

E2 205 Error-Control Coding

Lecture 17

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1 Performance of Convolutional Codes

In the trellis diagram, the metric of a path is the sum of path metrics in the case of a BSC.

This also holds in the case of the AWGN.

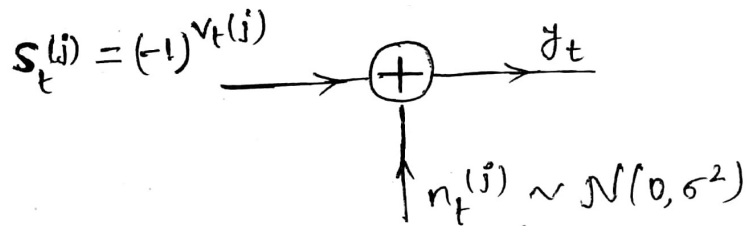


Fig. AWGN Channel

$$P(\vec{y}|\vec{c}_i) = \prod_{t=0}^{N-1} \prod_{j=1}^2 \frac{1}{\sqrt{2\pi\sigma^2}} \exp\left(\frac{-1}{2\sigma^2} (y_t(j) - s_t(j))^2\right).$$

So, the MLD reduces to a minimum Euclidean distance decoder.

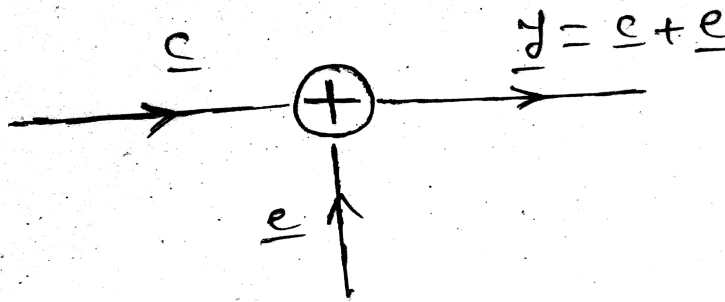
$$\min \sum_{t=0}^{N-1} \sum_{j=1}^2 (y_t(j) - s_t(j))^2$$

Thus it suffices to maximize

$$\sum_{t=0}^{N-1} \sum_{j=1}^2 y_t(j) s_t(j) \text{ (the inner product)}$$

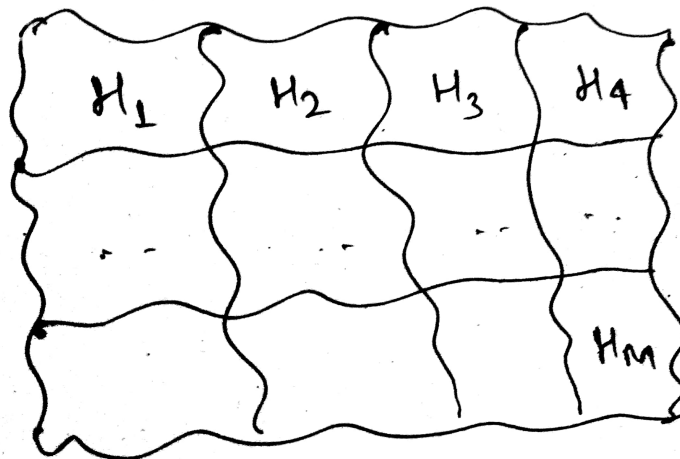
The rate of the procedure remains as before with min. replaced by max.
 Termination is not required in practice because the survivors at time $t + \tau$,
 tend to have a common prefix, up until time t , allowing code symbols upto
 time to be decoded by time $t + \tau$.

Consider the BSC in the form:



where $\vec{c} \in \mathcal{C}$ an $[n,k]$ linear code.

At the receiver let the decision regions be $H_i, 1 \leq i \leq M=2^k$.



$\vec{y} \in H_i$ requires that $d_H(\vec{y}, \vec{c}_i) \leq d_H(\vec{y}, \vec{c}_j)$.

Let us choose H_i s.t.

$\vec{y} \in H_i$, iff $d_H(\vec{y}, \vec{c}_i) \leq d_H(\vec{y}, \vec{c}_j)$.

