

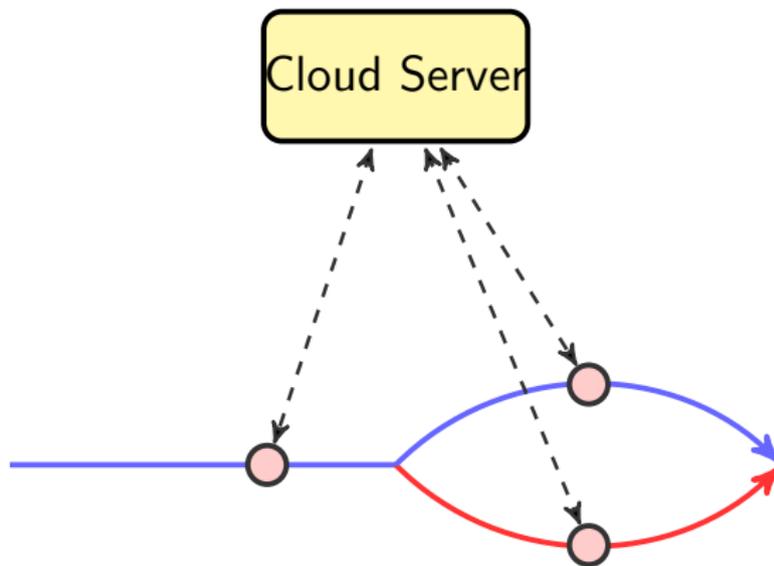
# Real-Time Status Updates for Markov Sources over Unreliable Channels

Parimal Parag

Joint work with Sanidhay Bhambay and Sudheer Poojary

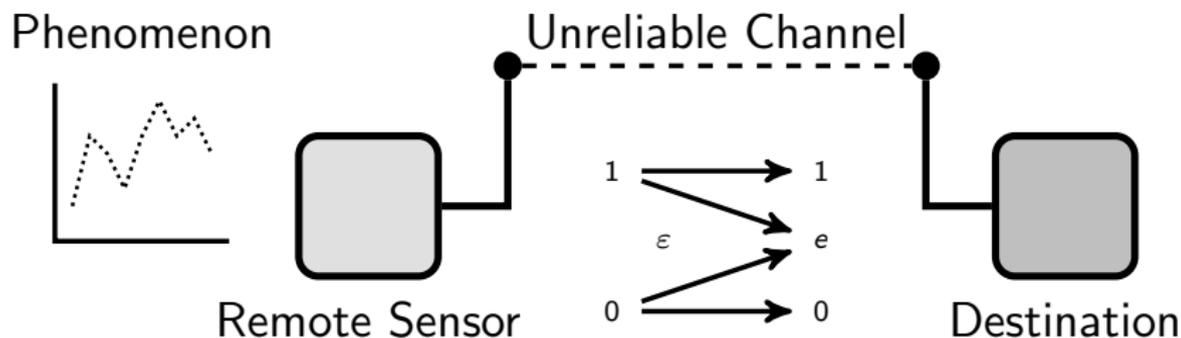


## Real-time decision systems



- ▶ Cyber-physical systems: Environmental/health monitoring
- ▶ Internet of Things: Real-time actuation/control
- ▶ Critical to know the status update before decision making

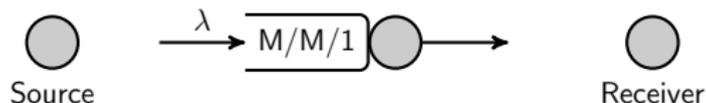
## Status updates over unreliable channels



### Question

How to encode message at the temporally correlated source for **timely** update over an iid binary symmetric erasure channel?

## Quantifying timeliness: delay versus age



### Information packet delay at transmitter

- ▶ If the arrival rate is higher than the service rate, then unstable queue and infinite delay

### Information age at receiver

- ▶ Generation time of last received information packet  $U(t)$
- ▶ Age of information  $A(t) = t - U(t)$
- ▶ Finite average age for always on source

