## Lecture-22: DTMC: Irreducibility and Aperiodicity

## 1 Communicating classes

**Definition 1.1.** For states  $x, y \in \mathcal{X}$ , it is said that state y is **accessible** from state x if  $p_{xy}^{(n)} > 0$  for some  $n \in \mathbb{Z}_+$ , and denoted by  $x \to y$ . If two states  $x, y \in \mathcal{X}$  are accessible to each other, they are said to **communicate** with each other, denoted by  $x \leftrightarrow y$ .

- (i) A set of states that communicate are called a communicating class.
- (ii) A communicating class C is called **closed** if no edges leave this class. That is, for all edges  $(x, y) \in C \times C^c$ , we have  $p_{xy} = 0$ .
- (iii) An **open** communicating class is not closed, i.e. there exist an edge that leaves this class. That is, there exists an edge  $(x, y) \in C \times C^c$  such that  $p_{xy} > 0$ .

**Proposition 1.2.** Communication is an equivalence relation.

*Proof.* Relation on state space  $\mathcal{X}$  is a subset of product of sets  $\mathcal{X} \times \mathcal{X}$ . Communication is a relation on state space  $\mathcal{X}$ , as it relates two states  $x, y \in \mathcal{X}$ . To show equivalence, we have to show reflexivity, symmetry, and transitivity of the relation.

- Symmetry: Further, if  $x \leftrightarrow y$ , then we know that  $x \rightarrow y$  and  $y \rightarrow x$  and hence  $y \leftrightarrow x$ . Hence, the symmetry of the relation follows.
- Transitivity: For transitivity, suppose  $x \leftrightarrow y$  and  $y \leftrightarrow z$ . Let  $m, n \in \mathbb{Z}_+$  such that  $p_{xy}^{(m)} > 0$  and  $p_{yz}^{(n)} > 0$ . Then by Chapman Kolmogorov equation, we have

$$p_{xz}^{(m+n)} = \sum_{w \in \mathcal{X}} p_{xw}^{(m)} p_{wz}^{(n)} \ge p_{xy}^{(m)} p_{yz}^{(n)} > 0.$$

This implies  $x \rightarrow z$ , and using similar arguments one can show that  $z \rightarrow x$ , and the transitivity follows.

Reflexivity: If this relation has a single element, then it is obvious. If not, then for  $x \leftrightarrow y$ , we have Since  $p_{xy}^{(n)} > 0$  and  $p_{yx}^{(m)} > 0$  for some  $m, n \in \mathbb{Z}_+$ . Therefore,  $p_{xx}^{(n+m)} \ge p_{xy}^{(n)} p_{yx}^{(m)} > 0$  and hence we have  $x \leftrightarrow x$ , implying the reflexivity of the relation.

*Remark* 1. Hence the communication relation partitions state space  $\mathcal{X}$  into equivalence classes.

**Definition 1.3.** Each equivalence class is called a **communicating class**. A property of states is said to be a **class property** if for each communicating class C, either all states in C have the property, or none do.

## **1.1** Irreducibility and periodicity

**Definition 1.4.** A Markov chain with a single class is called an **irreducible** Markov chain. That is, for any two states  $x, y \in \mathcal{X}$ , there exists an integer  $n \in \mathbb{N}$  such that  $p_{xy}^{(n)} > 0$ . In other words, any state y can be reached from any state x using transitions of positive probability.

**Definition 1.5.** Let  $\mathcal{T}(x) \triangleq \left\{ n \in \mathbb{N} : p_{xx}^{(n)} > 0 \right\}$  be the set of times when the chain can possibly return to the initial state *x*. The **period** of any state  $x \in \mathcal{X}$  is defined as

$$d(x) \triangleq \gcd \mathcal{T}(x) = \gcd \{ n \in \mathbb{N} : p_{xx}^{(n)} > 0 \}.$$

We define  $d(x) = \infty$ , if  $p_{xx}^{(n)} = 0$  for all  $n \in \mathbb{N}$ . A state  $x \in \mathcal{X}$  is called **aperiodic** if the period d(x) is 1.

**Proposition 1.6.** If  $x \leftrightarrow y$ , then d(x) = d(y). That is, periodicity is a class property.

*Proof.* Let  $m, n \in \mathbb{N}$  be such that  $p_{xy}^{(m)} p_{yx}^{(n)} > 0$ . Suppose  $s \in \mathcal{T}(x)$ , that is  $p_{xx}^{(s)} > 0$ . Then

$$p_{yy}^{(n+m)} \ge p_{yx}^{(n)} p_{xy}^{(m)} > 0,$$
  $p_{yy}^{(n+s+m)} \ge p_{yx}^{(n)} p_{xx}^{(s)} p_{xy}^{(m)} > 0.$ 

Hence d(y)|n + m and d(y)|n + s + m, and hence d(y)|s for any  $s \in \mathcal{T}(x)$ . In particular, it implies that d(y)|d(x). By symmetrical arguments, we get d(x)|d(y). Hence d(x) = d(y).

**Definition 1.7.** For an irreducible chain, the period of the chain is defined to be the period which is common to all states. An irreducible Markov chain is called **aperiodic** if the single communicating class is aperiodic.

