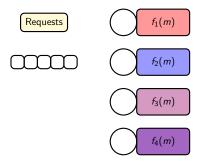
Parallel Coded Access

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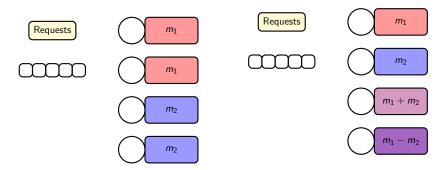
Problem Statement



Compute mean access time to download single message m

- with number of fragments k such that $m = (m_1, \ldots, m_k)$
- with encoding $(f_1(m), \ldots, f_n(m))$, and $f_i(m)$ stored at node i

Symmetric Codes



Replication (n, k)

Piece *i* stored at n/k servers

MDS (n, k)

Whole message can be decoded by any k out of n servers

Applications

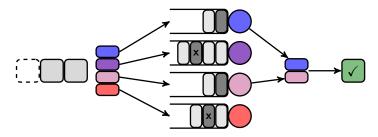
Distributed Storage

- Content streaming: NetFlix, HotStar, Eros Now, YouTube, Hulu, Amazon Prime Video
- Cloud storage: GitHub, DropBox, iCloud, OneDrive, UbuntuOne
- Cloud service: Facebook, Google Suite, Office365

Distributed Computation

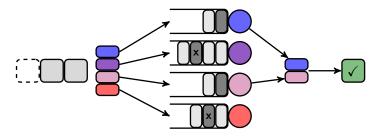
- Cloud computing: Amazon Web Services, Microsoft Azure, Google Search
- **Cluster computing:** Hadoop, Spark
- Distributed database: Aerospike, Cassandra, Couchbase, Druid

Information Retrieval



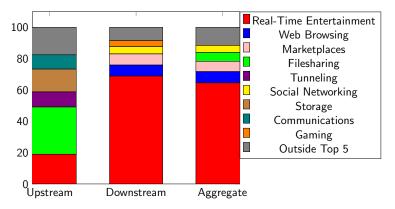
- Query *m* from sub-queries $(f_1(m), \ldots, f_n(m))$
- k sub-queries suffice
- Parallel processing for scaling and speed-up
- Redundancy for availability and speed-up

Distributed Computation



- Result *m* from sub-results $(f_1(m), \ldots, f_n(m))$
- k sub-results suffice
- Parallel processing for scaling and speed-up
- Redundancy for availability and speed-up

Dominant traffic on Internet



Peak Period Traffic Composition (North America)

 Real-Time Entertainment: 64.54% for downstream and 36.56 % for mobile access¹

 $[\]label{eq:linear} \begin{array}{c} 1 \\ https://www.sandvine.com/downloads/general/global-internet-phenomena/2015/global-internet-phenomena-report-latin-america-and-north-america.pdf \end{array}$

System Model

File storage

- Each media file divided into k pieces
- Pieces encoded and stored on n servers

Arrival of requests

- Each request wants entire media file
- Poisson arrival of requests with rate λ

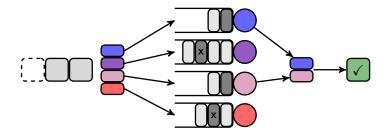
Time in the system

Till the reception of whole file

Service at each server

• IID exponential service time with rate $\mu = k/n$

Storage Coding -(n, k) Fork-Join Model



exempli gratia: Joshi, Liu, Soljanin (2012, 2014), Joshi, Soljanin, Wornell (2015), Sun, Zheng, Koksal, Kim, Shroff (2015), Kadhe, Soljanin, Sprintson (2016), Li, Ramamoorthy, Srikant (2016)

Prior Work and Contributions

Kannan et al: join k queues for replication and MDS codes

- Numerical bounds using block Markov chains
- Trade-off between numerical accuracy and computational effort

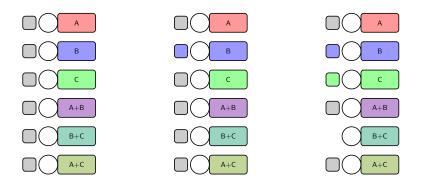
Soljanin, Wornell et al: fork-join (n, k) queues for MDS codes

- Closed-form upper and lower bounds
- Loose bounds for most of the rate region

This work: fork-join (n, k) queues for all symmetric codes

- Tight closed-form approximations for all rate regions
- Stability region for all symmetric codes
- Delay minimising symmetric code

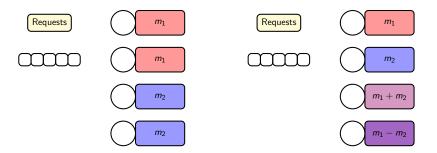
Coding Model



- ▶ Information sets $\mathcal{I} = \{S \subset [n] : |S| = k, f_S \text{ reconstructs } m\}$
- ▶ Observed servers $T \subset S$ for some info set $S \in \mathcal{I}$
- Useful servers $M(T) = \bigcup_{S \in \mathcal{I}} S \setminus T$

Symmetric codes: number useful servers $N_{|T|} = |M(T)|$

Symmetric Codes



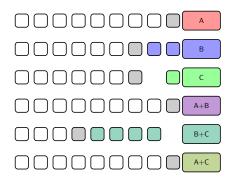
Replication (n, k)

Number of useful servers $N_i = (k - i)n/k$

MDS (n, k)

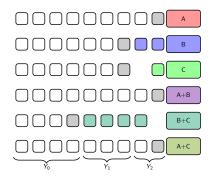
Number of useful servers $N_i = (n - i)$