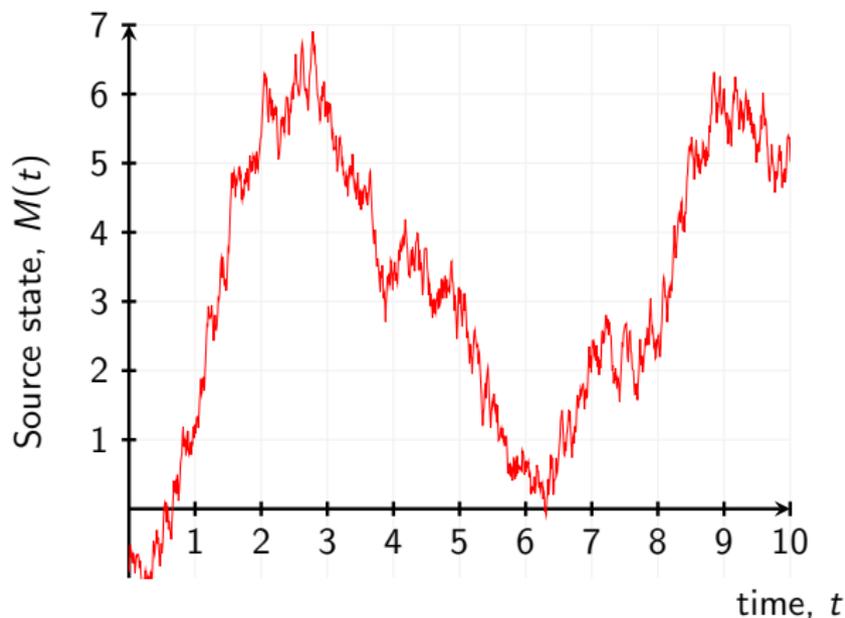
## Prior work

- ▶ 'Age of information' introduced by Kaul and R. Yates in 2012
- ▶ Different queueing models and queueing disciplines are used to analyze update system [A. Ephremides, N. B. Shroff, Y. Sun, E. Modiano, M. Costa]
- ▶ M/M/1 queue with FCFS/LCFS [Kaul, Gruster, S. Kompella, R. Yates]

## Summary

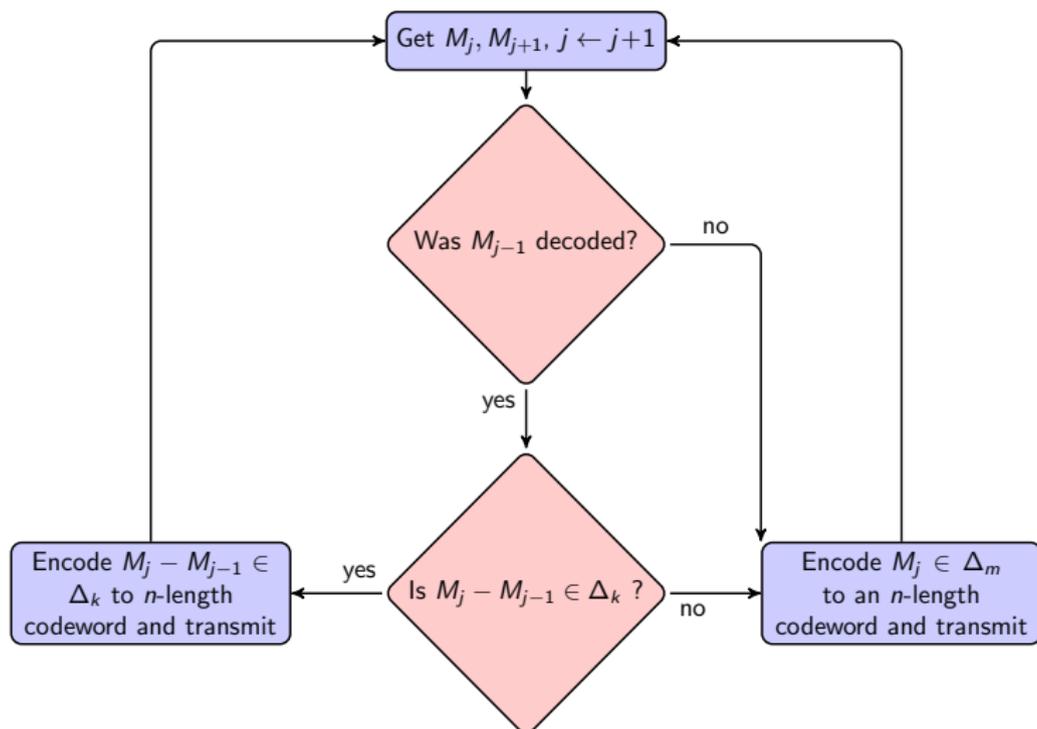
- ▶ Limiting average age minimizing update rate
- ▶ Channel uncertainty modeled by random service time
- ▶ A scheduled update is delivered successfully

## Source Model



- ▶ Discrete sampled source  $M_j = M(jn) \in \Delta_m$  Markov with transition matrix  $P$  and invariant distribution  $\nu$

