Large-MIMO: A Technology Whose Time Has Come

A. Chockalingam and B. Sundar Rajan

(achockal,bsrajan@ece.iisc.ernet.in)

April 2010



Department of Electrical Communication Engineering

(http://www.ece.iisc.ernet.in/)

Indian Institute of Science

Bangalore – 560 012. INDIA



Why Multiple Antennas?

 N_t : No. of transmit antennas, N_r : No. of receive antennas

# Antennas	Error Probability (P_e)	Capacity (C) , bps/Hz
$N_t = N_r = 1$ (SISO)	$P_e \propto SNR^{-1}$	$C = \log(SNR)$
$N_t=1, N_r>1$ (SIMO)	$P_e \propto SNR^{-N_r}$	$C = \log(SNR)$
$N_t > 1, N_r > 1$ (MIMO)	$P_e \propto SNR^{-N_tN_r}$	$C = \min(N_t, N_r) \log(SNR)$
	$N_t N_r$: Diversity Gain	$\min(N_t,N_r):$ Spatial Mux Gain

• Large N_t , $N_r \rightarrow$ increased spectral efficiency

^[1] I. E. Telatar, Capacity of multi-antenna Gaussian channels, *European Trans. Telecommun.*, vol. 10, no. 6, pp. 585-595, November 1999.

^[2] G. J. Foschini and M. J. Gans, On limits of wireless communications in a fading environment when using multiple antennas, *Wireless Pers. Commun.*, vol. 6, pp. 311-335, March 1998.

Large-MIMO Approach

- Employ large number (several tens) of antennas at the Tx and Rx
- Achieve high spectral efficiencies (tens to hundreds of bps/Hz)
 - Data rate (bps) = Spectral efficiency (bps/Hz) \times Bandwidth (Hz)
 - e.g., 100 bps/Hz \implies 1 Gbps rate in just 10 MHz bandwidth
- Limitation in current MIMO standards
 - spectral efficiency: $\sim 10~{\rm bps/Hz}$ only
 - 2 to 4 transmit antennas
 - e.g., 750 Mbps in 80 MHz in 802.11n using 4 Tx antennas
 - do not exploit the potential of large spatial dimensions



- Need low-complexity detectors
- Channel estimation
 - Estimation of large number of channel coefficients

Dept. of ECE, IISc, Bangalore, April 2010



[3] Gregory Breit *et al*, 802.11ac Channel Modeling, doc. IEEE 802.11-09/0088r0, submission to Task Group TGac, 19 January 2009.

Dept. of ECE, IISc, Bangalore, April 2010

 64×64 MIMO Indoor Channel Sounding (5 GHz)



(a) 64-Antenna/RF hardware at 5 GHz



(b) LOS setup

[4] Jukka Koivunen, "Characterisation of MIMO propagation channel in multi-link scenarios," MS Thesis, Helsinki University of Technology, December 2007.





Linear Vector Channels

 Several communication systems can be characterized by the following linear vector channel model

 $\mathbf{y}_c = \mathbf{H}_c \mathbf{x}_c + \mathbf{n}_c$

 $\mathbf{x}_c \in \mathbb{C}^{d_t}$, $\mathbf{H}_c \in \mathbb{C}^{d_r imes d_t}$, $\mathbf{y}_c \in \mathbb{C}^{d_r}$, $\mathbf{n}_c \in \mathbb{C}^{d_r}$

- Examples
 - MIMO

* $d_t = N_t$, #Tx antennas; $d_r = N_r$, #Rx antennas; \mathbf{x}_c : Tx symbol vector * \mathbf{H}_c : Channel gain matrix; \mathbf{y}_c : Rx signal vector; \mathbf{n}_c : Noise vector

- Coding

* $d_t = k$, # Information bits; $d_r = n$, # Coded bits; \mathbf{x}_c : Information bit vector * \mathbf{H}_c : Generator matrix; \mathbf{y}_c : Rx signal vector; \mathbf{n}_c : Noise vector

- CDMA

* $d_t = d_r = K$, # users; \mathbf{x}_c : Tx. bit vector; \mathbf{H}_c : Cross correlation matrix, * \mathbf{y}_c : Rx signal vector; \mathbf{n}_c : Noise vector



