)) - L(X(t-1)) | X(t-1) = x] &\leq \frac{1}{2} \sum_{i=1}^N \mathbb{E}[(x_i + A_i(t-1) - R_i(t))^2 - x_i^2 | X(t-1) = x] \\ &= \frac{1}{2} \sum_{i=1}^N \left[2x_i \mathbb{E}[A_i(t-1) - R_i(t) | X(t-1) = x] + \mathbb{E}[(A_i(t-1) - R_i(t-1))^2 | X(t-1) = x] \right] \\ &= \sum_{i=1}^N x_i \lambda_i - \sum_{i=1}^N x_i R_i(t-1) + \frac{N}{2}(1+C) \\ &= \sum_{i=1}^N x_i \lambda_i + \frac{N}{2}(1+C) - \max_i x_i \\ &\leq \frac{N}{2}(1+C) + (\max_i x_i) \left(\sum_{i=1}^N \lambda_i - 1 \right) \\ &= C_1 - \epsilon(\max_i x_i) \end{aligned}$$

Foster's theorem applies with $k = \max\{\frac{\|x\|_2^2}{2}\}$

□