# Update Protocol



- Special case when  $\Delta_k = 0$ , always send actual update

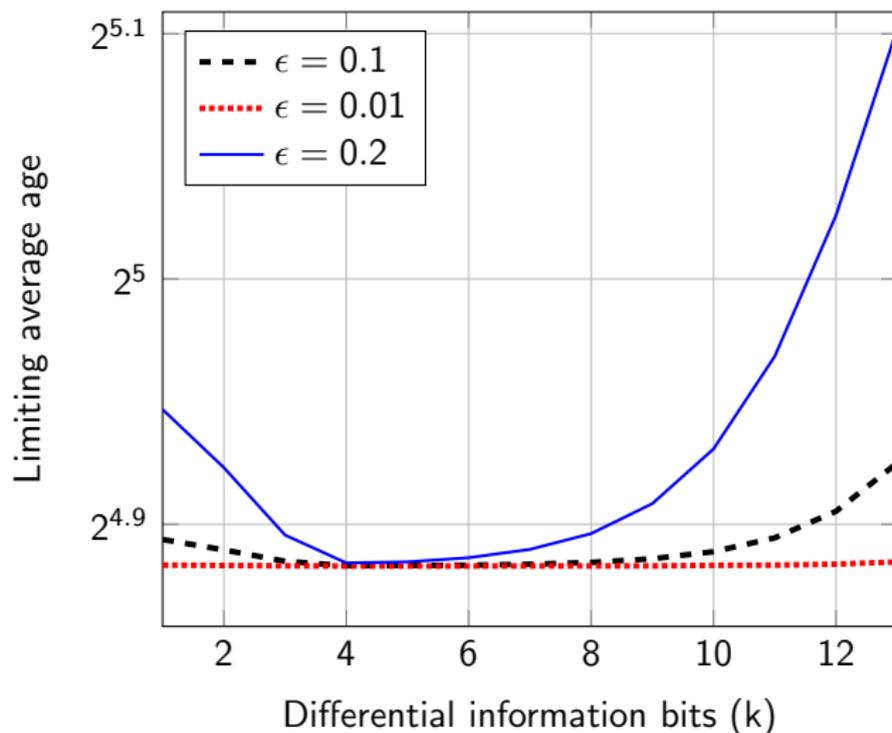
# Problem Statement

## Question

For a fixed code-length  $n$ , find the code-rate  $\frac{k}{n}$  for the given Markov source  $(M_j \in \Delta_m : j \in \mathbb{N})$  over iid binary symmetric erasure channel  $\sim \text{Bernoulli}(\epsilon)$  that minimizes the limiting average age

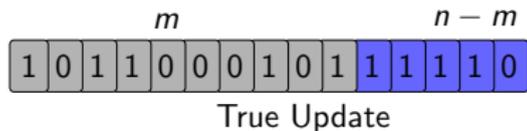
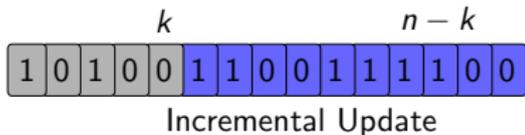
$$\lim_{T \rightarrow \infty} \frac{1}{T} \sum_{t=1}^T A(t).$$

## Optimal code-rate



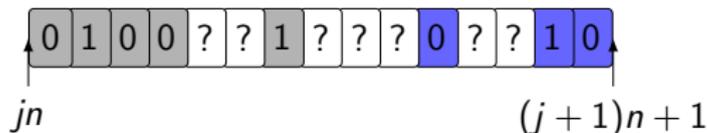
- ▶ Code length  $n = 20$ , number of information bits  $m = 15$

## Permutation invariant code



## Bit-wise erasure channel

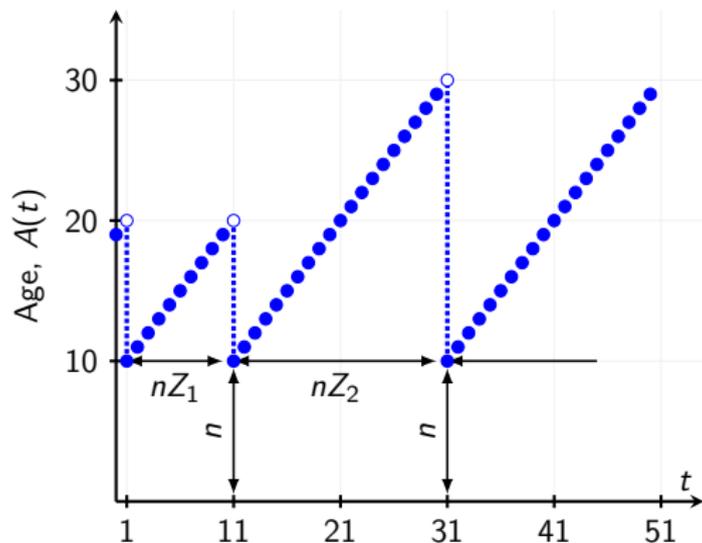
- ▶ Number of erasures per codeword  $E \sim \text{Binomial}(n, \epsilon)$