Optimum Detection

- Let \mathbbm{A} be $M\mbox{-}\mathsf{PAM}$ or $M\mbox{-}\mathsf{QAM}$ (Two PAMs in quadrature)
 - *M*-PAM symbols take values from $\{A_m, m = 1, \dots, M\}$, $A_m = (2m 1 M)$

-
$$\mathbf{y}_c = \mathbf{y}_I + j\mathbf{y}_Q$$
, $\mathbf{x}_c = \mathbf{x}_I + j\mathbf{x}_Q$, $\mathbf{n}_c = \mathbf{n}_I + j\mathbf{n}_Q$, $\mathbf{H}_c = \mathbf{H}_I + j\mathbf{H}_Q$

• Convert (1) into a real-valued system model

$$\mathbf{y} = \mathbf{H}\mathbf{x} + \mathbf{n} \tag{2}$$

 $\mathbf{H} \in \mathbb{R}^{2d_r imes 2d_t}$, $\mathbf{y} \in \mathbb{R}^{2d_r}$, $\mathbf{x} \in \mathbb{R}^{2d_t}$, $\mathbf{n} \in \mathbb{R}^{2d_r}$

$$\mathbf{H} = \begin{pmatrix} \mathbf{H}_I & -\mathbf{H}_Q \\ \mathbf{H}_Q & \mathbf{H}_I \end{pmatrix}, \quad \mathbf{y} = [\mathbf{y}_I^T \ \mathbf{y}_Q^T]^T, \quad \mathbf{x} = [\mathbf{x}_I^T \ \mathbf{x}_Q^T]^T, \quad \mathbf{n} = [\mathbf{n}_I^T \ \mathbf{n}_Q^T]^T.$$

• ML solution

$$\mathbf{x}_{ML} = \frac{\arg\min}{\mathbf{x} \in \mathbb{S}} \|\mathbf{y} - \mathbf{H}\mathbf{x}\|^2 \mathbf{x}^T \mathbf{H}^T \mathbf{H}\mathbf{x} - 2\mathbf{y}^T \mathbf{H}\mathbf{x}, \quad (3)$$

S: $2d_t$ -dimensional signal space (Cartesian product of \mathbb{A}_1 to \mathbb{A}_{2d_t} ; \mathbb{A}_i : *M*-PAM signal set from which x_i takes values, $i = 1, \dots, 2d_t$). ML Complexity: Exponential in d_t

Optimum Detection

- Maximum a posteriori (MAP) solution
 - Consider square M-QAM
 - Each entry of ${\bf x}$ belongs to a $\sqrt{M}\mbox{-}{\rm PAM}$ constellation
 - Let $b_i^{(0)}, b_i^{(1)}, \cdots, b_i^{(q-1)}$ denote the $q = \log_2(\sqrt{M})$ constituent bits of x_i
 - x_i can be written as

$$x_i = \sum_{j=0}^{q-1} 2^j b_i^{(j)}, \quad i = 0, 1, \cdots, 2d_t - 1$$

– Let the bit vector $\mathbf{b} \in \{\pm 1\}^{2qd_t}$ be written as

$$\mathbf{b} \stackrel{\triangle}{=} \left[b_0^{(0)} \cdots b_0^{(q-1)} b_1^{(0)} \cdots b_1^{(q-1)} \cdots b_{2d_t-1}^{(0)} \cdots b_{2d_t-1}^{(q-1)} \right]^T$$

- Defining $\mathbf{c} \stackrel{\triangle}{=} [2^0 \, 2^1 \cdots 2^{q-1}]$, \mathbf{x} can be written as

$$\mathbf{x} = (\mathbf{I}_{2d_t} \otimes \mathbf{c})\mathbf{b}$$



Sub-optimum Solutions

• Matched filter (MF)

$$\mathbf{x}_{MF} = \mathbf{H}^T \mathbf{y}$$

• Zero-forcing (ZF) solution

 $\mathbf{x}_{ZF} = \mathbf{H}^{-1}\mathbf{y}$

• Minimum mean square error (MMSE) solution

 $\mathbf{x}_{MMSE} = (\mathbf{H} + \sigma^2 \mathbf{I})^{-1} \mathbf{y}$

 These suboptimum solution vectors can be used as initial vectors in search algorithms to improve performance further