**Proposition 1.8.** *If the transition matrix P is aperiodic and irreducible, then there is an integer*  $r_0$  *such that*  $p_{xy}^{(r)} > 0$  *for all*  $x, y \in X$  *and*  $r \ge r_0$ .

## 1.2 Transient and recurrent states

**Proposition 1.9.** *Transience and recurrence are class properties.* 

*Proof.* Let us start with proving recurrence is a class property. Let *x* be a recurrent state and let  $x \leftrightarrow y$ . Hence there exist some m, n > 0, such that  $P_{xy}^{(m)} > 0$  and  $p_{yx}^{(n)} > 0$ . As a consequence of the recurrence,  $\sum_{s \in \mathbb{N}} p_{xx}^{(s)} = \infty$ . It follows that *y* is recurrent by observing

$$\sum_{s \in \mathbb{N}} p_{yy}^{(m+n+s)} \geqslant \sum_{s \in \mathbb{N}} p_{yx}^{(n)} p_{xx}^{(s)} P_{xy}^{(m)} = \infty.$$

Now, if *x* were transient instead, we conclude that *y* is also transient by the following observation

$$\sum_{s \in \mathbb{N}} p_{yy}^{(s)} \leqslant \frac{\sum_{s \in \mathbb{N}} p_{xx}^{(m+n+s)}}{p_{yx}^{(n)} P_{xy}^{(m)}} < \infty.$$

**Corollary 1.10.** If y is recurrent, then for any state x such that  $y \to x$ , then  $x \to y$  and  $f_{xy} = 1$ .

*Proof.* Let  $y \in \mathcal{X}$  be a recurrent state, and consider state  $x \in \mathcal{X}$  such that  $y \to x$ . We will show that  $f_{xy} = 1$  and hence  $f_{xy}^{(n)} > 0$  for some  $n \in \mathbb{Z}_+$  and  $x \to y$ . To this end, we observe that since  $y \to x$ , there exists an integer  $n \in \mathbb{Z}_+$  such that the probability of hitting state x for the first time starting from state y in n-steps is positive. That is,

$$f_{yx}^{(n)} \triangleq P_y \{X_n = x, X_{n-1} \neq x, \dots, X_1 \neq x\} = P_y \{H_x = n\} > 0.$$

Suppose  $f_{xy} < 1$ , then from the strong Markov property, we have

$$1 - f_{yy} = P_y \{ H_y = \infty \} \ge P_y \{ H_y = \infty, H_x = n \} = P_x \{ H_y = \infty \} P_y \{ H_x = n \} = f_{yx}^{(n)} (1 - f_{xy}) > 0.$$

This is a contradiction since state *y* is recurrent. This implies that  $f_{xy} = 1$  and hence  $x \to y$ .

**Corollary 1.11.** Let  $x, y \in \mathcal{X}$  be in the same communicating class and the state y is recurrent. Then,  $\lim_{n \in \mathbb{N}} \frac{\sum_{k=1}^{n} p_{xy}^{(k)}}{n} = \frac{1}{\mu_{yy}}$ . Furthermore, if the state y is aperiodic, then  $\lim_{n \in \mathbb{N}} p_{xy}^{(n)} = \frac{1}{\mu_{yy}}$ .

*Proof.* Since *y* is recurrent and  $y \to x$ , it follow that  $f_{xy} = 1$  from the previous Lemma. From the Theorem 1.7 in previous lecture, it follows that  $\lim_{n \in \mathbb{N}} \frac{\sum_{k=1}^{n} p_{xy}^{(k)}}{n} = \frac{1}{\mu_{yy}}$ . Let the period of the state *y* be *d*. Then we know that there exists a positive integer  $r_0$  such that for all

Let the period of the state *y* be *d*. Then we know that there exists a positive integer  $r_0$  such that for all  $n \ge r_0$ , we have  $p_{yy}^{(nd)} > 0$ .

**Theorem 1.12.** The states in a communicating class are of one of the following types; all transient, or all null recurrent, or all positive recurrent.

*Proof.* It suffices to show that if x, y belong to the same communicating class and y is null recurrent, then x is null recurrent as well. We take  $r, s \in \mathbb{N}$ , such that  $p_{yx}^{(r)} p_{xy}^{(s)} > 0$ . It follows that  $p_{yy}^{r+\ell+s} \ge p_{yx}^{(r)} p_{xx}^{(\ell)} P_{xy}^{(s)}$  for all  $\ell \in \mathbb{N}$ . Hence, for any n > r + s, we have

$$\frac{1}{n}\sum_{k=1}^{n}p_{yy}^{(k)} \ge \frac{1}{n}\sum_{k=r+s+1}^{n}p_{yy}^{(k)} \ge \left(\frac{n-r-s}{n}\right)\left(\frac{1}{n-r-s}\sum_{\ell=1}^{n-r-s}p_{xx}^{(\ell)}\right)p_{yx}^{(r)}P_{xy}^{(s)}.$$

Since *y* is null recurrent LHS goes to zero as *n* increases, which implies  $\lim_{n \in \mathbb{N}} \frac{1}{n} \sum_{\ell=1}^{n} p_{xx}^{(\ell)} = 0$ . Hence, *x* is null recurrent as well.