## Probability of decoding failure

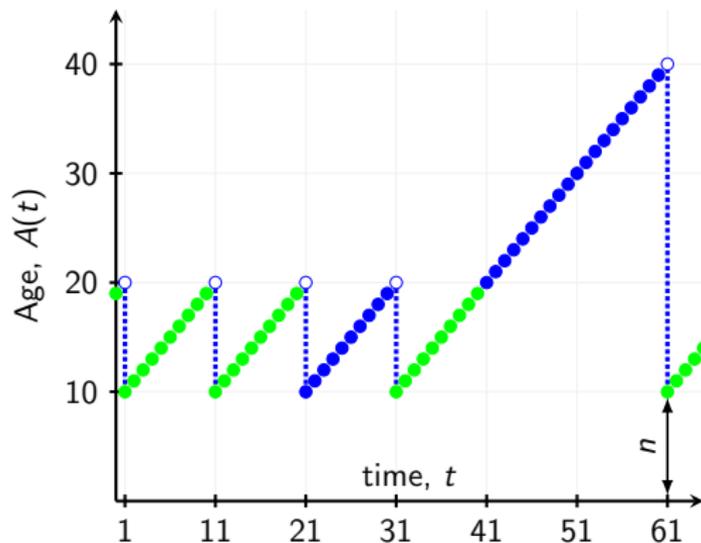
- ▶ Actual state updates:  $p_a = \mathbb{E}P(n, n - m, E)$
- ▶ Incremental updates:  $p_d = \mathbb{E}P(n, n - k, E)$

# Information age process



- ▶ Age determined by scaled samples ( $A_j = \frac{1}{n}A(jn + 1) : j \in \mathbb{N}$ )
- ▶ Mean average age is  $\frac{\mathbb{E} \sum_{t=1}^{T_1} A(t)}{\mathbb{E} T_1} = n\mathbb{E}A_j + \frac{(n-1)}{2}$

# Age process for incremental updates



- ▶ Decoding success iff  $A_j = 1$ , Bernoulli indicators with  $p_a, p_d$
- ▶ Incremental update if  $\{A_{j-1} = 1\} \cap \{M_j - M_{j-1} \in \Delta_k\}$
- ▶ Process  $((M_j, A_j) : j \in \mathbb{N})$  is jointly Markov

## Highly Correlated Source

- ▶  $M_j - M_{j-1} \in \Delta_k$  almost surely, then  
 $P_i(\Delta_k) = \sum_{l-i \in +\Delta_k} P_{il} = 1$  or  $\tilde{D} = (p_a - p_d)P$

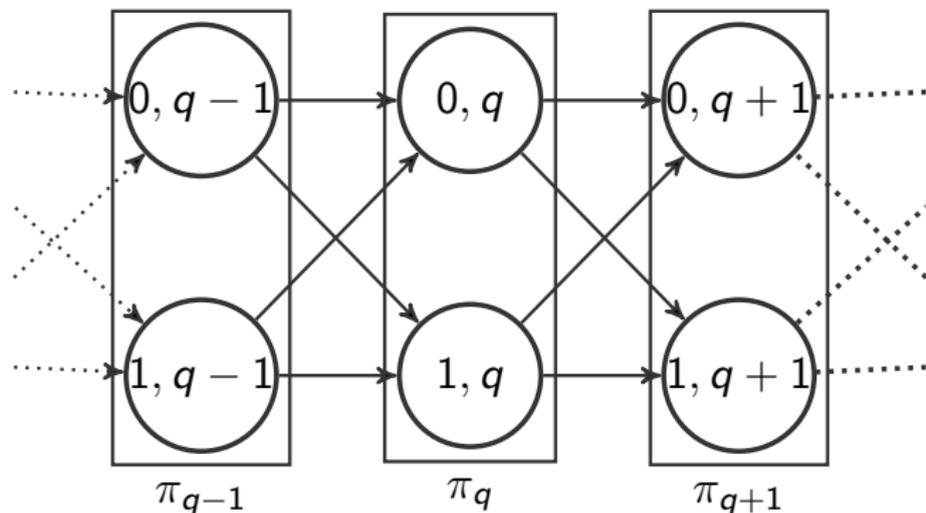
$$\begin{bmatrix} P_{00} & P_{01} & 0 & 0 & \dots & \dots \\ P_{10} & P_{11} & P_{12} & 0 & \dots & \dots \\ 0 & P_{21} & P_{22} & P_{23} & \dots & \dots \\ \vdots & \vdots & \ddots & \ddots & \ddots & \vdots \end{bmatrix}$$