Standard-Array Table

Codewords →	0000	0101	1010	1111
	0001	0100	1011	1110
	0010	0111	1000	1101
	0011	0110	1001	1100
↙ coset leaders				

Let $E = ([0000]^t, [0001]^t, [0010]^t, [0011]^t)$ be the set of coset leaders.

$H_0 = E(i^{th} \text{ column of standard array})$

$H_{0101} = \text{column containing } 0101 = H_0 + [0101]^t$.

$$\vec{H}_i = \vec{H}_0 + \vec{C}_i$$

Suppose $\vec{0}$ was transmitted and \vec{e} was the error pattern introduced by the BSC. Suppose the decoded codeword is \vec{C}_i . Then the residual error often decoding equals?

Next, suppose \vec{c}_j was transmitted and \vec{e} was again the error vector.

Then, $\vec{y} = \vec{c}_j + \vec{e}$

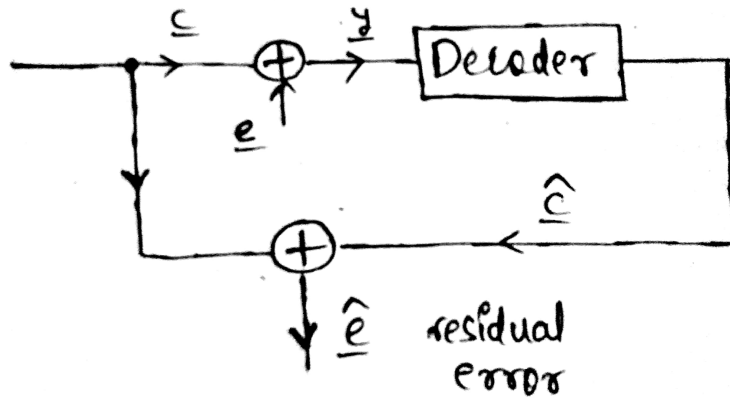
We know that, $\vec{e} \in H_i$.

So, $\vec{e} = \vec{c}_i + \vec{a}$, $\vec{a} \in E$

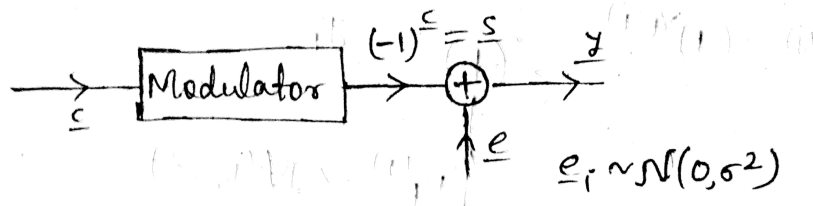
$\therefore \vec{y} = \vec{c}_j + \vec{c}_i + \vec{a}$

$\implies \vec{y} \in H\vec{c}_i + \vec{c}_j$

\therefore the residual error pattern =
 $(\vec{c}_i + \vec{c}_j) + \vec{c}_j = \vec{c}_i$
 (independent of the transmitted codeword)
 $\underline{c} + \hat{\underline{c}} = \underline{u}^T G + \hat{\underline{u}} G = (\underline{u}^T + \hat{\underline{u}}) G = \underline{\theta}$ (residual error)