LAS Algorithm

- Search for good solution vectors in the local neighborhood
- Neighborhood definition
 - Neighbors that differ in one coordinate
 - * e.g., Consider $\mathbb{A} = \{\pm 1\}; \quad \mathbf{x} = [-1, 1, 1, -1]$
 - $\ast\,$ 1-bit away neighbors of x:

 $\mathcal{N}_{1}(\mathbf{x}) = \left\{ [-1, 1, 1, 1], [-1, 1, -1], [-1, -1], [-1, -1], [1, 1, 1, -1] \right\}$

- Neighbors that differ in two coordinates
- 2-bit away neighbors of x:

$$\mathcal{N}_{2}(\mathbf{x}) = \left\{ \begin{bmatrix} -1, 1, -1, 1 \end{bmatrix}, \begin{bmatrix} -1, -1, -1, -1 \end{bmatrix}, \begin{bmatrix} 1, -1, 1, -1 \end{bmatrix}, \begin{bmatrix} 1, 1, 1, 1 \end{bmatrix}, \begin{bmatrix} 1, 1, 1, 1 \end{bmatrix}, \begin{bmatrix} 1, 1, 1, -1 \end{bmatrix}, \begin{bmatrix} 1, -1, 1, 1 \end{bmatrix} \right\}$$

• Choose best neighbor based on ML cost: $\phi(\tilde{\mathbf{x}}) = \tilde{\mathbf{x}}^T \mathbf{H}^T \mathbf{H} \tilde{\mathbf{x}} - 2 \mathbf{y}^T \mathbf{H} \tilde{\mathbf{x}}$















Complexity of 1-LAS in V-BLAST

- Consider $N_t = N_r$
- Total complexity comprises of 3 main parts
 - 1. Computing initial vector (e.g., ZF, MMSE): $O(N_t^2)$ per symbol
 - 2. Computing $\mathbf{H}^T \mathbf{H}$: $O(N_t^2)$ per symbol
 - 3. Search operation: $O(N_t)$ per symbol (through simulations)
- So, overall average per-symbol complexity: $O(N_t^2)$
- This low-complexity allows detection of V-BLAST signals in hundreds of spatial dimensions

Large # Dimensions: The Key

- Observation
 - In V-BLAST, LAS algorithm achieves near-ML performance, but only when the # antennas is in hundreds
 - hundreds of antennas may not be practical
- Note
 - LAS requires large # dimensions to perform well
 - but, all dimensions need not be in space alone
- Q1: Can large # dimensions be created with less # Tx antennas?
- A1: Yes. Use time dimension as well. Approach: Non-orthogonal STBCs
- Q2: Can LAS modified to work well for smaller (tens) dimensions?
- A2: Yes. Approach: Escape strategies from local minima

An Escape Strategy from Local Minima [7]

- Multistage LAS (M-LAS)
 - Start the algorithm as 1-LAS
 - On reaching the local minima,
 - * find 2-symbol away neighbors of the local minima
 - choose the best 2-symbol away neighbor if it has lesser cost than local minima
 - * run 1-LAS from this best neighbor till a local minima is reached
 - Expect better performance. Complexity is increased a little, but not by an order
 - Escape strategy with 3-symbol away neighborhood on reaching local minima
- Another promising strategy is reactive tabu search

^[7] S. K. Mohammed, A. Chockalingam, B. S. Rajan, A low-complexity near-ML performance achieving algorithm for large MIMO detection, *IEEE ISIT*²⁰⁰⁸, Toronto, June 2008.