State Space Collapse



- ▶ $L(t) = \{(\ell_1, \dots, \ell_r) : \ell_i = |T_i|, \ell_1 \ge \ell_2\}$ is a Markov process
- Arrival: $(\ell_1, \ldots, \ell_r) \rightarrow (\ell_1, \ldots, \ell_r, 0)$ at rate λ
- Departure: $(\ell_1, \ldots, \ell_r) \rightarrow (\ell_2, \ldots, \ell_r)$ at rate $N_{\ell_1}\mu$
- ▶ Service: $(\ldots, \ell_i, \ldots) \rightarrow (\ldots, \ell_i + 1, \ldots)$ at rate $(N_{\ell_i} N_{\ell_{i+1}})\mu$

State Space Transformation

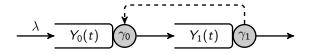


•
$$\mathbf{Y}(t) = \{Y_0, Y_1, \dots, Y_{k-1}\}$$
 is a Markov process

• Arrival: $Y_0 \rightarrow Y_0 + 1$ at rate λ

- Departure: $Y_{k-1} \rightarrow Y_{k-1} 1$ at rate $N_{k-1}\mu$
- ▶ Service: $(Y_{i-1}, Y_i) \rightarrow (Y_{i-1} 1, Y_i + 1)$ at rate $(N_{i-1} N_{l_{i-1}})\mu$

Tandem Queue Interpretation (General Case)



Tandem Queue with Pooled Resources

- Servers with empty buffers help upstream
- Aggregate service at level i becomes

$$\sum_{j=i}^{l_i(t)-1} \gamma_j$$
 where $l_i(t) = k \wedge \{l > i : Y_l(t) > 0\}$

 No explicit description of stationary distribution for multi-dimensional Markov process

Stability Region For Pooled Tandem Queues

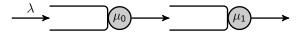
$$\xrightarrow{3\lambda} Y_0(t) [\Gamma_0 \xrightarrow{2\lambda} Y_1(t) [\Gamma_1 \xrightarrow{\lambda} Y_1(t)]\Gamma_2 \longrightarrow$$

For a distributed storage system with symmetric codes and fork-join queues with FCFS service, the stability region is

$$\lambda < \min\left\{\frac{\Gamma_i}{k-i}: i \in \{0,\ldots,k-1\}\right\},$$

where $\Gamma_i \triangleq \sum_{j=i}^{k-1} \gamma_j$ is the useful service rate for level *i*.

Bounding and Separating



Theorem[†]

When $\lambda < \min \mu_i$, tandem queue has product form distribution

$$\pi(y) = \prod_{i=0}^{k-1} rac{\lambda}{\mu_i} \left(1 - rac{\lambda}{\mu_i}
ight)^{y_i}$$

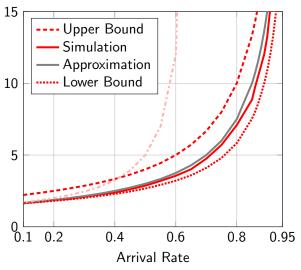
Uniform Bounds on Service Rate Transition rates are uniformly bounded by

$$\gamma_i \leq \sum_{j=i}^{l_i(y)-1} \gamma_j \leq \sum_{j=i}^{k-1} \gamma_j \triangleq \Gamma_i$$

[†]F. P. Kelly, Reversibility and Stochastic Networks. New York, NY, USA: Cambridge University Press, 2011.

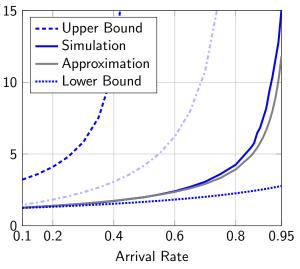
Mean Sojourn Time

Replication Coding

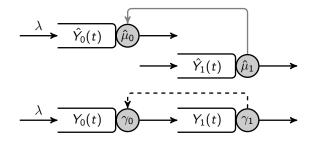


Mean Sojourn Time





Approximating Pooled Tandem Queue



Independence Approximation with Statistical Averaging Service rate is equal to base service rate γ_i plus cascade effect, averaged over time

$$\hat{\mu}_{k-1} = \gamma_{k-1} \qquad \hat{\pi}(y) = \prod_{i=0}^{k-1} \frac{\lambda}{\hat{\mu}_i} \left(1 - \frac{\lambda}{\hat{\mu}_i}\right)^{y_i}$$

Delay Minimizing Storage Code

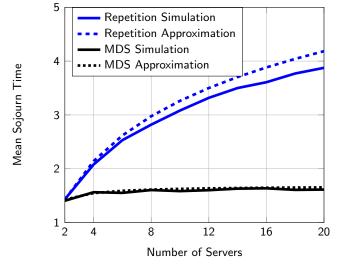
$$\xrightarrow{\lambda} \hat{Y}_0(t) \hat{\mu}_0 \longrightarrow \hat{Y}_1(t) \hat{\mu}_1 \longrightarrow$$

Optimizer to the objective function

$$\gamma^* = \arg \min \left\{ \sum_{i=1}^{k-1} \frac{1}{\Gamma_i - (k-i)\lambda} : \gamma \in \mathcal{A} \right\}.$$

The MDS coding scheme minimizes the approximate mean sojourn time for a fork-join queueing system with identical exponential servers among all symmetric codes.

Comparing Replication versus MDS Coding



Arrival rate 0.3 units and coding rate n/k = 2

Summary and Discussion

Main Contributions

- Analytical framework for study of distributed computation and storage systems
- Upper and lower bounds to analyze replication and MDS codes
- A tight closed-form approximation to study distributed storage codes
- MDS codes are better suited for large distributed systems
- Mean access time is better for MDS codes for all code-rates