Next consider the case of the AWGN channel

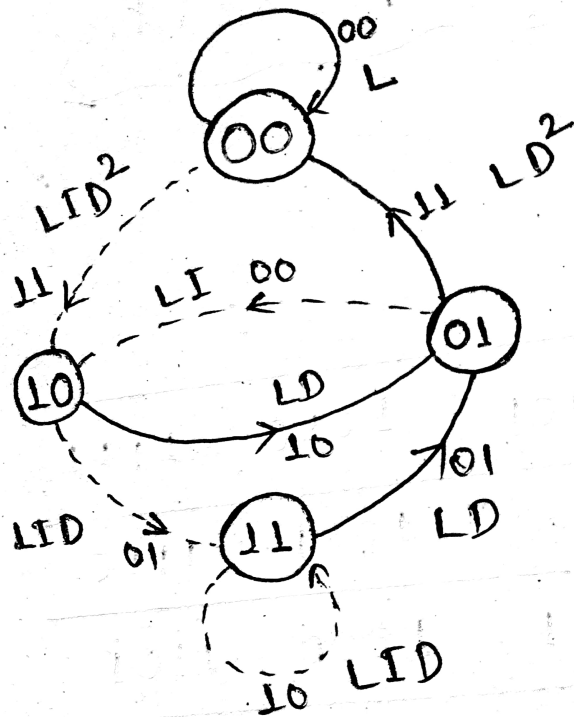
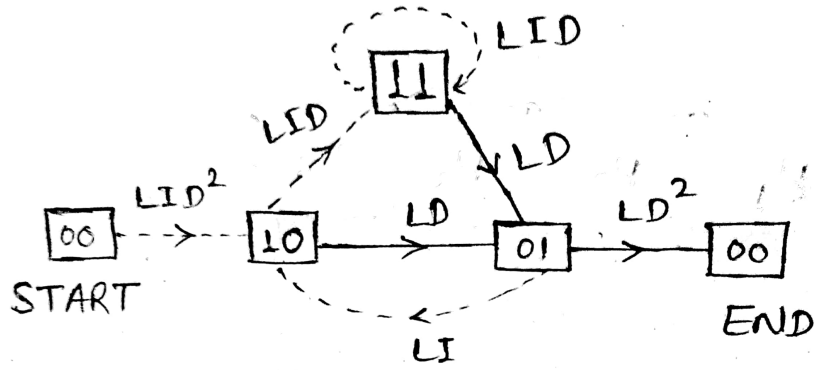


Suppose $\underline{c}_i = 0$ was transmitted. $\underline{y} \in H_0$ only.
 $\underline{y}^T (-1)^0 \geq \underline{y}^T (-1)^{c_i} \quad , \quad \underline{c}_i \neq 0$

Suppose we choose our decision regions such that
 $\underline{y} \in H_0$ iff $(\underline{y}^T) (-1)^0 \geq (\underline{y}^T) (-1)^{c_i} \quad \underline{c}_i \neq 0$

Next suppose $\underline{y} \in D_i$
 $\therefore \underline{y}^T (-1)^{c_i} \geq \underline{y}^T (-1)^{c_j} \quad , \quad c_j \neq c_i$
 $\therefore (\underline{y} \odot (-1)^{c_i})^T (-1)^0 \geq (\underline{y} \odot (-1)^{c_i})^T (-1)^{c_j + c_i}$

Modified State Diagram



$$\begin{bmatrix} A_{10} \\ A_{11} \\ A_{10} \end{bmatrix} = \begin{bmatrix} 0 & 0 & LI \\ LID & LID & 0 \\ LD & LD & 0 \end{bmatrix} \begin{bmatrix} A_{10} \\ A_{11} \\ A_{01} \end{bmatrix} + \begin{bmatrix} LID^2 \\ 0 \\ 0 \end{bmatrix}$$

$$\therefore \begin{bmatrix} 1 & 0 & -LI \\ -LID & 1 - LID & 0 \\ -LD & -LD & 1 \end{bmatrix} \begin{bmatrix} A_{10} \\ A_{11} \\ A_{01} \end{bmatrix} = \begin{bmatrix} LID^2 \\ 0 \\ 0 \end{bmatrix}$$

Let $T = \begin{bmatrix} 1 & 0 & -LI \\ -LID & 1 - LID & 0 \\ -LD & -LD & 1 \end{bmatrix}$

$$\begin{aligned} \det(T) &= 1(1 - LID) - LI(L^2ID^2 + LD(1 - LID)) \\ &= 1 - LID - L^2ID \\ &= 1 - LID(1 + L) \end{aligned}$$

$$\therefore A_{01} = \frac{(LID^2)(LD)}{1 - LID(1 + L)} \quad \text{where } \begin{vmatrix} -LID & 1 - LID \\ -LD & -LD \end{vmatrix} = LD$$

$$\therefore A_{END} = \frac{(LID^2)(LD)(LD^2)}{1 - LID(1 + L)} = \frac{L^3ID^5}{1 - LID(1 + L)}$$

Set $L=1, I=1$

$$\implies A_{END} = \frac{D^5}{1 - 2D} = D^5(1 + 2D + 4D^2 + \dots)$$

$$A_{END}(L, I, D) = \sum_{i,j,k} a_{i,j,k} L^i I^j D^k$$