Space-Time Block Codes

- Provide redundancy across space and time
- Goal of space-time coding
 - Achieve the maximum Tx-diversity of N_t (i.e., full-diversity), high rate, decoding at low-complexity
- An STBC is usually represented by a $p imes n_t$ matrix
 - rows: time slots; p: # time slots
 - columns: Tx. antennas; n_t : # Tx. antennas

$$\mathbf{X} = \begin{bmatrix} s_{11} & s_{12} & \cdot & s_{1n_t} \\ s_{21} & s_{22} & \cdot & s_{2n_t} \\ \cdot & \cdot & \cdot & \cdot \\ s_{p1} & s_{42} & \cdot & s_{pn_t} \end{bmatrix}$$

• s_{ij} denotes the complex number transmitted in the *i*th time slot on the *j*th Tx antenna

Space-Time Block Codes

- Rate of an STBC, $r = \frac{k}{p}$
 - k: number of information symbols sent in one STBC
 - p: number of time slots in one STBC
 - * Higher rate means more information carried by the code
- \bullet A matrix ${\bf X}$ is said to be a Orthogonal STBC if

$$\mathbf{X}^{H}\mathbf{X} = (|x_{1}|^{2} + |x_{2}|^{2} + \dots + |x_{k}|^{2})\mathbf{I}_{n_{t}}$$

- Elements of X are linear combinations of x_1, \dots, x_k and their conjugates
- x_1, x_2, \cdots, x_k are information symbols
- 2-Tx Antennas Codes (2×2 Alamouti Code)

$$\mathbf{X} = \begin{bmatrix} x_1 & x_2 \\ -x_2^* & x_1^* \end{bmatrix}, \quad k = 2, p = 2, r = 1, \text{ orthogonal STBC}$$

Linear-Complexity Decoding of OSTBCs

- Consider Alamouti code with $n_t = 2$, $n_r = 1$
- Received signal in *i*th slot, y_i , i = 1, 2, is

$$y_1 = h_1 x_1 + h_2 x_2 + n_1$$

$$y_2 = -h_1 x_2^* + h_2 x_1^* + n_2$$

- ML decoding amounts to
 - computing

$$\tilde{x}_1 = y_1 h_1^* + y_2^* h_2$$

 $\tilde{x}_2 = y_1 h_2^* - y_2^* h_1$

- decoding x_1 by finding the symbol in the constellation that is closest to \tilde{x}_1
- and decoding x_2 by finding the symbol that is closest to \tilde{x}_2
- This decoding feature is called Single-Symbol Decodability (SSD)



decoding complexity

2

Non-Orthogonal STBCs

• Golden code [8] (2×2 non-orthogonal STBC)

$$\mathbf{X} = \begin{bmatrix} x_1 + \tau x_2 & x_3 + \tau x_4 \\ i(x_3 + \mu x_4) & x_1 + \mu x_2 \end{bmatrix}, \quad k = 4, p = 2, r =$$

where $\tau = \frac{1+\sqrt{5}}{2}$ and $\mu = \frac{1-\sqrt{5}}{2}$

- Features
 - Information Losslessness (ILL)
 - Full Diversity (FD)
 - Coding Gain (CG)
- 'Perfect codes' [9] achieve all the above three features
 - Golden code is a perfect code

^[8] J.-C. Belfiore, G. Rekaya, and E. Viterbo, "The golden code: A 2 × 2 full-rate space-time code with non-vanishing determinants," *IEEE Trans. on Information Theory*, vol. 51, no. 4, pp. 1432-1436, April 2005.

^[9] F. E. Oggier, G. Rekaya, J.-C. Belfiore, and E. Viterbo, "Perfect space-time block codes," *IEEE Trans. on Information Theory*, vol. 52, no. 9, pp. 3885-3902, September 2006.



Linear Vector Channel Model for NO-STBC

- (n, p, k) STBC is a matrix $\mathbf{X}_c \in \mathbb{C}^{n \times p}$, n: # time slots, p: # tx antennas, k: # data symbols in one STBC; $(n = p \text{ and } k = n^2 \text{ for NO-STBC from CDA})$
- Received space-time signal matrix

$$\mathbf{Y}_c = \mathbf{H}_c \mathbf{X}_c + \mathbf{N}_c,$$

• Consider linear dispersion STBCs where \mathbf{X}_c can be written in the form

$$\mathbf{X}_c = \sum_{i=1}^k x_c^{(i)} \mathbf{A}_c^{(i)}$$

where $\mathbf{A}_c^{(i)} \in \mathbb{C}^{N_t imes p}$ is the weight matrix corresponding to data symbol $x_c^{(i)}$