## Uniform Markov Source

- ▶  $P_i(\Delta_k)$  is identical to all states  $i$

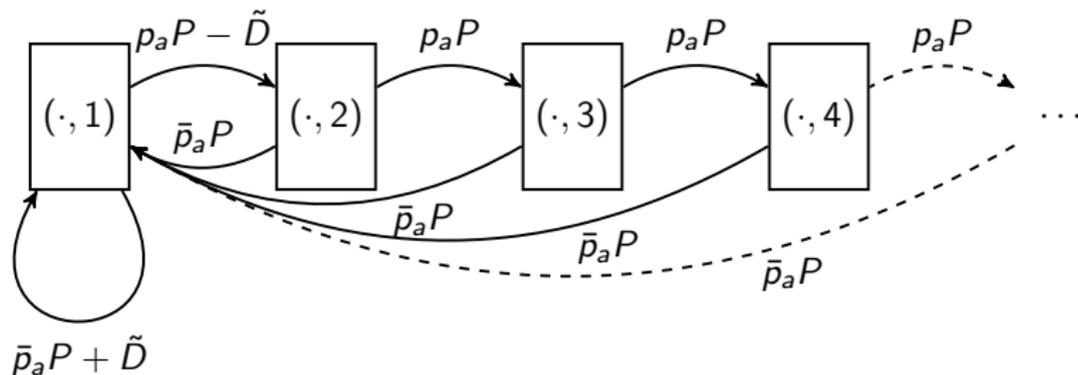
$$\begin{bmatrix} P_{00} & P_{01} & P_{02} & P_{03} & \dots & \dots \\ P_{10} & P_{11} & P_{12} & P_{13} & \dots & \dots \\ P_{20} & P_{21} & P_{22} & P_{23} & \dots & \dots \\ \vdots & \vdots & \ddots & \ddots & \ddots & \vdots \end{bmatrix}$$

## Quasi birth-death process



- ▶ Transition probability  $(i, q) \rightarrow (j, q+1)$ :  $p_a P_{ij}$
- ▶ Transition probability  $(i, 1) \rightarrow (i, 2)$ :  
 $-(p_a - p_d)P_{ij}1_{\{j-i \in \Delta_k\}} + p_a P_{ij}$

## Block transitions



- ▶  $\tilde{D}_{ij} = (p_a - p_d)P_{ij}1_{\{j-i \in \Delta_k\}}$
- ▶ Transition matrix for the joint Markov process

$$\begin{bmatrix} p_a P + \tilde{D} & p_a P - \tilde{D} & 0 & 0 & \dots & \dots \\ \bar{p}_a P & 0 & p_a P & 0 & \dots & \dots \\ \bar{p}_a P & 0 & 0 & p_a P & \dots & \dots \\ \bar{p}_a P & 0 & 0 & 0 & p_a P & \dots \\ \vdots & \vdots & \ddots & \ddots & \ddots & \vdots \end{bmatrix}$$

## Equilibrium distribution

Geometric form in sampled-scaled age

$$\pi_q = \bar{p}_a \nu (I - \tilde{D})^{-1} (p_a P - \tilde{D}) (p_a P)^{q-2}, \quad q \geq 2.$$

Mean sampled and scaled age

$$\mathbb{E}A = \langle -\nu (I - \tilde{D})^{-1}, \mathbf{1} \rangle + 1 + \frac{1}{\bar{p}_a}.$$

Tail decay-rate

$$\theta = - \lim_{q \rightarrow \infty} \frac{1}{q} \log \sum_{u=q}^{\infty} \langle \pi_u, \mathbf{1} \rangle = \log \frac{1}{p_a}.$$

# Discussion and Concluding Remarks

## Main Contributions

- ▶ Integration of coding and block Markov chain techniques to study timely communication for delay-sensitive traffic
- ▶ We model channel unreliability by the erasure channel
- ▶ Model source correlation by Markov process
- ▶ True and incremental updates are special cases