• Applying vec(.) operation

$$vec(\mathbf{Y}_{c}) = \sum_{i=1}^{k} x_{c}^{(i)} vec(\mathbf{H}_{c}\mathbf{A}_{c}^{(i)}) + vec(\mathbf{N}_{c})$$
$$= \sum_{i=1}^{k} x_{c}^{(i)}(\mathbf{I}_{p \times p} \otimes \mathbf{H}_{c}) vec(\mathbf{A}_{c}^{(i)}) + vec(\mathbf{N}_{c})$$

Linear Vector Channel Model for NO-STBC

- Define $\mathbf{y}_c \stackrel{\Delta}{=} vec\left(\mathbf{Y}_c\right) \in \mathbb{C}^{N_r p}$, $\widehat{\mathbf{H}}_c \stackrel{\Delta}{=} \left(\mathbf{I} \otimes \mathbf{H}_c\right) \in \mathbb{C}^{N_r p \times N_t p}$, $\mathbf{a}_c^{(i)} \stackrel{\Delta}{=} vec\left(\mathbf{A}_c^{(i)}\right) \in \mathbb{C}^{N_t p}$, $\mathbf{n}_c \stackrel{\Delta}{=} vec\left(\mathbf{N}_c\right) \in \mathbb{C}^{N_r p}$
- System model can then be written in vector form as

$$\mathbf{y}_{c} = \sum_{i=1}^{k} x_{c}^{(i)} \left(\widehat{\mathbf{H}}_{c} \mathbf{a}_{c}^{(i)} \right) + \mathbf{n}_{c}$$
$$= \widetilde{\mathbf{H}}_{c} \mathbf{x}_{c} + \mathbf{n}_{c}$$
(4)

 $\widetilde{\mathbf{H}}_c \in \mathbb{C}^{N_r p \times k}$, whose *i*th column is $\widehat{\mathbf{H}}_c \ \mathbf{a}_c^{(i)}$, $i = 1, \cdots, k$ $\mathbf{x}_c \in \mathbb{C}^k$, whose *i*th entry is the data symbol $x_c^{(i)}$

- Convert the complex system model in (4) into real system model as before
- Apply LAS algorithm on the resulting real system model


Special Issue on Managing Complexity in Multiuser MIMO Systems, vol. 3, no. 6, pp. 958-974, December 2009.

Comparison with Other Architectures/Detectors [11]No. $Nimo Architecture/Detector Combinations
(fixed <math>N_t = N_r = 16$ and 32 bps/HzComplexity
(in # real operations
per bit) at 5×10^{-2} SNR required
to achieve 5×10^{-2}
uncoded BER

	(fixed $N_t=N_r=16$ and 32 bps/Hz	per bit) at $5 imes 10^{-2}$	uncoded BER
	for all combinations)	uncoded BER	
	16 imes16 ILL-only CDA STBC (rate-16),		
i)	4-QAM and 1-LAS detection	$3.473 imes \mathbf{10^3}$	6.8 dB
	(Proposed scheme [11])		
ii)	16 imes16 ILL-only CDA STBC (rate-16),		
	4-QAM and ISIC algorithm [Choi, Cioffi]	1.187×10^5	11.3 dB
iii)	Four $4 imes 4$ stacked rate-1 QOSTBCs,		
	256-QAM and IC algorithm [Jafarkhani]	5.54×10^6	24 dB
iv)	Eight 2×2 stacked rate-1 Alamouti codes,		
	16-QAM and IC algorithm [Jafarkhani]	8.719×10^3	17 dB
v)	16 imes16 V-BLAST (rate-16) scheme,		
	4-QAM and sphere decoding	4.66×10^4	7 dB
vi)	16 imes16 V-BLAST (rate-16) scheme,		
	4-QAM and V-BLAST detector (ZF-SIC)	$1.75 imes 10^4$	13 dB









MIMO Capacity with Estimated CSIR

Hassibi-Hochwald (H-H) bound [13] on capacity with estimated CSIR:

$$C \geq \frac{T-\tau}{T} \mathbb{E} \left[\mathsf{logdet} \left(\mathbf{I}_{N_t} + \frac{\gamma^2 \beta_d \beta_p \tau}{N_t (1+\gamma \beta_d) + \gamma \beta_p \tau} \; \frac{\hat{\mathbf{H}}_c \hat{\mathbf{H}}_c^H}{N_t \sigma_{\hat{\mathbf{H}}_c}^2} \right) \right]$$



Figur4: H-H capacity bound [13] for 1P+8D ($T = 144, \tau = 16, \beta_p = \beta_d = 1$) and 1P+1D ($T = 32, \tau = 16, \beta_p = \beta_d = 1$) training for a 16×16 MIMO channel.

[13] B. Hassibi and B. M. Hochwald, "How much training is needed in multiple-antenna wireless links?," *IEEE Trans. on Information Theory*, vol. 49, no. 4, pp. 951-963, April 2003.



Figur**5**: Capacity as a function of N_t with SNR = 18 dB and $N_r = 12$. For a given N_r , SNR (γ) , and

coherence time (T), there is an optimum N_t [13].





^[14] N. Srinidhi, S. K. Mohammed, A. Chockalingam, and B. S. Rajan, Low-Complexity Near-ML Decoding of Large Non-Orthogonal STBCs using Reactive Tabu Search, *IEEE ISIT'2009*, Seoul, June 2009.

Dept. of ECE, IISc, Bangalore, April 2010

^[15] S. K. Mohammed, A. Chockalingam, B. S. Rajan, Low-complexity near-MAP decoding of large non-orthogonal STBCs using PDA, *IEEE ISIT'2009*, Seoul, June 2009.

^[16] S. Madhekar, P. Som, A. Chockalingam, B. S. Rajan, Belief Propagation Based Decoding of Large Non-Orthogonal STBCs, *IEEE ISIT'2009*, Seoul, June 2009.

^[17] P. Som, T. Datta, A. Chockalingam, B. S. Rajan, Improved Large-MIMO Detection using Damped Belief Propagation, *IEEE ITW*'2010, Cairo, January 2010.

Reactive Tabu Search

- Another iterative local search algorithm
 - A metaheuristics algorithm
 - cannot guarantee optimal solution, but generally gives near optimal solution
- Uses 'tabu' mechanism to escape from local minima or cycles
 - Certain vectors are prohibited (made tabu) from becoming solution vectors for certain number of iterations (called *tabu period*) depending on the search path
 - This is meant to ensure efficient exploration of the search space
- The reactive part adapts the tabu period





















[18] N. Srinidhi, S. K. Mohammed, A. Chockalingam, B. S. Rajan, *Near-ML Signal Detection in Large-Dimension Linear Vector Channels Using Reactive Tabu Search*, Online arXiv:0911.4640v1 [cs.IT] 24 Nov 2009.







* Turbo equalization (Yin et al 2004)



 \mathbf{h}_t : *t*th column of **H**

PDA Based Large-MIMO Detection [15]

- Define $p_i^{j+} \stackrel{\triangle}{=} P(b_i^{(j)} = +1)$ and $p_i^{j-} \stackrel{\triangle}{=} P(b_i^{(j)} = -1)$
- To compute $eta_i^{(j)}$, approximate the distribution of $\widetilde{\mathbf{n}}$ to be Gaussian
- \bullet Mean of \boldsymbol{y}

$$\boldsymbol{\mu}_{i}^{j+} \stackrel{\Delta}{=} \mathbb{E}(\mathbf{y}|b_{i}^{(j)} = +1) = \mathbf{h}_{qi+j} + \sum_{l=0}^{2k-1} \sum_{\substack{m=0\\m \neq q(i-l)+j}}^{q-1} \mathbf{h}_{ql+m}(2p_{l}^{m+j} - 1)$$

$$\boldsymbol{\mu}_{i}^{j-} \stackrel{\Delta}{=} \mathbb{E}(\mathbf{y}|b_{i}^{(j)}=-1) = \boldsymbol{\mu}_{i}^{j+}-2\mathbf{h}_{qi+j}$$

 \bullet Covariance of \boldsymbol{y}

$$\mathbf{C}_{i}^{j} = \sigma^{2} \mathbf{I}_{2N_{r}p} + \sum_{l=0}^{2k-1} \sum_{\substack{m=0\\m \neq q(i-l)+j}}^{q-1} \mathbf{h}_{ql+m} \, \mathbf{h}_{ql+m}^{T} \, 4p_{l}^{m+} (1-p_{l}^{m+})$$

PDA Based Large-MIMO Detection [15]

• Using $\mu_i^{j\pm}$ and \mathbf{C}_i^j , $P(\mathbf{y}|b_i^{(j)}=\pm 1)$ can be written as

$$P(\mathbf{y}|b_i^{(j)} = \pm 1) = \frac{e^{-(\mathbf{y} - \boldsymbol{\mu}_i^{j\pm})^T (\mathbf{C}_i^j)^{-1} (\mathbf{y} - \boldsymbol{\mu}_i^{j\pm})}}{(2\pi)^{N_r p} |\mathbf{C}_i^j|^{\frac{1}{2}}}$$

- Using (5), β_i^j can be written as $\beta_i^j \beta_i^j = e^{-\left((\mathbf{y} - \boldsymbol{\mu}_i^{j+})^T (\mathbf{C}_i^j)^{-1} (\mathbf{y} - \boldsymbol{\mu}_i^{j+}) - (\mathbf{y} - \boldsymbol{\mu}_i^{j-})^T (\mathbf{C}_i^j)^{-1} (\mathbf{y} - \boldsymbol{\mu}_i^{j-})\right)}$
- Compute $\Lambda_i^{(j)}$ using $\alpha_i^{(j)}$ and $\beta_i^{(j)}$
- Update the statistics of $b_i^{(j)}$ as

$$P(b_i^{(j)} = +1|\mathbf{y}) = \frac{\Lambda_i^{(j)}}{1 + \Lambda_i^{(j)}}, \quad P(b_i^{(j)} = -1|\mathbf{y}) = \frac{1}{1 + \Lambda_i^{(j)}}$$

This completes one iteration of the algorithm

PDA Based Large-MIMO Detection [15]

- Updated values of $P(b_i^{(j)} = +1|\mathbf{y})$ and $P(b_i^{(j)} = -1|\mathbf{y})$ for all i, j are fed back as a priori probabilities to the next iteration
- Algorithm terminates after a certain number of iterations
- At the end of the last iteration,
 - decide $\widehat{b}_i^{(j)}$ as +1 if $\Lambda_i^{(j)} \ge 1$, and -1 otherwise
- In coded systems
 - feed $\Lambda_i^{(j)}$'s as soft inputs to the decoder









- Damping significantly improves performance
- Order of per-symbol complexity: $O(N_t^2)$



Large-MIMO Detection using **BP on FGs** [17]

• With x_i 's $\in \{\pm 1\}$, the log-likelihood ratio (LLR) of x_k at observation node i, denoted by Λ_i^k , is

$$\Lambda_{i}^{k} = \log \frac{p(y_{i} | \mathbf{H}, x_{k} = 1)}{p(y_{i} | \mathbf{H}, x_{k} = -1)} = \frac{2}{\sigma_{z_{ik}}^{2}} \Re \left(h_{ik}^{*} (y_{i} - \mu_{z_{ik}}) \right)$$

- LLR values computed at observation nodes are passed to variable nodes.
- Using these LLRs, variable nodes compute the probabilities

$$p_i^{k+} \stackrel{\triangle}{=} p_i(x_k = +1|\mathbf{y}) = \frac{\exp(\sum_{l \neq i} \Lambda_l^{\kappa})}{1 + \exp(\sum_{l \neq i} \Lambda_l^{k})}$$

and pass them back to the observation nodes.

- This message passing is carried out for a certain number of iterations.
- At the end, x_k is detected as

$$\widehat{x}_k = \operatorname{sgn}\left(\sum_{i=1}^{2N_r} \Lambda_i^k\right)$$

9 M



Figure: Message passing between variable nodes and observation nodes.

Dept. of ECE, IISc, Bangalore, April 2010




Concluding Remarks

- Low-complexity detection
 - critical enabling technology for large-MIMO
 - no more a bottleneck
- Large-MIMO systems can be implemented
- Large-MIMO approach scores high on spectral efficiency and operating SNR compared to other approaches (e.g., increasing QAM size)
- Standardization efforts can consider reaping the benefits of large-MIMO in their evolution



